



# Spectral analogues of Erdős' theorem on Hamilton-connected graphs

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## ABSTRACT

A graph  $G$  is Hamilton-connected if for any pair of vertices  $v$  and  $w$ ,  $G$  has a spanning  $(v, w)$ -path. Extending theorems of Dirac and Ore, Erdős prove a sufficient condition in terms of minimum degree and the size of  $G$  to assure  $G$  to be Hamiltonian. We investigate the spectral analogue of Erdős' theorem for a Hamilton-connected graph with given minimum degree, and prove that there exist two graphs  $\{L_n^k, M_n^k\}$  such that each of the following holds for an integer  $k \geq 3$  and a simple graph  $G$  on  $n$  vertices.

(i) If  $n \geq 6k$ ,  $\delta(G) \geq k$ , and  $|E(G)| > \binom{n-k}{2} + k(k+1)$ , then  $G$  is Hamilton-connected if and only if  $C_{n+1}(G) \notin \{L_n^k, M_n^k\}$ .

(ii) If  $n \geq \max\{6k, \frac{1}{2}k^3 - \frac{1}{2}k^2 + k + 4\}$ ,  $\delta(G) \geq k$  and spectral radius  $\lambda(G) \geq n - k$ , then  $G$  is Hamilton-connected if and only if  $G \notin \{L_n^k, M_n^k\}$ .

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## 1. Introduction

We consider finite and simple graphs, with undefined notation and term following [3]. We normally use  $e(G)$ ,  $n$ ,  $\delta(G)$  and  $A(G)$  to denote  $|E(G)|$ ,  $|V(G)|$ , the minimum degree and the adjacency matrix of a graph  $G$ , respectively. The largest eigenvalue of  $A(G)$ , called the *spectral radius* of  $G$ , is denoted by  $\lambda(G)$ . Let  $H$  be a subgraph of a graph  $G$ , and let  $u \in V(G)$ . The set of neighbors of a vertex  $u$  in  $H$  is denoted by  $N_H(u)$ . Thus

$$N_H(u) = \{v \in V(H) : uv \in E(G)\}.$$

Define  $d_H(u) = |N_H(u)|$ . A *clique* is a subset of vertices of an undirected graph whose induced subgraph is a complete graph. The maximum size of a clique of a graph is called *clique number*, denoted by  $\omega(G)$ . For  $S \subseteq V(G)$ , the *induced subgraph*  $G[S]$  is the graph with vertex set  $S$  and edge set  $\{uv \in E(G) \mid u, v \in S\}$ .

The *disjoint union* of two graphs  $G_1$  and  $G_2$ , denoted by  $G_1 + G_2$ , is the graph with the vertex set  $V(G_1) \cup V(G_2)$  and edge set  $E(G_1) \cup E(G_2)$ . The disjoint union of  $k$  copies of a graph  $G$  is denoted by  $kG$ . The *join* of  $G_1$  and  $G_2$ , denoted by  $G_1 \vee G_2$ , has vertex set  $V(G_1) \cup V(G_2)$  and edge set  $E(G_1) \cup E(G_2) \cup \{xy \mid x \in V(G_1), y \in V(G_2)\}$ .

A path (or a cycle, respectively) of a graph  $G$  is called a *Hamilton path* (or *Hamilton cycle*, respectively) if it passes through all the vertices of  $G$ . A graph is *Hamilton-connected* if any two vertices are connected by a Hamilton path. The investigation of hamiltonian graphs has a long history. Dirac and Ore proved the following.

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**Theorem 1.1.** Let  $G$  be a graph of order  $n$ .

- (i) (Dirac [6]) If  $\delta(G) \geq \frac{n}{2}$ , then  $G$  is Hamiltonian.
- (ii) (Ore [14]) If  $e(G) > \binom{n-1}{2} + 1$ , then  $G$  is Hamiltonian.

Motivated by these results, Erdős [7] later extended Theorem 1.1 (ii) by utilizing the minimum degree as a new parameter.

**Theorem 1.2.** (Erdős [7]) Let  $G$  be a graph of order  $n$  and the minimum degree  $\delta$  and  $k$  be an integer with  $1 \leq k \leq \delta \leq \frac{n-1}{2}$ . If

$$e(G) > \max \left\{ \binom{n-k}{2} + k^2, \binom{\lceil \frac{n+1}{2} \rceil}{2} + \left\lfloor \frac{n+1}{2} \right\rfloor^2 \right\},$$

then  $G$  is Hamiltonian.

How many edges can ensure a graph to be Hamilton-connected with a given number of vertices? In 1963, Ore [15] answered the question.

**Theorem 1.3.** [15] Let  $G$  be a graph of order  $n$ , if

$$e(G) \geq \binom{n-1}{2} + 3,$$

then  $G$  is Hamilton-connected.

**Theorem 1.4.** ([16], Theorem 1.8) Let  $G$  be a graph of order  $n \geq 6k^2 - 8k + 5$  with  $\delta(G) \geq k \geq 2$ . If  $e(G) \geq \frac{n^2 - (2k-1)n + 2k - 2}{2}$ , then  $G$  is Hamilton-connected unless  $cl_{n+1}(G) = K_2 \vee (K_{n-k-1} \cup K_{k-1})$  or  $cl_{n+1}(G) = K_k \vee (K_{n-2k-1} \cup \bar{K}_{k-1})$ .

**Theorem 1.5.** ([16], Corollary 1.10) Let  $G$  be a graph of order  $n \geq \max\{6k^2 - 8k + 5, \frac{k^3 - k^2 + 4k - 1}{2}\}$  with  $\delta(G) \geq k \geq 2$ . If  $\rho(G) \geq n - k$ , then  $G$  is Hamilton-connected unless  $G = K_2 \vee (K_{n-k-1} \cup K_{k-1})$  or  $G = K_k \vee (K_{n-2k+1} \cup \bar{K}_{k-1})$ .

