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# Note Hamiltonian *s*-properties and eigenvalues of *k*-connected graphs

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### A R T I C L E I N F O

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## ABSTRACT

Chvátal and Erdös (1972) [5] proved that, for a *k*-connected graph *G*, if the stability number  $\alpha(G) \le k - s$ , then *G* is Hamilton-connected (*s* = 1) or Hamiltonian (*s* = 0) or traceable (*s* = -1). Motivated by the result, we focus on tight sufficient spectral conditions for *k*-connected graphs to possess Hamiltonian *s*-properties. We say that a graph possesses Hamiltonian *s*-properties, which means that the graph is Hamilton-connected if *s* = 1, Hamiltonian if *s* = 0, and traceable if *s* = -1.

For a real number  $a \ge 0$ , and for a *k*-connected graph *G* with order *n*, degree diagonal matrix D(G) and adjacency matrix A(G), we have identified best possible upper bounds for the spectral radius  $\lambda_1(aD(\Gamma) + A(\Gamma))$ , where  $\Gamma$  is either *G* or the complement of *G*, to warrant that *G* possesses Hamiltonian *s*-properties. Sufficient conditions for a graph *G* to possess Hamiltonian *s*-properties in terms of upper bounds for the Laplacian spectral radius as well as lower bounds of the algebraic connectivity of *G* are also obtained. Other best possible spectral conditions for Hamiltonian *s*-properties are also discussed.

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#### 1. Introduction

We consider simple, undirected and connected graphs with undefined terms and notation reference to [3]. As in [3],  $\overline{G}$ ,  $\alpha(G)$ ,  $\kappa(G)$ ,  $\delta(G)$  and  $d(\nu)$  denote the complement, the **stability number** (also call the **independence number**), **the connectivity**, the minimum degree of a graph *G* and the degree of vertex  $\nu$  in *G*, respectively. Let  $K_{a,b}$  denote complete bipartite graphs on *n* vertices, where a + b = n.

A well-known result of Whitney [28] states that  $\kappa(G) \leq \delta(G)$  for any graph *G*. A graph *G* is *k*-connected if  $\kappa(G) \geq k$ . A cycle (path, respectively) passing through all the vertices of a graph is called a Hamilton cycle (Hamilton path, respectively). A graph *G* is called **Hamilton-connected** if every two vertices of *G* are connected by a Hamilton path. A graph containing a Hamilton cycle is called a **Hamiltonian graph**. It is known that all Hamilton-connected graphs are Hamiltonian. A graph containing a Hamilton path is said to be **traceable**.

For any graph *G* with the adjacency matrix A(G) and the diagonal degree matrix D(G), we define  $\lambda_1(aD(G) + bA(G))$  to be the spectral radius of aD(G) + bA(G), where  $a \ge 0$  and b > 0 are two real numbers. When a = 0 and b = 1, the value  $\lambda_1(aD(G) + bA(G))$  is called the **spectral radius** of a graph *G*, denoted by  $\rho_1(G)$ . If a = 1 and b = 1, then  $\lambda_1(aD(G) + bA(G))$  is called the **graph** *G*, and is denoted by  $q_1(G)$ . Furthermore, for a real number  $\alpha \in [0, 1)$ ,  $\lambda_1(A_\alpha(G)) =$ 

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 $\lambda_1(\alpha D(G) + (1 - \alpha)A(G))$  is called the  $A_\alpha$ -spectral radius of *G*, formerly introduced by Nikiforov in [26]. We denote by  $\lambda_n(G)$  the **least eigenvalue** of *G*. The matrix L(G) = D(G) - A(G) is known as the Laplacian matrix of *G*. Let  $\mu_1(G) \ge \mu_2(G) \ge \cdots \ge \mu_{n-1}(G) \ge \mu_n(G)$  be the Laplacian eigenvalues of *G*. It is known that  $\mu_n(G) = 0$ . The values  $\mu_1(G)$  and  $\mu_{n-1}(G)$  are called the Laplacian spectral radius of *G* and the algebraic connectivity of *G*, respectively.

The investigation on sufficient spectral conditions which warrant Hamiltonian *s*-properties of a graph was initiated by Fielder and Nikiforov [11]. However, the results in the literature mainly focus on the spectral radius and the Q-index of dense graphs. Hardly any of them involve graphs with uniform edge density and the Laplacian eigenvalues. Recently, Li [17] initially proved sufficient conditions of  $\rho_1(G)$  based on the connectivity to assure a connected graph to be Hamiltonian and traceable. Inspired by a well-known theorem of Chvatál and Erdös [5], we present tight sufficient spectral conditions on certain matrices arisen from graphs by taking a unified approach to assure a *k*-connected graph to possess Hamiltonian *s*-properties. For an integer *s* with  $s \in \{1, 0, -1\}$ , we say that a graph *G* possesses **Hamiltonian** *s*-properties if each of the following holds: if s = 1, then *G* is Hamilton-connected; if s = 0, then *G* is Hamiltonian; and if s = -1, then *G* is traceable.

Our ideas are also motivated by the literatures [6,8,13,19,18,20–24]. One of our goals is to investigate the relationship between Hamiltonian *s*-properties, the spectral radius  $\lambda_1(aD(G) + bA(G))$  and  $\lambda_1(aD(\overline{G}) + bA(\overline{G}))$  of a *k*-connected graph *G*. This provides a mechanism to take a unified approach to the adjacency spectral radius, the signless Laplacian spectral radius, and the  $A_{\alpha}$ -spectral radius of *G*. Another goal of this research is to initiate studies to find tight bounds of  $\mu_1(G)$ ,  $\mu_{n-1}(G)$  and  $\mu_1(G) + \lambda_n(G)$  to predict *k*-connected graphs to possess Hamiltonian *s*-properties. The main results are as follows.

