# ${\rm Mod}\;(2p+1){\rm -orientation}\;{\rm of}\;{\rm Graphs}$

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#### **Motivation: Score Sequence**

Suppose that *n* teams v<sub>1</sub>, v<sub>2</sub>, ··· , v<sub>n</sub> are in a tournament. A non negative integer sequence s<sub>1</sub>, s<sub>2</sub>, ··· , s<sub>n</sub> is a score sequence if it is a possible outcome that there is a permutation π on {1, 2, ··· , n} such that for each *i*, the team v<sub>i</sub> will win exactly s<sub>π(i)</sub> games in the tournament.

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- Landau's score sequence problem: Given a non negative integer sequence  $s_1, s_2, \dots, s_n$ , how do you know if this sequence is a score sequence?

#### **Score Sequence**

Theorem 1 (Landau 1953) Let  $n \ge 1$  be an integer. A nondecreasing sequence  $(s_1, s_2, \dots, s_n)$  of nonnegative integers is a score sequence if and only if

$$\sum_{i=1}^{k} s_i \ge \begin{pmatrix} k \\ 2 \end{pmatrix}, \forall k \text{ with } 1 \le k \le n,$$

where equality holds if and only if n = k.

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- $\forall v \in V(G)$ ,  $d_D^+(v) = \#$  of edges directed from v(out-degree),  $d_D^-(v) = \#$  of edges directed into v(in-degree).

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- $\forall v \in V(G)$ ,  $d_D^+(v) = \#$  of edges directed from v(out-degree),  $d_D^-(v) = \#$  of edges directed into v(in-degree).
- Formulation: A function  $c: V(G) \mapsto \mathbb{Z}$  represents a score sequence if and only if the complete graph  $K_n$  has an orientation D such that at each v,  $d_D^+(v) = c(v)$ .

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Problem: Suppose that *G* is a graph (not necessarily complete). Given an integer valued function  $c: V(G) \mapsto \mathbb{Z}$ , can we find an orientation D(G) such that at each v,  $d_D^+(v) = c(v)$ ?

# Orientation

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- Get an orientation D on E(G): each edge e is oriented so that the face with greater color number is on the right side of the oriented edge  $e \Leftrightarrow$  the greater – the smaller = 1.

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The resulted orientation D satisfied that  $\forall v$ ,  $d_D^+(v) - d_D^-(v) \equiv 0 \pmod{3}$ , (called a mod 3-orientation of G).

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### **Conjectures**

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#### **Conjectures**

- Conjecture (Tutte 1960) Every 4-edge-connected graph has a mod 3-orientation.
- Fact: There exists a 3-edge-connected graph that does not have a mod 3-orientation. (For example, K<sub>4</sub>).
- Conjecture (Jaeger 1988) There exists an integer k ≥ 4 such that every k-edge-connected graph has a mod 3-orientation.

Tutte's mod 3-orientation problem: (also known as the 3-flow Problem) Given a graph G, can we find an orientation D(G) such that at v,  $d_D^+(v) - d_D^-(v) \equiv 0$  (mod 3)? Or: does G have a mod 3-orientation? (NP-complete even within planar graphs).

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- Problem: Let *A* be an abelian (additive) group. Given an *A*-valued function  $b: V(G) \mapsto A$ , can we find an orientation D(G) such that at v,  $d_D^+(v) - d_D^-(v) = b(v)$  in *A*?

G: = a graph, with vertex set

$$V = V(G) = \{v_1, v_2, \cdots, v_n\},\$$

and edge set

$$E = E(G) = \{e_1, e_2, \cdots, e_m\}.$$

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 $\square$   $D = (d_{ij})_{n \times m}$ : = vertex-edge incidence matrix, where

$$d_{ij} = \begin{cases} 1 & \text{if } e_j \text{ is oriented away from } v_i \\ -1 & \text{if } e_j \text{ is oriented into } v_i \\ 0 & \text{otherwise} \end{cases}$$

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The *i*th row (component) of D1 equals the net out degree of  $v_i$ .