The results above, as well as the recent advances in [9,13,16], motivate the current research. In this paper, we present a spectral analogous of Erdős theorem for a Hamilton-connected graph with a given minimum degree. For a graph  $G$ , notice that  $\delta(G) \geq 3$  is a necessary condition for  $G$  to be Hamilton-connected. A sufficient condition for a Hamilton-connected graph in terms of spectral radius is also justified. This paper is independently research work with Chen and Zhang's ([16]) results.

Throughout this paper, for  $2 \leq k \leq \frac{n}{2}$ , let

$$L_n^k = K_2 \vee (K_{n-k-1} + K_{k-1}) \text{ and } M_n^k = K_k \vee (K_{n-2k+1} + (k-1)K_1).$$

In Section 2, extremal sizes of graphs to ensure Hamilton-connectedness are investigated. These will be applied in Section 3 to find an optimal spectral sufficient condition for a graph  $G$  to be Hamilton-connected.

## 2. Extremal sizes of Hamilton-connected graphs

Let  $X, Y$  be vertex subsets of a graph  $G$ . Following [3], we adopt these notation:  $e(X) = |E(G[X])|$ ,

$$E_G[X, Y] = \{xy \in E(G) : x \in X \text{ and } y \in Y\}, \text{ and } e(X, Y) = |E_G[X, Y]|.$$

Throughout this section, if  $J$  is a subgraph of  $G$  and  $v \in V(G) - V(J)$ , define  $d_J(v) = |E_G[\{v\}, V(J)]|$ .

The purpose of this section is to prove two extremal results, namely, Theorems 2.2 and 2.5 in this section, on the optimal sizes to assure a graph to be Hamilton-connected. We state some known results as our tools.

**Theorem 2.1.** (Erdős, Gallai, [8]) Let  $G$  be a graph of order  $n \geq 3$ , and  $u, v$  are any pair distinct and nonadjacent vertices. If

$$d_G(u) + d_G(v) \geq n + 1,$$

then  $G$  is Hamilton-connected.

**Lemma 2.1.** [1] Let  $G$  be a graph of order  $n \geq 3$  with the degree sequence  $(d_1, d_2, \dots, d_n)$ , where  $d_1 \leq d_2 \leq \dots \leq d_n$ . If there is no integer  $2 \leq t \leq \frac{n}{2}$  such that  $d_{t-1} \leq t$  and  $d_{n-t} \leq n - t$ , then  $G$  is Hamilton-connected.

**Theorem 2.2.** Let  $G$  be a graph with order  $n$  and the minimum degree  $\delta$ , and let  $k$  be an integer with  $2 \leq k \leq \delta$ . If

$$e(G) > \max \left\{ \binom{n-k+1}{2} + k(k-1), \binom{\lceil \frac{n}{2} \rceil + 1}{2} + \left\lfloor \frac{n}{2} \right\rfloor \left( \left\lfloor \frac{n}{2} \right\rfloor - 1 \right) \right\}, \tag{2.1}$$

then  $G$  is Hamilton-connected.

**Proof.** Suppose that  $G$  is not Hamilton-connected. By Lemma 2.1, there exists an integer  $t$  such that  $d_{t-1} \leq t$ , where  $k \leq t \leq \frac{n}{2}$ . Without loss of generality, let  $d(v_i) = d_i$  for  $1 \leq i \leq t - 1$ . The number of edges which are not incident to any vertex in

$\{v_1, v_2, \dots, v_{t-1}\}$  does not exceed  $\binom{n-t+1}{2}$ , and the number of edges incident to any vertex in  $\{v_1, v_2, \dots, v_{t-1}\}$  is at most  $t(t-1)$ . It follows that

$$e(G) \leq \binom{n-t+1}{2} + t(t-1). \tag{2.2}$$

The bound in (2.2) is best possible in the sense that the graph  $M_n^t = K_t \vee (K_{n-2t+1} + (t-1)K_1)$  is not Hamilton-connected. For  $k \leq t \leq \frac{n}{2}$ , by (2.1),

$$\begin{aligned} e(G) &> \max \left\{ \binom{n-k+1}{2} + k(k-1), \binom{\lceil \frac{n}{2} \rceil + 1}{2} + \lfloor \frac{n}{2} \rfloor (\lfloor \frac{n}{2} \rfloor - 1) \right\} \\ &\geq \binom{n-t+1}{2} + t(t-1), \end{aligned}$$

contrary to (2.2). Hence  $G$  must be Hamilton-connected.  $\square$

In [2], Bondy and Chvátal introduced the closure concept which plays an important role in cycle theory. For a graph  $G$  of order  $n$  and an integer  $k = k(n) > 0$ , the  $k$ -closure of  $G$ , denoted by  $C_k(G)$ , is obtained from  $G$  by sequentially joining pairs of nonadjacent vertices whose degree sum is at least  $k$  until no such vertex pairs exist.

**Theorem 2.3.** (Bondy and Chvátal [2]) *A graph  $G$  is Hamilton-connected if and only if  $C_{n+1}(G)$  is Hamilton-connected.*

**Lemma 2.2.** *Let  $k \geq 2$  be an integer,  $G$  be a graph of order  $n \geq 6k$ , and  $G = C_{n+1}(G)$ . Let  $\omega(G)$  denote the clique number of  $G$ . If*

$$e(G) > \binom{n-k}{2} + k(k+1),$$

then  $\omega(G) \geq n - k + 1$ .