For real number *a*, integers *k* and  $\delta$  with  $a \ge 0$ ,  $1 \le k \le \delta$ , and  $s \in \{1, 0, -1\}$ , define

$$f(a, n, k, \delta, s) = \begin{cases} \delta \sqrt{\frac{k-s+1}{n-k+s-1}} & \text{if } a = 0, \\ \frac{\delta n}{n-k+s-1} & \text{if } a = 1, \\ \frac{a\delta n}{n-k+s-1} & \text{if } a \in (0, 1), \\ \max\{a\delta, \frac{a\delta(k-s+1)}{n-k+s-1}\} & \text{if } a \in (1, +\infty). \end{cases}$$
(1)

**Theorem 1.1.** Let *G* be a *k*-connected graph of order  $n \ge 10$  and minimum degree  $\delta = \delta(G)$ . Suppose that

$$\lambda_1(aD(G) + A(G)) \le f(a, n, k, \delta, s).$$
<sup>(2)</sup>

Then each of the following holds.

(i) If  $a \in \{0, 1\}$ , then G possesses Hamiltonian s-properties if and only if  $G \cong K_{k,k-s+1}$ . (ii) If 0 < a < 1 or if  $1 < a < +\infty$ , then G possesses Hamiltonian s-properties.

It can be seen that when a = 0 or 1, the upper bound on  $\lambda_1(aD(G) + A(G))$  is tight in some sense in Theorem 1.1. For graphs *G* and *H*, we use  $H \subseteq G$  to denote the fact that *H* is a subgraph of *G*. Let  $\mathcal{F}(p,q) = \{G : K_{p,q} \subseteq G \subseteq K_p \lor qK_1\}$  be a family of graphs.

**Theorem 1.2.** *Let G be a k-connected graph of order*  $n \ge 10$ *. If* 

$$\lambda_1(aD(\overline{G}) + A(\overline{G})) \le (a+1)(k-s),$$

then *G* possesses Hamiltonian *s*-properties if and only if  $G \notin \mathcal{F}(k, k - s + 1)$ .

For the Laplacian matrix, tight bounds on  $\mu_1(G)$  and  $\mu_{n-1}(G)$  to assume a *k*-connected graph to possess Hamiltonian *s*-properties are proved as follows.

**Theorem 1.3.** Let G be a k-connected graph of order  $n \ge 3$  and minimum degree  $\delta = \delta(G)$ . Each of the following holds. (i) If  $\mu_1(G) < \frac{n\delta}{n-k+s-1}$ , then G possesses Hamiltonian s-properties. (ii) If  $\delta = n-k+s-1$  and  $\mu_1(G) \le \frac{n\delta}{n-k+s-1}$ , then G possesses Hamiltonian s-properties if and only if  $G \notin \mathcal{F}(k, k-s+1)$ .

**Theorem 1.4.** Let G be a k-connected graph of order n > 3. If

$$\mu_{n-1}(G) > n-k+s-1$$
,

then *G* possesses Hamiltonian s-properties if and only if  $G \notin \mathcal{F}(k, k - s + 1)$ .

**Theorem 1.5.** Let *G* be a *k*-connected graph of order  $n \ge 3$  and minimum degree  $\delta = \delta(G)$ . If

$$\mu_1(G) + \lambda_n(G) \leq \frac{n\delta}{n-k+s-1} - \sqrt{\left\lfloor \frac{n}{2} \right\rfloor \left\lceil \frac{n}{2} \right\rceil},$$

then G possesses Hamiltonian s-properties if and only if  $G \ncong K_{k,k-s+1}$  for  $s \in \{1, 0\}$ .

It can be observed that the upper bound on  $\mu_1(G) + \lambda_n(G)$  in Theorem 1.5 is tight for  $s \in \{1, 0\}$  in some sense. In the next section, we display some tools to be employed in our arguments. The proofs of the main results are in the subsequent section.

#### 2. Preliminaries

We in this section will present some important results that will be used in our arguments. Recall that a bipartite graph is called balanced if its two partite sets *A* and *B* have equal number of vertices.

**Theorem 2.1.** (Moon and Moser [25]) Let *G* be a balanced bipartite graph of order 2*n* with bipartition (*A*, *B*). If d(x) + d(y) > n for every pair of nonadjacent vertices  $x \in A$  and  $y \in B$ , then *G* is Hamiltonian.

**Theorem 2.2.** (Jackson [16]) Let *G* be a 2-connected bipartite graph with bipartition (*A*, *B*), where  $|A| \ge |B|$ . If each vertex in *A* has degree at least *k* and each vertex in *B* has degree at least *l*, then *G* contains a cycle of length at least  $2 \min(|B|, k + l - 1, 2k - 2)$ . Moreover, if |A| = |B| and k = l, then *G* contains a cycle of length at least  $2 \min(|B|, 2k - 1)$ .

**Theorem 2.3.** (Dirac [9], Ore [27]) Let *G* be a graph of order  $n \ge 3$  and minimum degree  $\delta(G)$ . If

$$\delta(G) \geq \frac{n+s}{2},$$

then G possesses Hamiltonian s-properties.

**Theorem 2.4.** (Chvatál and Erdös [5]) Let G be a k-connected graph of order  $n \ge 3$ . If

$$\alpha(G)\leq k-s,$$

then G possesses Hamiltonian s-properties.

Note that  $k \ge 2$  is a trivial condition in Theorem 2.4 for  $s \in \{1, 0\}$ .

**Theorem 2.5.** (Anderson and Morely [1]) Let G be a graph of order  $n \ge 2$ . Then  $\mu_1(G) \le n$  with equality if and only if  $\overline{G}$  is disconnected.

**Theorem 2.6.** (Fiedler [10]) Let G with n vertices contain an independent set of size t. Then  $\mu_{n-1}(G) \le n-t$ .

**Theorem 2.7.** (Godsil and Newman [12]) Let G be a loopless graph, and  $\mu_1(G)$  be the Laplacian spectral radius. For any independent set I of size t, we have  $t \le n \frac{\mu_1 - \overline{d}_I}{\mu_1}$ , where  $\overline{d}_I = \frac{1}{t} \sum_{i \in I} d_i$ .