## **Non-Homogeneous System Formulation**

■ A: = an abelian (additive) group with identity 0, and with  $|A| \ge 3$ .

# **Description of an Orientation**

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- Suppose that G has an orientation D (whose adjacency matrix is also denoted by D).
- Suppose we reverse the orientation of  $e_i$  in D to obtain a new orientation D'. If  $\exists f : E \mapsto \{1, -1\}$  such that  $f^{-1}(-1) = \{e_i\}$ , then  $D\mathbf{1} = D'f$ . Therefore, with an arbitrarily given orientation of G, any other orientation of G can be viewed as a function  $\exists f : E \mapsto \{1, -1\}$ .

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- Assume that G has a fixed orientation. If  $\exists f: E \mapsto \{1, -1\}$  such that Df = 0 over  $\mathbb{Z}_{2p+1}$ , then we say that G has a mod (2p + 1)-orientation.
- For an undirected graph G, whether G has a mod (2p+1)-orientation or not is independent of the choice of the orientation of G.

A necessary Condition: If  $\forall b : V(G) \mapsto \mathbb{Z}_{2p+1}$ ,  $\exists f : E \mapsto \{1, -1\}$  such that Df = b over  $\mathbb{Z}_{2p+1}$ , then

$$\sum_{v \in V(G)} b(v) \equiv 0 \pmod{2p+1}.$$

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**Proof:** View f as an orientation D. Then

$$\sum_{v \in V(G)} b(v) = \sum_{v \in V(G)} [d_D^+(v) - d_D^-(v)].$$

Each edge is counted on the right hand side exactly twice, once positive and once negative.

#### Graphs That Are Mod (2p + 1)-Contractible

If  $\forall b : V(G) \mapsto \mathbb{Z}_{2p+1}$  with  $\sum_{v \in V(G)} b(v) \equiv 0 \pmod{2p+1}$ ,  $\exists f : E \mapsto \{1, -1\}$  such that Df = b over  $\mathbb{Z}_{2p+1}$ , then we say that G is mod (2p+1)-contractible.

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- Every mod (2p + 1)-contractible graph has a mod (2p + 1)-orientation.
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- Every mod (2p+1)-contractible graph has a mod (2p+1)-orientation.
- For an undirected graph G, whether G is mod (2p+1)-contractible or not is independent of the choice of the orientation of G.
- **Examples:**  $2K_2$  and  $K_5$  are mod 3-contractible.

- Proposition: If H is a subgraph of G, and if H is mod (2p+1)-contractible, then the following are equivalent:
  (i) G has mod (2p+1)-orientation,
  - (ii) the contraction G/H has mod (2p + 1)-orientation.

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Proof: Note that mod (2p + 1)-contractible implies mod (2p + 1)-orientation. This follows from Theorem 2 of next page.

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- Theorem 1 (Z.-H. Chen, H.-J. Lai, H. Y. Lai 2001) Suppose that  $H \subseteq G$  and that H is mod 3-contractible. Then G is mod 3-contractible if and only if G/H is mod 3-contractible.

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- Theorem 2 (Lai, Shao, Wu and Zhou 2006) Suppose that  $H \subseteq G$  and that H is mod (2p + 1)-contractible. Then G is mod (2p + 1)-contractible if and only if G/His mod (2p + 1)-contractible.

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- Theorem 7 (Steinburg and D. H. Younger 1989, Thomassen 1994) Every 4-edge-connected projective planar graph has a mod 3-orientation.

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- Theorem 7 (Steinburg and D. H. Younger 1989, Thomassen 1994) Every 4-edge-connected projective planar graph has a mod 3-orientation.
- Theorem 8 (Lai and Zhang 1992) Every 4 log<sub>2</sub>(|V(G)|)-edge-connected graph has a mod 3-orientation.

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- $\delta_D^-(S)$  = edges oriented into *S*.