**Proof.** It suffices to show that  $G$  contains a clique  $C$  with  $|C| \geq n - k + 1$ . Define

$$F = \left\{ u \in V(G) : d_G(u) \geq \frac{n+1}{2} \right\}.$$

As  $G = C_{n+1}(G)$ ,  $F$  is a clique, and so there exists a maximal clique  $C$  of  $G$  with  $F \subseteq C$ . Let  $s = |C|$  and  $H = G - C$ . As  $C$  is a maximal clique and as  $F \subseteq C$ , for any  $v \in V(H)$ , we have

$$d_C(v) \leq s - 1 \text{ and } d_G(v) \leq \frac{n}{2}. \tag{2.3}$$

**Claim 1.**  $s \geq \frac{n}{3} + k + 1$ .

By contradiction, we assume that  $s < \frac{n}{3} + k + 1$ . It follows by  $|V(H)| = n - s$  and by (2.3) that

$$\begin{aligned} e(H) + e(V(H), C) &= \frac{\sum_{v \in V(H)} d_G(v) + \sum_{v \in V(H)} |N_C(v)|}{2} = \frac{\sum_{v \in V(H)} d_G(v) + \sum_{v \in V(H)} d_C(v)}{2} \\ &\leq \frac{(n-s)\frac{n}{2} + (n-s)(s-1)}{2} = \frac{(n-s)(n+2s-2)}{4}. \end{aligned} \tag{2.4}$$

As  $C$  is a clique,  $e(G[C]) = \binom{s}{2}$  and so by (2.4) and by  $s < \frac{n}{3} + k + 1$ , we have

$$\begin{aligned} e(G) &= e(G[C]) + e(H) + e(V(H), C) \leq \binom{s}{2} + \frac{(n-s)(n+2s-2)}{4} \\ &= \frac{n(n+s-2)}{4} < \frac{n(n+\frac{n}{3}+k+1-2)}{4} \\ &= \frac{1}{3}n^2 + \frac{k-1}{4}n \leq \binom{n-k}{2} + k(k+1) < e(G), \end{aligned}$$

a contradiction. Hence Claim 1 must hold.

**Claim 2.**  $s \geq n - k + 1$ .

By contradiction, we assume that  $s \leq n - k$ . Since  $G = C_{n+1}(G)$ , if  $uv \notin E(G)$ , then  $d_G(u) + d_G(v) \leq n$ . As  $C$  is a clique, every vertex  $u \in C$  satisfies  $d_G(u) \geq s - 1$ . For each  $v \in V(H)$ , as  $v \notin C$ , we have  $d_G(v) + d_G(u) \leq n$ , and so  $d_G(v) \leq n - d_G(u) \leq n - s + 1$ . As  $H = G - C$ , we have  $\sum_{v \in V(H)} d_G(v) = 2e(H) + e(V(H), C)$ . Thus

$$e(H) + e(V(H), C) = \frac{\sum_{v \in V(H)} d_H(v)}{2} + \sum_{v \in V(H)} d_C(v) \leq \sum_{v \in V(H)} d_G(v) \leq (n-s)(n-s+1),$$

and so

$$e(G) = e(G[C]) + e(H) + e(V(H), C) \leq \binom{s}{2} + (n-s)(n-s+1) = \frac{3}{2}s^2 - \left(2n + \frac{3}{2}\right)s + n^2 + n.$$

Let  $f(x) = \frac{3}{2}x^2 - (2n + \frac{3}{2})x + n^2 + n$ . It is routine to show that  $f(x)$  is increasing on  $x$  for  $x \geq \frac{2}{3}n + \frac{1}{2}$  and decreasing on  $x$  for  $x \leq \frac{2}{3}n + \frac{1}{2}$ . As  $f(n-k) = f(\frac{n}{3} + k + 1) = \binom{n-k}{2} + k(k+1)$ , it follows that

$$e(G) = \frac{3}{2}s^2 - (2n + \frac{3}{2})s + n^2 + n \leq \binom{n-k}{2} + k(k+1) < e(G),$$

a contradiction. Hence Claim 2 holds and so  $\omega(G) \geq s \geq n-k+1$ .  $\square$

**Theorem 2.4.** (Dirac [6]) *If a simple graph  $G$  has minimum degree  $d > 1$ , then  $G$  contains a cycle of length at least  $d + 1$ .*

Let  $H$  be a subgraph of a graph  $G$ . Define the **vertices of attachment** of  $H$  in  $G$  to be the vertex set:

$$A_G(H) = \{v \in V(H) : \exists u \in V(G) - V(H) \text{ such that } uv \in E(G)\}.$$

**Theorem 2.5.** *Let  $k \geq 3$  be an integer,  $G$  be a graph with order  $n \geq 6k$  and  $\delta(G) \geq k$ . Suppose that*

$$e(G) > \binom{n-k}{2} + k(k+1).$$

*Then  $G$  is Hamilton-connected if and only if  $C_{n+1}(G) \notin \{L_n^k, M_n^k\}$ .*

**Proof.** Let  $G' = C_{n+1}(G)$ . By Theorem 2.3,  $G$  is Hamilton-connected if and only if  $G'$  is Hamilton-connected. It is routine to verify that neither  $L_n^k$  nor  $M_n^k$  is Hamilton-connected, and so it remains to assume that  $G' \notin \{L_n^k, M_n^k\}$  to prove that  $G'$  is Hamilton-connected.

We argue by contradiction and assume that  $G' \notin \{L_n^k, M_n^k\}$  and  $G'$  is not Hamilton-connected. Since  $\delta(G') \geq \delta(G) \geq k$  and  $e(G') \geq e(G)$ , it follows by Lemma 2.2 that  $\omega(G') \geq n-k+1$ . Let  $C$  be a maximum clique of  $G'$ ,  $w = \omega(G')$  and  $H = G' - C$ .