**Theorem 2.8.** (Constantine [7]) If G is a graph of order n, then

$$\lambda_n(G) \geq -\sqrt{\left\lfloor \frac{n}{2} \right\rfloor \left\lceil \frac{n}{2} \right\rceil},$$

with equality if and only if  $G \cong K_{\lfloor \frac{n}{2} \rfloor, \lceil \frac{n}{2} \rceil}$ .

Let  $A = (a_{ij})$  and  $B = (b_{ij})$  be two  $n \times n$  matrices. Define  $A \leq B$  if  $a_{ij} \leq b_{ij}$  for all i and j, and A < B if  $A \leq B$  and  $A \neq B$ .

**Theorem 2.9.** (Berman and Plemmons [2], Horn and Johnson [15]) Let  $A = (a_{ij})$  and  $B = (b_{ij})$  be two  $n \times n$  matrices with the spectral radii  $\lambda_1(A)$  and  $\lambda_1(B)$ . If  $0 \le A \le B$ , then  $\lambda_1(A) \le \lambda_1(B)$ . Furthermore, if B is irreducible and  $0 \le A < B$ , then  $\lambda_1(A) < \lambda_1(B)$ .

The main tool in our paper is the eigenvalue interlacing technique described below. Given two non-increasing real sequences  $\theta_1 \ge \theta_2 \ge \cdots \ge \theta_n$  and  $\eta_1 \ge \eta_2 \ge \cdots \ge \eta_m$  with n > m, the second sequence is said to **interlace** the first one if  $\theta_i \ge \eta_i \ge \theta_{n-m+i}$  for  $i = 1, 2, \ldots, m$ . The interlacing is **tight** if exists an integer  $k \in [0, m]$  such that  $\theta_i = \eta_i$  for  $1 \le i \le k$  and  $\theta_{n-m+i} = \eta_i$  for  $k + 1 \le i \le m$ .

Consider an  $n \times n$  real symmetric matrix

$$M = \begin{pmatrix} M_{1,1} & M_{1,2} & \cdots & M_{1,m} \\ M_{2,1} & M_{2,2} & \cdots & M_{2,m} \\ \vdots & \vdots & \ddots & \vdots \\ M_{m,1} & M_{m,2} & \cdots & M_{m,m} \end{pmatrix},$$

whose rows and columns are partitioned according to a partitioning  $X_1, X_2, ..., X_m$  of  $\{1, 2, ..., n\}$ . The **quotient matrix** R(M) of the matrix M is the  $m \times m$  matrix whose entries are the average row sums of the blocks  $M_{i,j}$  of M. The partition is **equitable** if each block  $M_{i,j}$  of M has constant row (and column) sum.

**Theorem 2.10.** (Brouwer and Haemers [4,14]) Let M be a real symmetric matrix. Then the eigenvalues of every quotient matrix of M interlace the ones of M. Furthermore, if the interlacing is tight, then the partition is equitable.

#### 3. Proofs

In the proofs of Theorems 1.1-1.5, we say that a graph possesses Hamiltonian *s*-properties, which means that the graph is Hamilton-connected if s = 1, Hamiltonian if s = 0, and traceable if s = -1. Before proceeding further, we present a technical lemma for the spectral radius of nonnegative matrices of bipartite graphs.

**Lemma 3.1.** Let *H* be a bipartite graph with bipartition (X, Y). If |X| = x, |Y| = y and |E(H)| = r, then

$$\lambda_1(aD(H) + A(H)) \ge \frac{1}{2} \left( a(\frac{r}{x} + \frac{r}{y}) + \sqrt{(a^2 - 1)(\frac{r}{x} - \frac{r}{y})^2 + (\frac{r}{x} + \frac{r}{y})^2} \right).$$

**Proof.** Let R(aD(H) + A(H)) be the quotient matrix of aD(H) + A(H) with respect to the partition (X, Y). One can see that

$$R(aD(H) + A(H)) = \begin{pmatrix} \frac{ar}{x} & \frac{r}{y} \\ \frac{r}{y} & \frac{ar}{y} \end{pmatrix}.$$

A direct computation shows that the characteristic polynomial of R(aD(H) + A(H)) is

$$\lambda^2 - a(\frac{r}{x} + \frac{r}{y})\lambda + (a^2 - 1)\frac{r^2}{xy} = 0,$$

which yields

$$\lambda_1(R(aD(H) + A(H))) = \frac{1}{2} \left( a(\frac{r}{x} + \frac{r}{y}) + \sqrt{a^2(\frac{r}{x} + \frac{r}{y})^2 - 4(a^2 - 1)\frac{r^2}{xy}} \right)$$
$$= \frac{1}{2} \left( a(\frac{r}{x} + \frac{r}{y}) + \sqrt{(a^2 - 1)(\frac{r}{x} - \frac{r}{y})^2 + (\frac{r}{x} + \frac{r}{y})^2} \right).$$

The result follows from Theorem 2.10.  $\Box$ 

We are now in a position to present the proofs of Theorems 1.1-1.5.

**Proof of Theorem 1.1.** (i) It is routine to verify that  $G \cong K_{k,k-s+1}$  does not possess Hamiltonian *s*-properties. Therefore, it suffices to prove the sufficiency. We argue by contradiction and assume that

 $G \ncong K_{k,k-s+1}$  and G does not possess Hamiltonian s-properties.

We shall justify two claims below.

**Claim 1.**  $n \ge 2k - s + 1$ .

In fact, if  $n \le 2k - s$ , then  $\delta(G) \ge \kappa(G) \ge k \ge \frac{n+s}{2}$ . By Theorem 2.3, *G* possesses Hamiltonian *s*-properties, a contradiction. Claim 1 holds.