Necessity: Let  $c: V(G) \mapsto \mathbb{Z}$ . If G has an orientation D such that  $d_D^+(v) = c(v), \forall v \in V(G)$ , then  $\forall S \subseteq V(G)$ 

$$|E(S)| \le \sum_{v \in S} c(v) \le |E(S)| + |\partial_G(S)|.$$

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A function  $c: V(G) \mapsto \mathbb{Z}$  satisfying inequality above will be called a feasible function of G.

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$$|E(S)| = \sum_{v \in S} d_D^+(v) - |\delta_D^+(S)| \le \sum_{v \in S} d_D^+(v)$$
$$\le \sum_{v \in S} d_D^+(v) + |\delta_D^-(S)| = |E(S)| + |\partial_G(S)|.$$

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• (i) *G* has an orientation *D* such that  $d_D^+(v) = c(v), \forall v \in V(G)$ .

• (ii) c is a feasible function of G. That is,  $\forall S \subseteq V(G)$ 

$$|E(S)| \le \sum_{v \in S} c(v) \le |E(S)| + |\partial_G(S)|.$$

Corollary (Landau 1953) Let  $n \ge 1$  be an integer. A nondecreasing sequence  $(s_1, s_2, \dots, s_n)$  of nonnegative integers is a score sequence if and only if

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**Proof** Directly verify that the function *s* is feasible.

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• (ii)  $\forall b : V(G) \mapsto \mathbf{Z}$  satisfying both

$$\sum_{v \in V(G)} b(v) \equiv 0 \pmod{2p+1}$$

and

$$b(v) \equiv d_G(v) \pmod{2}, \forall v \in V(G),$$

G has an orientation D such that  $d_D^+(v) - d_D^-(v) \equiv b(v)$ (mod 2p + 1),  $\forall v \in V(G)$ .

Corollary (Lai, Shao, Wu and Zhou, 2006) Let G be a (4p+1)-regular graph. Then G has a mod (2p+1)-orientation if and only if V(G) has a partition  $(V^+, V^-)$  such that  $\forall U \subseteq V(G)$ ,

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Proof of the "only if" part Suppose *G* has a mod 3-orientation *D*. Since *G* is 5-regular,  $\forall v \in V(G)$ , either  $d_D^+(v) = 4p$  or  $d_D^+(v) = 1$ . Define  $V^+ = \{v \in V(D) : d_D^+(v) = 4p\}$  and  $V^- = V(D) - V^+$ . Apply Theorem 10.

Proof of the "if" part Define a map  $b : V(G) \mapsto \mathbb{Z}$ satisfying  $b(V^+) = \{2p+1\}$  and  $b(V^-) = \{-2p-1\}$ . Since *G* is (4p+1)-regular,  $\forall v \in V(G)$ ,  $b(v) \equiv d_G(v)$ (mod 2).

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- When U = V(G), we have  $|V^+| = |V^-|$ , and so  $\sum_{v \in V(G)} b(v) = 0$ .

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Apply Theorem 10.

Corollary (Da Silva and Dahad, 2005) Let G be a 5-regular graph. Then G has a mod 3-orientation if and only if V(G) has a partition  $(V^+, V^-)$  such that  $\forall U \subseteq V(G)$ ,

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**Proof** Let p = 1.
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(i) V(G) is a disjoint union V(G) = V<sub>1</sub>UV<sub>2</sub> with |V<sub>1</sub>| = k, |V<sub>2</sub>| = n - k, and

$$\left\lceil \frac{|E(V_1, V_2)| + 1}{k} \right\rceil + \left\lceil \frac{|E(V_1, V_2)| + 1}{n - k} \right\rceil \le 4p + 2.$$

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$$\left\lceil \frac{|E(V_1, V_2)| + 1}{k} \right\rceil + \left\lceil \frac{|E(V_1, V_2)| + 1}{n - k} \right\rceil \le 4p + 2.$$

• (ii) V(G) is a disjoint union  $V(G) = V_1 \bigcup V_2$  with  $|V_1| = k$ ,  $|V_2| = n - k$ , and

$$|E(V_1, V_2)| \le \frac{(4p+2)k(n-k)}{n}$$

Example: For any positive  $p \in \mathbb{Z}$ ,  $K_{4p+1}$  is mod (2p+1)-contractible.