**Claim 1:** Each of the following holds:

- (i) If  $v \in C$  satisfying  $d_{G'}(v) \geq \omega(G')$ , then for any  $u \in V(H)$ ,  $uv \in E(G')$ .
- (ii) There is no vertex  $u \in V(H)$  satisfying  $d_{G'}(u) \geq n - \omega(G') + 2$ .

Let  $v \in C$  be a vertex with  $d_{G'}(v) \geq \omega(G')$ . For any  $u \in V(H)$ , by  $\delta(G') \geq \delta(G) \geq k$  and by Lemma 2.2, we have  $d_{G'}(v) + d_{G'}(u) \geq \omega(G') + k \geq n - k + 1 + k = n + 1$ , and so as  $G' = C_{n+1}(G)$ , we have  $uv \in E(G')$ . This proves (i).

We argue by contradiction to prove (ii) and assume that there exists a vertex  $u \in V(H)$  satisfying  $d_{G'}(u) \geq n - \omega(G') + 2$ , then as  $C$  is a clique, for any vertex  $v \in C$ ,  $d_{G'}(v) \geq \omega(G') - 1$ . Hence  $d_{G'}(v) + d_{G'}(u) \geq \omega(G') - 1 + n - \omega(G') + 2 = n + 1$ . It follows by  $G' = C_{n+1}(G)$  that  $uv \in E(G')$ , and so every vertex in  $H$  is adjacent to every vertex in  $C$ , contrary to the fact that  $C$  is a maximum clique. This verifies Claim 1(ii), and so Claim 1 is justified.

**Claim 2:**  $\omega(G') = n - k + 1$  and for any vertex  $u \in V(H)$ ,  $d_{G'}(u) = k$ .

If  $\omega(G') \geq n - k + 2$ , then for any vertex  $v \in C$ ,  $d_{G'}(v) \geq n - k + 1$ . Since  $G'$  is not Hamilton-connected,  $G'$  is not a clique, and so  $V(H) \neq \emptyset$ . For any vertex  $u \in V(H)$ , as  $d_{G'}(u) \geq \delta(G) \geq k$ , we obtain a contradiction to Claim 1. Hence by Lemma 2.2, we have  $\omega(G') = n - k + 1$ . As  $\omega(G') = n - k + 1$  and  $\delta(G) \geq k$ , it follows by Claim 1 that for any vertex  $u \in V(H)$ ,  $d_{G'}(u) = k$ . This justifies Claim 2.

Denote  $F = A_{G'}(C) = \{u_1, u_2, \dots, u_s\}$ . As  $C$  is a maximum clique of  $G'$ , it follows from Claim 2 that  $d_{G'}(u_i) \geq n - k + 1 = \omega(G')$ , and so by Claim 1(i) that  $d_{G'}(u_i) = n - 1$ . This implies that for any  $u \in V(H)$ ,  $A_{G'}(C) \subseteq N_{G'}(u)$ , and so by Claim 2,  $s \leq k$ . As  $\omega(G') = n - k + 1$ , for any  $u \in H$ ,  $d_H(u) \leq k - 2$ . Hence  $2 \leq s \leq k$ .

By inspection, if  $s = 2$ , then  $G' = K_2 \vee (K_{k-1} + K_{n-k-1}) = L_n^k$ ; and if  $s = k$ , then  $G' = K_k \vee (K_{n-2k+1} + (k-1)K_1) = M_n^k$ . As we assume that  $G' \notin \{L_n^k, M_n^k\}$ , we must have  $3 \leq s \leq k - 1$ .

For any vertex  $u \in V(H)$ , since  $|A_{G'}(C)| = s$  and  $A_{G'}(C) \subseteq N_{G'}(u)$ , we have  $d_H(u) = k - s$ . By Theorem 2.4,  $H$  has a cycle  $C_1$  with  $q = |C_1| \geq k - s + 1$ . Let  $C_1 = x_1x_2 \cdots x_qx_1$ . For any pair of distinct vertices  $x_i$  and  $x_j$  on  $C_1$  with  $i \neq j$ , we use  $x_iC_1x_j$  ( $x_i \overleftarrow{C_1} x_j$ ) to denote the subpath  $x_ix_{i+1} \cdots x_j$  ( $x_ix_{i-1} \cdots x_{j+1}x_j$ ) on  $C_1$ , where the subscripts are taken modulo  $q$ .

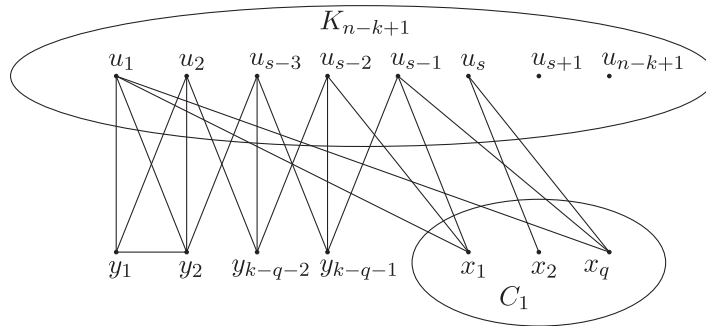
In the rest of the proof of this theorem, we denote  $V(H - C_1) = \{y_1, y_2, \dots, y_{k-q-1}\}$  and  $C - F = \{u_{s+1}, u_{s+2}, \dots, u_{n-k+1}\}$ .