By Theorem 2.4,  $\alpha(G) \ge k - s + 1$ , and then there exists an independent set  $X = \{u_i \in V(G) | 1 \le i \le k - s + 1\}$ . Let  $Y = V(G) \setminus X = \{v_j | 1 \le j \le n - k + s - 1\}$ . Consider the bipartite spanning subgraph *H* of *G* with the partitions *X* and *Y*. Let *r* be the number of edges with one end-vertex in *X* and the other in *Y*. Then  $r \ge \delta(k - s + 1)$ . For simplicity, define  $\xi = k - s + 1$ . By Theorem 2.9 and Lemma 3.1, we have

$$\lambda_1(aD(G) + A(G)) \ge \lambda_1(aD(H) + A(H))$$

$$\ge \frac{1}{2} \left( a(\frac{r}{\xi} + \frac{r}{n-\xi}) + \sqrt{(a^2 - 1)(\frac{r}{\xi} - \frac{r}{n-\xi})^2 + (\frac{r}{\xi} + \frac{r}{n-\xi})^2} \right).$$
(4)

**Claim 2.** If  $a \in \{0, 1\}$ , then  $2k - s + 1 \le n \le 2(k - s + 1)$ .

(3)

Assume first that a = 0. By (4) and the assumption of Theorem 1.1, we have

$$\lambda_{1}(aD(G) + A(G)) \geq \lambda_{1}(aD(H) + A(H)) \geq \lambda_{1}(R(aD(H) + A(H)))$$
$$\geq r\sqrt{\frac{1}{\xi(n-\xi)}} \geq \delta\sqrt{\frac{\xi}{n-\xi}} = \delta\sqrt{\frac{k-s+1}{n-k+s-1}}$$
$$\geq \lambda_{1}(aD(G) + A(G)).$$
(5)

It follows that all the inequalities in (5) must be equalities. Hence  $G \cong H$  and  $r = \delta(k - s + 1)$ . Furthermore,  $\lambda_n(aD(H) + A(H)) = -\lambda_1(aD(H) + A(H)) = -r\sqrt{\frac{1}{\xi(n-\xi)}} = \lambda_2(R(aD(H) + A(H)))$ , and thus the interlacing is tight. By Theorem 2.10, the partition is equitable. That is to say, each vertex  $v_j$  of Y in G has the same number of neighbors in X, and thus  $\frac{r}{n-k+s-1} \ge \delta(G)$ . Then we have  $n \le 2(k-s+1)$ . By Claim 1, we have  $2k - s + 1 \le n \le 2(k-s+1)$ .

Next, we assume that a = 1. By (4) and assumption of Theorem 1.1, we have

$$\lambda_1(aD(G) + A(G)) \ge \lambda_1(aD(H) + A(H)) \ge \lambda_1(R(aD(H) + A(H)))$$
  
$$\ge \frac{r}{\xi} + \frac{r}{n-\xi} \ge \frac{\delta n}{n-\xi} = \frac{\delta n}{n-k+s-1}$$
  
$$\ge \lambda_1(aD(G) + A(G)).$$
(6)

It follows that all the inequalities in (6) must be equalities. Hence  $G \cong H$  and  $r = \delta(k - s + 1)$ . Furthermore,  $\lambda_n(aD(H) + A(H)) = 0 = \lambda_2(R(aD(H) + A(H)))$ , and hence the interlacing is tight. By Theorem 2.10, the partition is equitable. That is, each vertex  $v_j$  of Y in G has the same number of neighbors in X, and thus  $\frac{r}{n-k+s-1} \ge \delta(G)$ . Then we have  $n \le 2(k-s+1)$ . By Claim 1, we have  $2k - s + 1 \le n \le 2(k - s + 1)$ . This proves Claim 2.

By the assumption of Theorem 1.1,  $s \in \{1, 0, -1\}$ . If s = 1, then n = 2k, and so both n - (k - s + 1) = k and  $d(v_j) = \delta(G)$ . As  $d(u_i) = d(v_j) = \delta(G) \ge \kappa(G) \ge k$ , we observe that  $G \cong K_{k,k} = K_{k,k-s+1}$ , contrary to (3).

Assume that s = 0. Then by Claim 2,  $n \in \{2k + 1, 2k + 2\}$ . If n = 2k + 1, then n - (k - s + 1) = k and  $d(v_j) > \delta(G) \ge k$ . As  $d(u_i) = \delta(G) \ge \kappa(G) \ge k$ , we have  $G \cong K_{k,k+1} = K_{k,k-s+1}$ , contrary to (3). If n = 2k + 2, then n - (k - s + 1) = k + 1 and  $d(v_j) = \delta(G)$ . As  $d(u_i) = d(v_j) = \delta(G) \ge \kappa(G) \ge k$ , it follows that G is a balanced bipartite graph of order 2k + 2 such that, as  $k \ge 2$ ,  $d(u_i) + d(v_j) \ge 2k > k + 1$  for any  $u_i v_j \notin E(G)$ . By Theorem 2.1, G is Hamiltonian, which is contrary to (3).

Finally we assume that s = -1. Then by Claim 2,  $n \in \{2k + 2, 2k + 3, 2k + 4\}$ . If n = 2k + 2, then n - (k - s + 1) = k and  $d(v_j) > \delta(G)$ , and in this case,  $d(u_i) = \delta(G) \ge \kappa(G) \ge k$ . Then  $G \cong K_{k,k+2} = K_{k,k-s+1}$ , contrary to (3). If n = 2k + 3, then n - (k - s + 1) = k + 1 and  $d(v_j) > \delta(G) \ge k$  with  $d(u_i) = \delta(G) \ge \kappa(G) \ge k$ . Since  $n \ge 9$ , we have  $k \ge 3$ , and so  $k + 1 \le 2k - 2$ . By Theorem 2.2, *G* contains a cycle of length 2k + 2. Since n = 2k + 3 and  $k \ge 3$ , it follows that *G* has a path containing all the vertices of *G*, and therefore *G* is traceable, contrary to (3). Assume that n = 2k + 4. Then n - (k - s + 1) = k + 2 = k - s + 1 and  $d(v_j) = \delta(G)$ . By  $n \ge 10$ , we have  $k \ge 3$ . As  $d(u_i) = d(v_j) = \delta(G) \ge \kappa(G) \ge k$ , it follows by Theorem 2.2 that *G* contains a cycle of length 2k + 4 which implies that *G* is traceable, contrary to (3). This completes the proof of Theorem 1.1(i).