- Example: For any positive  $p \in \mathbb{Z}$ ,  $K_{4p+1}$  is mod (2p+1)-contractible.
- Proof: n = 4p + 1. By Theorem 11,  $V(K_n)$  can be partitioned into two subsets  $V_1$  and  $V_2$  with  $|V_1| = k$  and  $|V_2| = n - k$  satisfying inequality Theorem 11(ii). Since  $|E(V_1, V_2)| = k(n - k)$ , we have

$$\begin{bmatrix} |E(V_1, V_2)| + 1 \\ k \end{bmatrix} + \begin{bmatrix} |E(V_1, V_2)| + 1 \\ 4p - k \end{bmatrix}$$
$$= \begin{bmatrix} \frac{k(n-k) + 1}{k} \end{bmatrix} + \begin{bmatrix} \frac{|k(n-k) + 1}{n-k} \end{bmatrix}$$
$$= (n-k+1) + (k+1) = n+2 > 4p+2.$$

Theorem 12 (Lai, Shao, Wu and Zhou 2006) Let n, pbe positive integers, and let  $f(n) = \frac{(2p+1)n \log_2(n)}{2}$ be a function. If *G* is a graph with *n* vertices and if  $|E(G)| \ge f(n)$ , then *G* has a subgraph *H* with  $E(H) \ne \emptyset$  which is mod (2p+1)-contractible.

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- Theorem 13 (Lai, Shao, Wu and Zhou 2006) Let G be a graph with n vertices. If G is (2p+1) log<sub>2</sub>(n)-edge-connected, then G is mod (2p+1)-contractible.
- Proof: Use connectivity to count the number of edges and use Theorem 12 to find a contractible subgraph.

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- Theorem 8 (Lai and Zhang 1992) Let G be a graph with n vertices. If G is  $4 \log_2(n)$ -edge-connected, then G has a mod 3-orientation.



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A connected loopless graph with 3 edges and a vertex of degree 3 is called a generalized claw.

A graph G with |E(G)| ≡ 0 (mod 3) has a claw-decomposition if E(G) can be partitioned into disjoint unions E(G) = X<sub>1</sub> ∪ X<sub>2</sub> ∪ · · · ∪ X<sub>k</sub> such that for each i with 1 ≤ i ≤ k, G[X<sub>i</sub>] is a generalized claw.

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- Theorem 14 (Barat and Thomassen 2004) If there exists an integer k such that every k-edge-connected graph G with  $|E(G)| \equiv 0 \pmod{3}$  has a claw-decomposition, then every k-edge-connected graph G has a mod 3-orientation.

Conjecture 15 (Barat and Thomassen 2004) Every 4-edge-connected simple planar graph *G* with  $|E(G)| \equiv 0 \pmod{3}$  has a claw-decomposition.

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Question in our minds: How do we approach this conjecture?

A connected loopless graph with 2p + 1 edges and a vertex of degree 2p + 1 is called a generalized  $K_{1,2p+1}$ .

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A graph G with  $|E(G)| \equiv 0 \pmod{2p+1}$  has a  $K_{1,2p+1}$ -decomposition if E(G) can be partitioned into disjoint unions  $E(G) = X_1 \cup X_2 \cup \cdots \cup X_k$  such that for each i with  $1 \leq i \leq k$ ,  $G[X_i]$  is a generalized  $K_{1,2p+1}$ .

Theorem 16 (Lai, Shao, Wu and Zhou 2006) Fix k > 0. The every k-edge-connected (planar) graph G is mod (2p+1)-contractible if and only if every k-edge-connected (planar) graph G with  $|E(G)| \equiv 0$ (mod 2p + 1) has a  $K_{1,2p+1}$ -decomposition.

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- **Proof** Apply Theorems 16 and 17 when p = 1.

Question 18 Is there an integer k such that every k-edge-connected planar graph G with  $|E(G)| \equiv 0$  (mod 3) has a  $K_{1,3}$ -decomposition?

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**Proof** Apply Theorems 16 and 19 when p = 1.