To obtain a contradiction, we are to show that a Hamiltonian  $(v, w)$ -path always exists in  $G'$  for any  $v, w \in V(G')$ . If  $v, w \in F$ , then without loss of generality, we assume that  $v = u_1$  and  $w = u_s$ . If  $|C_1| = q \geq k - s + 2$ , then  $k - q \leq s - 2$ , and so  $u_1y_1u_2y_2 \cdots u_{k-q-1}y_{k-q-1}u_{k-q}x_1 \overrightarrow{C_1} x_q u_{k-q+1} \cdots u_{s-1}u_{s+1} \cdots u_{n-k+1}u_s$  is a Hamiltonian  $(v, w)$ -path.

Hence we may assume that  $|C_1| = q = k - s + 1$ , or equivalently,  $k - q = s - 1$ . If there exists no edge in  $G'$  linking a vertex in  $V(C_1)$  to a vertex  $V(H - C_1)$ , then  $E(H - C_1) \neq \emptyset$ . By symmetry, we assume that  $y_1y_2 \in E(H - C_1)$ . In this case,

**Table 1**  
The existence of a Hamiltonian  $(v, w)$ -path.

Cases	$v$	$w$	Hamiltonian $(v, w)$ -path in $G'$
$v, w \in V(C_1)$	$x_i$	$x_j$	$x_i C_1 x_{j-1} u_1 y_1 \cdots u_{k-q-1} u_{k-q} \cdots u_{s-1}$ $u_{s+1} \cdots u_{n-k+1} u_s x_{i-1} \tilde{C}_1 x_j$
$v \in V(C_1), w \in F$	$x_1$	$u_1$	$x_1 C_1 x_q u_s u_{s+1} \cdots u_{n-k+1} u_{k-q} \cdots u_{s-1} y_{k-q-1}$ $u_{k-q-1} \cdots y_2 u_2 y_1 u_1$
$v \in V(C_1), w \in V(H - C_1)$	$x_1$	$y_1$	$x_1 C_1 x_q u_s u_{s+1} \cdots u_{n-k+1} u_{k-q} \cdots u_{s-1} u_{k-q-1}$ $y_{k-q-1} \cdots u_2 y_2 u_1 y_1$
$v \in V(C_1), w \in C - F$	$x_1$	$u_{n-k+1}$	$x_1 C_1 x_q u_1 y_1 u_2 y_2 \cdots u_{k-q-1} y_{k-q-1} u_{k-q} \cdots u_{s-1} u_s$ $u_{s+1} \cdots u_{n-k+1}$
$v \in F, w \in C - F$	$u_1$	$u_{n-k+1}$	$u_1 y_1 u_2 y_2 \cdots u_{k-q-1} y_{k-q-1} u_{k-q} \cdots u_{s-1} x_1 C_1 x_q$ $u_s u_{s+1} \cdots u_{n-k+1}$
$v \in F, w \in V(H - C_1)$	$u_1$	$y_1$	$y_1 u_2 y_2 u_3 y_3 \cdots u_{k-q-1} y_{k-q-1} u_{k-q} \cdots u_{s-1} x_1 C_1 x_q$ $u_s u_{s+1} \cdots u_{n-k+1} u_1$
$v \in V(H - C_1), w \in V(H - C_1)$	$y_1$	$y_{k-q-1}$	$y_1 u_1 y_2 u_2 \cdots y_{k-q-2} u_{k-q-2} u_{k-q} \cdots u_{s-1} x_1 C_1 x_q$ $u_s \cdots u_{n-k+1} u_{k-q-1} y_{k-q-1}$
$v \in V(H - C_1), w \in C - F$	$y_1$	$u_{n-k+1}$	$y_1 u_1 y_2 u_2 \cdots y_{k-q-1} u_{k-q-1} u_{k-q} \cdots u_{s-1} x_1 C_1 x_q$ $u_s \cdots u_{n-k+1}$
$v \in C - F, w \in C - F$	$u_{s+1}$	$u_{n-k+1}$	$u_{s+1} u_1 y_1 u_2 y_2 \cdots u_{k-q-1} y_{k-q-1} u_{k-q} \cdots u_{s-1}$ $x_1 P x_q u_s u_{s+2} \cdots u_{n-k+1}$



**Fig. 1.** The underlying graph to constructing  $P_i$  ( $i = 1, \dots, 10$ ).

$u_s \cdots u_{n-k+1} u_{s-1} x_1 C_1 x_q u_{k-q-1} y_{k-q-1} \cdots u_2 y_2 y_1 u_1$  is a Hamiltonian  $(v, w)$ -path. Hence we assume that there exists an edge  $x_1 y_1$  (say) linking  $V(C_1)$  to  $V(H - C_1)$ . Then  $u_s \cdots u_{n-k+1} u_{s-1} y_1 x_1 C_1 x_q u_{k-q-1} y_{k-q-1} \cdots u_2 y_2 u_1$  is a Hamiltonian  $(v, w)$ -path.

Therefore, in the discussions below, we assume that  $|\{v, w\} \cap F| \leq 1$ . As  $V(G)$  is partitioned into  $F, C - F, V(C_1)$  and  $V(H - C_1)$ , Table 1 indicates that for any other choices of  $v, w \in V(G')$ , by symmetry,  $G'$  always have a Hamiltonian  $(v, w)$ -path (Fig. 1).

As for any  $v, w \in V(G')$ , we have shown that  $G'$  always has a Hamiltonian  $(v, w)$ -path, leading to contradiction to the assumption that  $G'$  is not Hamilton-connected. This contradiction completes the proof of the theorem.  $\square$

### 3. Spectral radius and Hamilton-connected graphs

The goal of this section is to show a relationship between the spectral radius of a graph  $G$  and the Hamilton-connectedness of  $G$ .

Given two distinct vertices  $u, v$  in a graph  $G$ , obtain a new graph  $G' = G'(u, v)$  by replacing all edges  $vw$  by  $uw$  for each  $w \in N_G(v) \setminus (N_G(u) \cup \{u\})$ . This operation is called the *Kelmans transformation* [11]. We start with some lemmas.