(ii) We argue by contradiction and assume that G does not possess Hamiltonian s-properties. If 0 < a < 1, then  $a^2 - 1 < 0$ , and so by (4),

$$\begin{split} \lambda_1(aD(H) + A(H)) &\geq \frac{1}{2} \left( a(\frac{r}{\xi} + \frac{r}{n-\xi}) + \sqrt{(a^2 - 1)(\frac{r}{\xi} - \frac{r}{n-\xi})^2 + (\frac{r}{\xi} + \frac{r}{n-\xi})^2} \right) \\ &> \frac{1}{2} \left( a(\frac{r}{\xi} + \frac{r}{n-\xi}) + \sqrt{(a^2 - 1)(\frac{r}{\xi} + \frac{r}{n-\xi})^2 + (\frac{r}{\xi} + \frac{r}{n-\xi})^2} \right) \\ &= a(\frac{r}{\xi} + \frac{r}{n-\xi}). \end{split}$$

It follows that

$$\begin{split} \lambda_1(aD(G)+A(G)) &\geq \lambda_1(aD(H)+A(H)) \\ &> a(\frac{r}{\xi}+\frac{r}{n-\xi}) \geq \frac{a\delta n}{n-\xi} = \frac{a\delta n}{n-k+s-1}, \end{split}$$

contrary to (2).

If  $1 < a < +\infty$ , then  $a^2 - 1 > 0$ , and so by (4),

$$\begin{split} \lambda_1(aD(H) + A(H)) &\geq \frac{1}{2} \left( a(\frac{r}{\xi} + \frac{r}{n-\xi}) + \sqrt{(a^2 - 1)(\frac{r}{\xi} - \frac{r}{n-\xi})^2 + (\frac{r}{\xi} + \frac{r}{n-\xi})^2} \right) \\ &> \frac{1}{2} \left( a(\frac{r}{\xi} + \frac{r}{n-\xi}) + \sqrt{(a^2 - 1)(\frac{r}{\xi} - \frac{r}{n-\xi})^2 + (\frac{r}{\xi} - \frac{r}{n-\xi})^2} \right) \\ &= \max\{\frac{ar}{\xi}, \frac{ar}{n-\xi}\}. \end{split}$$

It follows that

$$\lambda_1(aD(G) + A(G)) \ge \lambda_1(aD(H) + A(H))$$
  
> max{ $\frac{ar}{\xi}, \frac{ar}{n-\xi}$ }  $\ge$  max{ $a\delta, \frac{a\delta(k-s+1)}{n-k+s-1}$ },

contrary to (2). We complete the proof of Theorem 1.1(ii).  $\Box$ 

By definition, any graph  $G \in \mathcal{F}(k, k-s+1)$  is not Hamilton-connected if s = 1, not Hamiltonian if s = 0, and not traceable if s = -1. Thus we have the following observation.

If 
$$G \in \mathcal{F}(k, k - s + 1)$$
, then G does not possess Hamiltonian s-properties. (7)

Proof of Theorem 1.2. By (7), it suffices to prove the sufficiency. We assume that

 $G \notin \mathcal{F}(k, k-s+1)$  and G does not possess Hamiltonian s-properties. (8)

By Theorem 2.4,  $\alpha(G) \ge k - s + 1$ , and thus there exists an independent set  $X = \{u_i \in V(G) | 1 \le i \le k - s + 1\}$  in *G*. Let  $Y = V(G) \setminus X = \{v_j | 1 \le j \le n - k + s - 1\}$ .

If G[Y] is a clique in G, then  $K_{k-s+1} \cup (n-k+s-1)K_1$  is a spanning subgraph of  $\overline{G}$ . It follows by the hypothesis of Theorem 1.2 and Theorem 2.9 that

$$(a+1)(k-s) \ge \lambda_1 (aD(\overline{G}) + A(\overline{G})) \ge \lambda_1 (aD(K_{k-s+1} \cup (n-k+s-1)K_1) + A(K_{k-s+1} \cup (n-k+s-1)K_1)) = (a+1)(k-s).$$
(9)

Thus all the inequalities in (9) must be equalities. Hence  $\overline{G} \cong K_{k-s+1} \cup (n-k+s-1)K_1$ , and so  $G \cong K_{n-k+s-1} \vee (k-s+1)K_1$ . Since *G* does not possess Hamiltonian *s*-properties, by Theorem 2.3, we must have  $n-k+s-1 = \delta(G) < \frac{n+s}{2}$ , and so  $2(n-k+s-1) \leq n+s-1$ . This implies  $n-k+s-1 \leq k$ . Note that *G* is *k*-connected. Then  $n-k+s-1 = \delta(G) \geq k$ . Thus  $G \cong K_k \vee (k-s+1)K_1$ , which contradicts (8).

If G[Y] is not a clique in G, then  $K_{k-s+1} \cup \overline{G[Y]}$  is a spanning subgraph of  $\overline{G}$ . By the assumption of Theorem 1.2 and Theorem 2.9, we have

$$(a+1)(k-s) \ge \lambda_1 (aD(\overline{G}) + A(\overline{G})) \ge \lambda_1 (aD(K_{k-s+1} \cup \overline{G[Y]}) + A(K_{k-s+1} \cup \overline{G[Y]})).$$
(10)

Note that (10) holds for any G[Y]. Then we have  $|Y| = n - k + s - 1 \le k - s + 1$ , and so  $n \le 2(k - s + 1)$ . By Claim 1 in the proof of Theorem 1.1, we obtain that  $n \ge 2k - s + 1$ . Therefore,  $2k - s + 1 \le n \le 2(k - s + 1)$ .