**Lemma 3.1.** (Hong et al. [10], and Nikiforov [12]) Let  $G$  be a graph of order  $n$  with the minimum degree  $\delta \geq k$ . Then

$$\lambda(G) \leq \frac{k - 1 + \sqrt{(k + 1)^2 + 4(2e(G) - nk)}}{2}.$$

**Lemma 3.2.** (Csikvári [4]) Let  $G$  be a graph and  $G'$  be the graph obtained from  $G$  by some Kelmans transformation. Then

$$\lambda(G) \leq \lambda(G').$$

Since  $K_{n-k+1}$  is a proper subgraph of both  $L_n^k$  and  $M_n^k$ , it follows that  $\lambda(L_n^k) > \lambda(K_{n-k+1}) = n - k$  and  $\lambda(M_n^k) > \lambda(K_{n-k+1}) = n - k$ . Motivated by the ideas in [13] and [9], we establish the following theorem.

**Theorem 3.1.** Let  $G$  be a graph of order  $n \geq \max\{6k, \frac{1}{2}k^3 - \frac{1}{2}k^2 + k + 4\}$  with the minimum degree  $\delta \geq k \geq 3$ .

(i) If  $G$  is a subgraph of  $L_n^k$ , then  $\lambda(G) < n - k$  if and only if  $G \neq L_n^k$ .

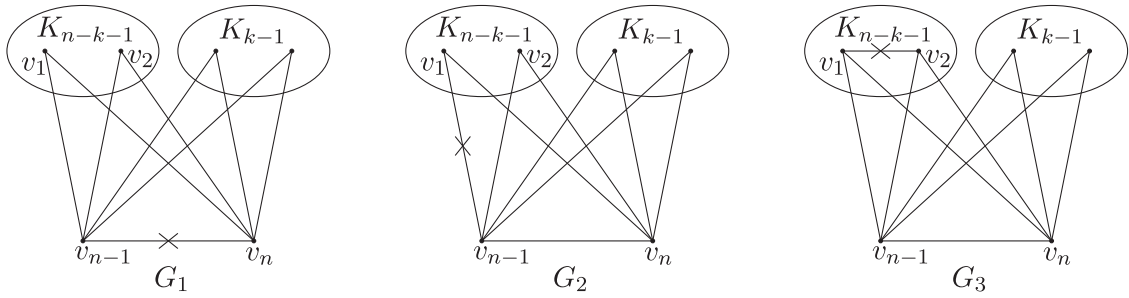


Fig. 2. The graphs obtained from  $L_n^k$  by deleting one edge.

(ii) If  $G$  is a subgraph of  $M_n^k$ , then  $\lambda(G) < n - k$  if and only if  $G \neq M_n^k$ .

**Proof.** As we have observed above, both  $\lambda(L_n^k) > \lambda(K_{n-k+1}) = n - k$  and  $\lambda(M_n^k) > \lambda(K_{n-k+1}) = n - k$ , it suffices to assume that  $G \neq L_n^k$  to prove  $\lambda(G) < n - k$  in (i); and that  $G \neq M_n^k$  to prove  $\lambda(G) < n - k$  in (ii).

Let  $\mathbf{x} = (x_1, x_2, \dots, x_n)^T$  be a positive unit eigenvector of  $\lambda(G)$ . By Rayleigh's quotient inequality [5],

$$\lambda(G) = \mathbf{x}^T A(G) \mathbf{x} = \langle A(G) \mathbf{x}, \mathbf{x} \rangle.$$

Throughout the rest of the proof, we often use  $\lambda$  for  $\lambda(G)$ , when  $G$  is understood from the context.

We argue by contradiction to prove (i), and assume that

$$\lambda(G) \geq n - k \text{ and } G \neq L_n^k. \tag{3.5}$$

Then  $G$  is a proper subgraph of  $L_n^k$ . Clearly, we only need to consider  $G$  with the maximum spectral radius which can be obtained from  $L_n^k$  by deleting one edge. By symmetry, there are only three such graphs:  $G_1 = L_n^k - v_{n-1}v_n$ ,  $G_2 = L_n^k - v_1v_{n-1}$  and  $G_3 = L_n^k - v_1v_2$ .

We claim that  $\lambda(G_1) \leq \lambda(G_2) \leq \lambda(G_3)$ . Using the notation in Fig. 2, let  $u = v_n, v = v_1$  in  $G_1$ . Thus  $N_G(v) \setminus (N_G(u) \cup \{u\}) = \{v_{n-1}\}$ . Let  $G'_1 = G'_1(u, v)$  be a Kelmans transformation of  $G_1$ . Then  $G'_1 = G_2$ . By Lemma 3.2,  $\lambda(G_1) \leq \lambda(G'_1) = \lambda(G_2)$ . Now let  $u = v_{n-1}, v = v_2$  in  $G_2$ . We have  $N_{G_2}(v) \setminus (N_{G_2}(u) \cup \{u\}) = \{v_1\}$ . Then the Kelmans transformation  $G'_2 = G'_2(u, v)$  is isomorphic to  $G_3$ , and so Lemma 3.2,  $\lambda(G_2) \leq \lambda(G'_2) = \lambda(G_3)$ . This justifies the claim.