Assume that s = 1. Then n = 2k, that is, |Y| = |X| = k. As  $d(u_i) \ge \delta(G) \ge k$  for any  $u_i \in X$ , we obtain that  $G \in \mathcal{F}(k, k) = \mathcal{F}(k, k - s + 1)$ , which is contrary to (8).

Next we assume that s = 0. Then  $n \in \{2k + 1, 2k + 2\}$ . If n = 2k + 1, then |Y| = k and |X| = k + 1. Since  $d(u_i) \ge \delta(G) \ge k$ , then we have  $G \in \mathcal{F}(k, k + 1) = \mathcal{F}(k, k - s + 1)$ , contrary to (8). If n = 2k + 2, then |Y| = |X| = k + 1. For any G[Y], we assume that G is not Hamiltonian in (8). However, when G[Y] is an independent set, we observe that G is a balanced bipartite graph of order 2k + 2 such that, as  $k \ge 2$ ,  $d(u_i) + d(v_j) \ge 2k > k + 1$  for any  $u_i v_j \notin E(G)$ . By Theorem 2.1, G is Hamiltonian, a contradiction.

Finally we suppose that s = -1. Then  $n \in \{2k + 2, 2k + 3, 2k + 4\}$ . If n = 2k + 2, then |Y| = k and |X| = k + 2. Note that  $d(u_i) \ge \delta(G) \ge k$ . Then  $G \in \mathcal{F}(k, k + 2) = \mathcal{F}(k, k - s + 1)$ , contrary to (8). If n = 2k + 3, then |Y| = k + 1 and |X| = k + 2. As  $n \ge 9$ , then we have  $k \ge 3$ , and therefore  $k + 1 \le 2k - 2$ . For any G[Y], we always assume that G is not traceable in (8). However, when G[Y] is an independent set, we observe that G is a 3-connected bipartite graph with bipartition (X, Y) and |X| > |Y|. Note that  $d(u_i) \ge \delta(G) \ge k$  and  $d(v_j) \ge \delta(G) \ge k$ . By Theorem 2.2, G contains a cycle of length 2k + 2. Note that n = 2k + 3 and  $k \ge 3$ , it follows that G has a path containing all the vertices of G, and therefore G is traceable, a contradiction. If n = 2k + 4, then |Y| = |X| = k + 2. By  $n \ge 10$ , we have  $k \ge 3$  and  $k + 2 \le 2k - 1$ . For any G[Y], we always assume that G is not traceable in (8). However, when G[Y] is an independent set, G is a 3-connected bipartite graph with bipartition (X, Y) and |X| = |Y|. Note that  $d(u_i) \ge k(G) \ge k$  and  $d(u_j) \ge \delta(G) \ge k$ . It follows by Theorem 2.2 that G contains a cycle of length 2k + 4, then G is traceable, a contradiction. This completes the proof of Theorem 1.2.  $\Box$ 

Before proving Theorem 1.3, we shall indicate that it suffices to consider graphs satisfying the inequality  $\delta \le n - k + s - 1$ when discussing Theorem 1.3. In fact, assume that *G* is a graph with  $\delta > n - k + s - 1$ . Then as  $\delta \ge \kappa(G) \ge k$ , we have  $2\delta \ge n + s$ , and so  $\delta \ge \frac{n+s}{2}$ . By Theorem 2.3, *G* possesses Hamiltonian *s*-properties. Thus we only need to consider  $\delta \le n - k + s - 1$ , or equivalently,  $\frac{n\delta}{n-k+s-1} \le n$ . By Theorem 2.5, the upper bound on  $\mu_1(G)$  in Theorem 1.3 is well-defined. **Proof of Theorem 1.3.** Suppose, to the contrary, that *G* does not possess Hamiltonian *s*-properties. By Theorems 2.3 and 2.4, then  $n \ge 2\delta - s + 1$  and  $\alpha(G) \ge k - s + 1$ . Let *X* be an independent set in *G* such that |X| = k - s + 1. Let *r* be the number of edges between *X* and  $V(G) \setminus X$ . Then  $r \ge \delta |X|$ . Accordingly, the quotient matrix R(L) of L(G) on the partition  $(X, V(G) \setminus X)$  becomes:

$$R(L) = \left(\begin{array}{cc} \frac{r}{k-s+1} & -\frac{r}{k-s+1} \\ -\frac{r}{n-k+s-1} & \frac{r}{n-k+s-1} \end{array}\right).$$

Let  $\mu_1(R(L)), \mu_2(R(L))$  be the eigenvalues of R(L). Then  $\mu_1(R(L)) \ge \mu_2(R(L)) = 0$ . By algebraic manipulations, we have

$$\mu_1(R(L)) = \frac{r}{k-s+1} + \frac{r}{n-k+s-1}.$$

By Theorem 2.10, then

$$\mu_1(G) \ge \mu_1(R(L)) = \frac{r}{k-s+1} + \frac{r}{n-k+s-1}$$

$$\ge (\frac{1}{k-s+1} + \frac{1}{n-k+s-1})(k-s+1)\delta = \frac{n\delta}{n-k+s-1},$$
(11)

which contradicts the assumption of this theorem.