Define  $Z = \{v \in V(L_n^k) : d_{L_n^k}(v) = n - 1\}$ ,  $X = \{v \in V(L_n^k) : d_{L_n^k}(v) = n - k\}$ , and  $Y = \{v \in V(L_n^k) : d_{L_n^k}(v) = k\}$ . Hence, it suffices to assume  $G = L_n^k - uv$  for some edge  $uv$  with  $\{u, v\} \subset X$  to prove (i). Therefore, in the rest of the proof for (i), we shall have such an assumption. Let  $\mathbf{x} = (x_1, x_2, \dots, x_n)^T$  denote a positive unit eigenvector of  $\lambda(G)$ , and define

$$\begin{aligned} x &= x_i, v_i \in X \setminus \{u, v\}, \\ y &= x_j, v_j \in Y, \\ z &= x_k, v_k \in Z, \\ s &= x_u = x_v. \end{aligned}$$

Then the  $n$  eigenequations of  $G$  can be reduced to the following four equations:

$$\begin{aligned} \lambda x &= (n - k - 4)x + 2z + 2s, \\ \lambda y &= (k - 2)y + 2z, \\ \lambda z &= (n - k - 3)x + (k - 1)y + z + 2s, \\ \lambda s &= (n - k - 3)x + 2z. \end{aligned}$$

It follows from algebraic manipulations that

$$\begin{aligned} y &= \frac{2}{\lambda - k + 2}z, \\ x &= \left(1 - \frac{2(k - 1)}{(\lambda + 1)(\lambda - k + 2)}\right)z, \\ s &= \frac{\lambda + 1}{\lambda + 2} \left(1 - \frac{2(k - 1)}{(\lambda + 1)(\lambda - k + 2)}\right)z. \end{aligned}$$

By the definition of  $G$ , we have

$$G - \{yz : y \in Y \text{ and } z \in Z\} + uv \cong K_{n-k+1} + K_{k-1}.$$

Let  $\mathbf{x}'$  be the restriction of  $\mathbf{x}$  to  $K_{n-k+1}$ , then

$$\begin{aligned} \langle A(K_{n-k+1})\mathbf{x}', \mathbf{x}' \rangle &= \langle A(G)\mathbf{x}, \mathbf{x} \rangle + 2s^2 - 4(k-1)yz - (k-1)(k-2)y^2 \\ &= \lambda + 2s^2 - 4(k-1)yz - (k-1)(k-2)y^2. \end{aligned}$$

By Rayleigh's quotient inequality,

$$\frac{\langle A(K_{n-k+1})\mathbf{x}', \mathbf{x}' \rangle}{\|\mathbf{x}'\|^2} < \lambda(K_{n-k+1}) = n - k.$$

By (3.5),  $\lambda \geq n - k$ . This, together with  $\|\mathbf{x}'\|^2 = \|\mathbf{x}\|^2 - (k-1)y^2$ , implies that

$$2s^2 + \lambda(k-1)y^2 < 4(k-1)yz + (k-1)(k-2)y^2. \tag{3.6}$$

Since  $k \geq 3$  and  $\lambda \geq n - k \geq 5k > k - 2$ , we have

$$\begin{aligned} s^2 &= \left(\frac{\lambda + 1}{\lambda + 2}\right)^2 \left(1 - \frac{2(k-1)}{(\lambda + 1)(\lambda - k + 2)}\right)^2 z^2 \\ &> \left(1 - \frac{2}{\lambda + 2}\right) \left(1 - \frac{4(k-1)}{(\lambda + 1)(\lambda - k + 2)}\right) z^2 \\ &> \left(1 - \frac{2}{\lambda + 2} - \frac{4(k-1)}{(\lambda + 1)(\lambda - k + 2)}\right) z^2 \\ &> \left(1 - \frac{2}{\lambda - k + 2} - \frac{1}{\lambda - k + 2}\right) z^2 = \left(\frac{\lambda - k - 1}{\lambda - k + 2}\right) z^2 \\ &> \left(\frac{4k - 4}{\lambda - k + 2}\right) z^2 = 2(k-1)yz, \end{aligned}$$

and  $\lambda(k-1)y^2 > (k-1)(k-2)y^2$ . It follows that

$$2s^2 + \lambda(k-1)y^2 > 4(k-1)yz + (k-1)(k-2)y^2,$$

contrary to (3.6). This completes the proof of (i).

The proof for (ii) follows a similar proving strategy as in that of (i), and so we also argue by contradiction. Assume that

$$\lambda(G) \geq n - k \text{ and } G \neq M_n^k. \tag{3.7}$$

Then  $G$  is a proper subgraph of  $M_n^k$ .

Define  $Z = \{v \in V(M_n^k) : d_{M_n^k}(v) = n - 1\}$ ,  $X = \{v \in V(M_n^k) : d_{M_n^k}(v) = n - k\}$ , and  $Y = \{v \in V(M_n^k) : d_{M_n^k}(v) = k\}$ .

As in the proof of (i), we only need to consider the case that  $G = M_n^k - uv$  for an edge  $uv$  with  $\{u, v\} \subset X$ . Let  $\mathbf{x} = (x_1, x_2, \dots, x_n)^T$  denote a positive unit eigenvector of  $\lambda(G)$ , and define

$$\begin{aligned} x &= x_i, v_i \in X \setminus \{u, v\}, \\ y &= x_j, v_j \in Y, \\ z &= x_k, v_k \in Z, \\ s &= x_u = x_v. \end{aligned}$$

Then the  $n$  eigenequations of  $G$  can be reduced to the following four equations:

$$\begin{aligned} \lambda x &= (n - 2k - 2)x + kz + 2s, \\ \lambda y &= kz, \\ \lambda z &= (n - 2k - 1)x + (k - 1)y + (k - 1)z + 2s, \\ \lambda s &= (n - 2k - 1)x + kz. \end{aligned}$$

It follows by algebraic manipulations that

$$\begin{aligned} y &= \frac{k}{\lambda}z, \\ x &= \left(1 - \frac{k(k-1)}{\lambda(\lambda+1)}\right)z, \\ s &= \frac{\lambda+1}{\lambda+2} \left(1 - \frac{k(k-1)}{\lambda(\lambda+1)}\right)z. \end{aligned}$$