Furthermore, if  $\mu_1(G) \le \frac{n\delta}{n-k+s-1}$ , then all the inequalities in (11) must be equalities. So  $r = (k-s+1)\delta$ , and the partition is equitable. That is to say, each vertex of *X* has  $\delta$  neighbors in *V*(*G*)\*X*, and each vertex of *V*(*G*)\*X* has  $\frac{(k-s+1)\delta}{n-k+s-1}$  neighbors in *X*. Note that  $\delta = n - k + s - 1$ . As *G* is *k*-connected, we have  $\delta \ge k$ . Thus  $2\delta - s + 1 \ge \delta + k - s + 1 = n \ge 2\delta - s + 1$ , forcing  $\delta = k$  and n = 2k - s + 1. Hence  $G \in \mathcal{F}(k, k - s + 1)$ , a contradiction. Conversely, it is obvious that  $G \in \mathcal{F}(k, k - s + 1)$  does not possess Hamiltonian *s*-properties.  $\Box$ 

**Proof of Theorem 1.4.** By (7), it suffices to prove the sufficiency. Let  $G \notin \mathcal{F}(k, k - s + 1)$  be a graph. We assume that G does not possess Hamiltonian *s*-properties. As  $n \ge 3$ , G cannot be a complete graph. By Theorems 2.3 and 2.4, then  $n \ge 2\delta - s + 1$  and  $\alpha(G) \ge k - s + 1$ . Let I be an independent set with size  $\alpha(G)$  in G. By Theorem 2.6 and the hypothesis of Theorem 1.4, we have

$$n - k + s - 1 \le \mu_{n-1}(G) \le n - \alpha(G) \le n - k + s - 1.$$

Hence

$$n-k+s-1=n-\alpha(G)=\mu_{n-1}(G)\leq\kappa(G)\leq\delta$$

and so

$$\delta + k - s + 1 \ge n \ge 2\delta - s + 1 \ge \delta + \kappa(G) - s + 1 \ge \delta + k - s + 1.$$

Therefore  $n = 2\delta - s + 1$ ,  $\delta = \kappa(G) = k$  and  $\alpha(G) = k - s + 1$ , implying that  $|V(G)\setminus I| = k = \delta$ . Note that for each vertex  $v \in I$ ,  $d(v) \ge \delta$ . Then  $G \in \mathcal{F}(k, k - s + 1)$ , a contradiction.  $\Box$ 

**Proof of Theorem 1.5.** It is routine to verify that  $G \cong K_{k,k-s+1}$  does not possess Hamiltonian *s*-properties, where  $s \in \{1, 0\}$  be an integer. Therefore, it suffices to prove the sufficiency. We assume that

$$G \ncong K_{k,k-s+1}, \text{ where } s \in \{1,0\}, \tag{12}$$

and

*G* does not possess Hamiltonian *s*-properties, where 
$$s \in \{1, 0, -1\}$$
. (13)

By Theorems 2.3 and 2.4, then  $n \ge 2\delta - s + 1 \ge 2k - s + 1$ , and  $\alpha(G) \ge k - s + 1$ . Let *I* be an independent set with size  $\alpha(G)$  in *G*. By Theorem 2.7,  $\alpha(G) \le n \frac{\mu_1 - \overline{d_I}}{\mu_1}$ , and hence  $\mu_1 \ge \frac{n\overline{d_I}}{n - \alpha(G)}$ . Note that  $\overline{d_I} \ge \delta$  and  $\alpha(G) \ge k - s + 1$ . Then  $\mu_1 \ge \frac{n\overline{d_I}}{n - \alpha(G)} \ge \frac{n\delta}{n - k + s - 1}$ . Combining Theorem 2.8 and the condition of Theorem 1.5, we have

$$\frac{n\delta}{n-k+s-1} - \sqrt{\left\lfloor \frac{n}{2} \right\rfloor \left\lceil \frac{n}{2} \right\rceil} \le \frac{n\overline{d}_I}{n-\alpha(G)} - \sqrt{\left\lfloor \frac{n}{2} \right\rfloor \left\lceil \frac{n}{2} \right\rceil} \le \mu_1 + \lambda_n \le \frac{n\delta}{n-k+s-1} - \sqrt{\left\lfloor \frac{n}{2} \right\rfloor \left\lceil \frac{n}{2} \right\rceil},$$

and therefore all the inequalities must be equalities, that is,

$$\mu_1 = \frac{n\overline{d}_I}{n - \alpha(G)} = \frac{n\delta}{n - k + s - 1}, \ \lambda_n = -\sqrt{\left\lfloor \frac{n}{2} \right\rfloor \left\lceil \frac{n}{2} \right\rceil}$$

So  $\alpha(G) = k - s + 1$ ,  $G \cong K_{\frac{n}{2}, \frac{n}{2}}$  if *n* is even, and  $G \cong K_{\frac{n-1}{2}, \frac{n+1}{2}}$  if *n* is odd. Next we consider three different values of *s*, respectively.

• s = 1. Note that *G* is not Hamiltonian-connected by (13). Then  $G \cong K_{\frac{n}{2},\frac{n}{2}}$  if *n* is even, and  $G \cong K_{\frac{n-1}{2},\frac{n+1}{2}}$  if *n* is odd. Consider even *n*. Note that  $\alpha(G) = k$ . Then  $\frac{n}{2} = k$ , and hence  $G \cong K_{k,k}$ , contrary to (12). For odd *n*. Note that  $\alpha(G) = k$ . Then  $\frac{n+1}{2} = k$ , and hence n = 2k - 1, which contradicts  $n \ge 2k$ .

• s = 0. Note that *G* is not Hamiltonian by (13). Then  $G \cong K_{\frac{n-1}{2},\frac{n+1}{2}}$ . Note that  $\alpha(G) = k+1$ . Then  $\frac{n+1}{2} = k+1$ , and hence n = 2k + 1. So  $G \cong K_{k,k+1}$ , which is contrary to (12).

• s = -1. It is obvious that  $K_{\frac{n}{2},\frac{n}{2}}$  and  $G \cong K_{\frac{n-1}{2},\frac{n+1}{2}}$  are traceable, which contradicts (13).

#### 4. Corollaries of Theorems 1.1 and 1.2

Throughout this section, we assume that *a* and *b* are real numbers with  $a \ge 0$  and b > 0, *k*, *s* and  $\delta$  are integers with  $1 \le k \le \delta$  and  $s \in \{1, 0, -1\}$ . Next we consider the nonnegative matrix aD(G) + bA(G). Theorems 1.1 and 1.2 have the following more general forms.