By the definition of  $G$ , we have

$$G - \{yz : y \in Y \text{ and } z \in Z\} + uv \cong K_{n-k+1} + (k-1)K_1.$$

Let  $\mathbf{x}'$  be the restriction of  $\mathbf{x}$  to  $K_{n-k+1}$ , then  $\|\mathbf{x}'\|^2 = \|\mathbf{x}\|^2 - (k-1)y^2 = 1 - (k-1)y^2$ . By Rayleigh's quotient inequality,

$$\frac{\langle A(K_{n-k+1})\mathbf{x}', \mathbf{x}' \rangle}{\|\mathbf{x}'\|^2} < \lambda(K_{n-k+1}) = n - k.$$

By (3.7),  $\lambda \geq n - k$ , and so,

$$\langle A(K_{n-k+1})\mathbf{x}', \mathbf{x}' \rangle < \lambda(1 - (k-1)y^2). \tag{3.8}$$

Since

$$\langle A(K_{n-k+1})\mathbf{x}', \mathbf{x}' \rangle = \langle A(G)\mathbf{x}, \mathbf{x} \rangle + 2s^2 - 2k(k-1)yz = \lambda + 2s^2 - 2k(k-1)yz, \tag{3.9}$$

by (3.8) and (3.9),

$$s^2 + \frac{k-1}{2}\lambda y^2 - k(k-1)yz < 0. \tag{3.10}$$

Since  $\lambda \geq n - k \geq \frac{1}{2}k^3 - \frac{1}{2}k^2 + 4$ ,

$$\begin{aligned} & \lambda \left( s^2 + \frac{k-1}{2}\lambda y^2 - k(k-1)yz \right) \\ &= \left( \frac{\lambda+1}{\lambda+2} \right)^2 \left( 1 - \frac{k(k-1)}{\lambda(\lambda+1)} \right)^2 \lambda z^2 + \frac{k^2(k-1)}{2} z^2 - k^2(k-1)z^2 \\ &> \left( 1 - \frac{2}{\lambda+2} \right) \left( 1 - \frac{2k(k-1)}{\lambda(\lambda+1)} \right) \lambda z^2 - \frac{k^2(k-1)}{2} z^2 \\ &> \left( \lambda - \frac{2\lambda}{\lambda+2} - \frac{2k(k-1)}{\lambda+1} - \frac{k^2(k-1)}{2} \right) z^2 \\ &> \left( \lambda - 2 - \frac{2k^2}{\frac{1}{2}k^3 - \frac{1}{2}k^2 + 5} - \frac{k^2(k-1)}{2} \right) z^2 \\ &> \left( \frac{1}{2}k^3 - \frac{1}{2}k^2 + 2 - \frac{2k^2}{\frac{1}{2}k^3 - \frac{1}{2}k^2} - \frac{k^3}{2} + \frac{k^2}{2} \right) z^2 \\ &= \left( 2 - \frac{4}{k-1} \right) z^2. \end{aligned}$$

Let  $f(x) = 2 - \frac{4}{x-1}$ . It is routine to show that  $f(x)$  is increasing on  $x$  for all real numbers  $x \geq 2$ . If  $k \geq 3$  and  $k$  is an integer, then we have  $f(k) \geq f(3) = 0$ . It follows that

$$s^2 + \frac{k-1}{2}\lambda y^2 - k(k-1)yz > 0,$$

contrary to (3.10). This completes the proof of (ii).  $\square$

**Theorem 3.2.** Let  $k \geq 3$  be an integer, and let  $G$  be a graph with order  $n \geq \max\{6k, \frac{1}{2}k^3 - \frac{1}{2}k^2 + k + 4\}$ ,  $\delta(G) \geq k$  and spectral radius  $\lambda(G) \geq n - k$ . Then  $G$  is Hamilton-connected if and only if  $G \notin \{L_n^k, M_n^k\}$ .

**Proof.** It routine to verify that if  $G \in \{L_n^k, M_n^k\}$ , then  $G$  is not Hamilton-connected. Therefore, in the rest of the proof, we assume that  $G \in \{L_n^k, M_n^k\}$  to prove that  $G$  is Hamilton-connected. To the end of this section, we set  $\lambda = \lambda(G)$  and  $\delta = \delta(G)$ .

Since  $\lambda \geq n - k$  and by Lemma 3.1, we have

$$n - k \leq \lambda \leq \frac{k-1 + \sqrt{(k+1)^2 + 4(2e(G) - nk)}}{2},$$

which implies that

$$e(G) \geq \frac{n^2 - (2k-1)n + 2k^2 - 2k}{2}. \tag{3.11}$$

Since  $n \geq \max\{6k, \frac{1}{2}k^3 - \frac{1}{2}k^2 + k + 4\} \geq \frac{k^2 + 5k + 2}{2}$ , it follows from (3.11) that

$$e(G) \geq \frac{n^2 - (2k-1)n + 2k^2 - 2k}{2} \geq \binom{n-k}{2} + k(k+1) + 1, \tag{3.12}$$

where the equality in (3.12) holds if and only if  $n = \frac{k^2 + 5k + 2}{2}$ .



Since  $n \geq \max\{6k, \frac{1}{2}k^3 - \frac{1}{2}k^2 + k + 4\}$  and since  $G \notin \{L_n^k, M_n^k\}$ , it follows by [Theorem 2.5](#) (when  $n \geq \max\{6k, \frac{1}{2}k^3 - \frac{1}{2}k^2 + k + 4\}$ ) or by [Theorem 3.1](#) (when  $\lambda(G) \geq n - k$ ) then  $G$  is Hamilton-connected. The proof is complete.  $\square$

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