**Corollary 4.1.** *Let G be a k*-connected graph of order  $n \ge 10$  and minimum degree  $\delta(G)$ *. If* 

$$\lambda_1(aD(G) + bA(G)) \le bf(\frac{a}{b}, n, k, \delta, s),$$

then each of the following holds.

(i) a = 0 or a = b. G possesses Hamiltonian s-properties if and only if  $G \ncong K_{k,k-s+1}$ . (ii) 0 < a < b or a > b. G possesses Hamiltonian s-properties.

**Proof.** Suppose that *G* does not possess Hamiltonian *s*-properties.

(i) a = 0 or a = b. That is,  $\frac{a}{b} = 0$  or 1. Notice that  $K_{k,k-s+1}$  does not possess Hamiltonian *s*-properties. Therefore, it suffices to prove the sufficiency. Assume that  $G \ncong K_{k,k-s+1}$ . By Theorem 1.1(i),  $\lambda_1(aD(G) + A(G)) > f(a, n, k, \delta, s)$ . Note that  $aD(G) + bA(G) = b(\frac{a}{b}D(G) + A(G))$ . It follows that  $\lambda_1(aD(G) + bA(G)) = b\lambda_1(\frac{a}{b}D(G) + A(G)) > bf(\frac{a}{b}, n, k, \delta, s)$ , a contradiction. (ii) 0 < a < b or a > b, i.e.,  $\frac{a}{b} \in (0, 1)$  or  $(1, +\infty)$ . Similar to the last part of proof of (i), (ii) follows immediately.  $\Box$ 

As one of main results of this paper, Corollary 4.1 can be applied to obtain sufficient condition in terms of the  $A_{\alpha}$ -spectral radius  $\lambda_1(A_{\alpha}(G))$  for a *k*-connected graph *G* to possess Hamiltonian *s*-properties.

**Corollary 4.2.** Let G be a k-connected graph of order  $n \ge 10$  and minimum degree  $\delta(G)$ , and let  $\alpha$  be a real number with  $\alpha \in [0, 1)$ . If

$$\lambda_1(A_{\alpha}(G)) \leq (1-\alpha)f(\frac{\alpha}{1-\alpha}, n, k, \delta, s),$$

then each of the following holds. (i)  $\alpha = 0$  or  $\alpha = \frac{1}{2}$ . *G* possesses Hamiltonian s-properties if and only if  $G \ncong K_{k,k-s+1}$ . (ii)  $\alpha \in (0, \frac{1}{2})$  or  $\alpha \in (\frac{1}{2}, 1)$ . *G* possesses Hamiltonian s-properties.

Particularly, sufficient conditions on  $\rho_1(G)$  and  $q_1(G)$  are as follows.

**Corollary 4.3.** *Let G be a k*-connected graph of order  $n \ge 4$  with minimum degree  $\delta(G)$ . If

$$\rho_1(G) \le \delta \sqrt{\frac{k-s+1}{n-k+s-1}},$$

then each of the following holds.

(i) *G* is Hamiltonian-connected if and only if  $G \ncong K_{k,k-s+1}$  for s = 1. (ii) (R. Li [17]) *G* is Hamiltonian if and only if  $G \ncong K_{k,k-s+1}$  for s = 0, where  $n \ge 6$ . (iii) (R. Li [17]) *G* is traceable if and only if  $G \ncong K_{k,k-s+1}$  for s = -1, where  $n \ge 10$ .

**Corollary 4.4.** Let *G* be a *k*-connected graph of order  $n \ge 4$  with minimum degree  $\delta(G)$ . If

$$q_1(G) \leq \frac{\delta n}{n-k+s-1},$$

then each of the following holds. (i) *G* is Hamiltonian-connected if and only if  $G \ncong K_{k,k-s+1}$  for s = 1. (ii) *G* is Hamiltonian if and only if  $G \ncong K_{k,k-s+1}$  for s = 0, where  $n \ge 6$ . (iii) *G* is traceable if and only if  $G \ncong K_{k,k-s+1}$  for s = -1, where  $n \ge 10$ .

At the end of this paper, tight upper bound on  $\lambda_1(aD(\overline{G}) + bA(\overline{G}))$  are proposed.

**Corollary 4.5.** Let *G* be a *k*-connected graph of order  $n \ge 10$  and minimum degree  $\delta(G)$ . If  $\lambda_1(aD(\overline{G}) + bA(\overline{G})) \le (a+b)(k-s)$ , then *G* possesses Hamiltonian *s*-properties if and only if  $G \notin \mathcal{F}(k, k-s+1)$ .

**Proof.** Note that  $\mathcal{F}(k, k - s + 1)$  does not possess Hamiltonian *s*-properties. Therefore, it suffices to prove the sufficiency. For  $G \notin \mathcal{F}(k, k - s + 1)$ . Suppose that *G* does not possess Hamiltonian *s*-properties. By Theorem 1.2,  $\lambda_1(aD(\overline{G}) + A(\overline{G})) > (a + 1)(k - s)$ . Since  $aD(\overline{G}) + bA(\overline{G}) = b(\frac{a}{b}D(\overline{G}) + A(\overline{G}))$ , it follows that  $\lambda_1(aD(\overline{G}) + bA(\overline{G})) = b\lambda_1(\frac{a}{b}D(\overline{G}) + A(\overline{G})) > b(\frac{a}{b} + 1)(k - s) = (a + b)(k - s)$ , a contradiction.  $\Box$ 

By choosing different values of *a* and *b*, sufficient conditions in terms of  $\rho_1(\overline{G})$ ,  $q_1(\overline{G})$  and  $\lambda_1(A_\alpha(\overline{G}))$  for a *k*-connected graph *G* to possess Hamiltonian *s*-properties are easily obtained.

#### **Declaration of competing interest**

The authors declare that they have no known competing financial interests or personal relationships that could have appeared to influence the work reported in this paper.

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