

On k -Maximal Strength Digraphs

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Abstract: Let $k > 0$ be an integer and let D be a simple digraph on $n > k$ vertices. We prove that if

$$|A(D)| > k(2n - k - 1) + \binom{n - k}{2},$$

then D must have a nontrivial subdigraph H such that the strong arc connectivity of H is at least $k + 1$. We also show that this bound is best possible and present a constructive characterization for the extremal graphs. © 2015

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1. INTRODUCTION

We consider finite simple graphs and simple digraphs. Usually, we use G to denote a graph and D a digraph. Undefined terms and notation will follow [3] for graphs and [2] for digraphs. In particular, $\kappa'(G)$ denotes the edge connectivity of a graph G and $\lambda(D)$ denotes the arc-strong connectivity of a digraph D . If G is a simple graph, then G^c denotes the complement of G . If $X \subseteq E(G^c)$, then $G + X$ is the simple graph with vertex set $V(G)$ and edge set $E(G) \cup X$. We will use $G + e$ for $G + \{e\}$. Likewise, if D is a simple digraph, let D^c denotes the complement of D . For $X \subseteq A(D^c)$ and $e \in A(D^c)$, we similarly define the simple digraphs $D + X$ and $D + e$, respectively. Throughout this article, we use the notation (u, v) to denote an arc oriented from u to v in a digraph. If $W \subseteq V(D)$ or if $W \subseteq A(D)$, then $D[W]$ denotes the subdigraph of D induced by W . For $v \in V(D)$, we use $D - v$ for $D[V(D) - \{v\}]$. For graphs H and G , we denote $H \subseteq G$ when H is a subgraph of G . Similarly, for digraphs H and D , $H \subseteq D$ when H is a subdigraph of D . We use $D \cong H$ when the digraphs D and H are isomorphic.

Given a graph G , Matula [5] first studied the quantity

$$\bar{\kappa}'(G) = \max\{\kappa'(H) : H \subseteq G\}.$$

He called $\bar{\kappa}'(G)$ the **strength** of G . Mader [4] considered an extremal problem related to $\bar{\kappa}'(G)$. For an integer $k > 0$, a simple graph G with $|V(G)| \geq k + 1$ is **k -maximal** if $\bar{\kappa}'(G) \leq k$ but for any edge $e \in E(G^c)$, $\bar{\kappa}'(G + e) > k$. In [4], Mader proved the following.

Theorem 1.1. (Mader [4]) *If G is a k -maximal graph on $n > k \geq 1$ vertices, then*

$$|E(G)| \leq (n - k)k + \binom{k}{2}.$$

Furthermore, this bound is best possible.

The purpose of this article is to investigate the same for simple digraphs. Naturally, for a digraph D , we define

$$\bar{\lambda}(D) = \max\{\lambda(H) : H \subseteq D\}.$$

Let $k \geq 0$ be an integer. A simple digraph D with $|V(D)| \geq k + 1$ is k -**maximal** if $\bar{\lambda}(D) \leq k$ but for any arc $e \in A(D^c)$, $\bar{\lambda}(D + e) \geq k + 1$. For positive integer n and k with $n \geq k + 1$, define

$$\mathcal{D}(n, k) = \{D : D \text{ is a simple digraph with } |V(D)| = n \text{ and } D \text{ is } k\text{-maximal}\}.$$

Our goal is to determine $\max\{|A(D)| : D \in \mathcal{D}(n, k)\}$. If $h < k$, we define $\binom{h}{k} = 0$. Our main result is the following.

Theorem 1.2. *Let $k \geq 0$ and $n \geq k + 1$ be nonnegative integers. If $D \in \mathcal{D}(n, k)$, then*

$$|A(D)| \leq k(2n - k - 1) + \binom{n - k}{2}.$$

Furthermore, the bound is best possible.

Corollary 1.3. *Let $k > 0$ be an integer and let D be a simple digraph on $n > k$ vertices. If*

$$|A(D)| > k(2n - k - 1) + \binom{n - k}{2},$$

then D must have a subdigraph H such that $\lambda(H) \geq k + 1$.

The corollary follows immediately from Theorem 1.2. In the next section, we investigate properties of k -maximal digraphs. The main result will be proved in the last section.

2. PROPERTIES OF k -MAXIMAL DIGRAPHS

Throughout this section, let $k \geq 0$ be an integer. Define $\mathcal{D}(k) = \cup_{n \geq k+1} \mathcal{D}(n, k)$. Thus $\mathcal{D}(k)$ is the family of all k -maximal digraphs. Recall that a tournament on n vertices is an orientation of the complete graph K_n on n vertices. The following lemma indicates that the $k = 0$ case has a clear structure.

Lemma 2.1. *([1]) A digraph $D \in \mathcal{D}(0)$ if and only if D is an acyclic tournament.*

The easy proof is left to the reader (and found in [1]).

For any integer $n \geq 0$, let K_n^* denote the complete digraph on n vertices. By definition, we have

$$K_{k+1}^* \in \mathcal{D}(k) \text{ and if } H \in \mathcal{D}(k) \text{ and } |V(H)| = k + 1, \text{ then } H \cong K_{k+1}^*. \tag{1}$$

Following [2], if D is a digraph and if $X, Y \subseteq V(D)$, then define

$$(X, Y)_D = \{(x, y) \in A(D) : x \in X, y \in Y\}.$$

We further define that, for $X \subseteq V(D)$,

$$\partial_D^+(X) = (X, V(D) - X)_D \text{ and } \partial_D^-(X) = (V(D) - X, X)_D.$$

For each $v \in V(D)$, we define

$$N_D^+(v) = \{u \in V(D) : (v, u) \in A(D)\} \text{ and } N_D^-(v) = \{u \in V(D) : (u, v) \in A(D)\}.$$

When the digraph D is understood from the context, we sometimes omit the subscript D in the notation above. By the definition of arc-strong connectivity in [2], a digraph D satisfies $\lambda(D) \geq k$ if and only if for any nonempty proper subset $X \subset V(D)$, $|\partial_D^+(X)| \geq k$.

Lemma 2.2. *If $D \in \mathcal{D}(k)$ and if D is not a complete digraph, then for any proper nonempty subset $X \subset V(D)$ such that $|\partial_D^+(X)| \leq k$, each of the following holds:*

- (i) $(X, V(D) - X)_{D^c} \neq \emptyset$.
- (ii) $|\partial_D^+(X)| = k$.
- (iii) $(V(D) - X, X)_D = \{(y, x) : \text{for any } y \in V(D) - X \text{ and for any } x \in X\}$.

Proof. Let $Y = V(D) - X$. Suppose that $(X, Y)_{D^c} = \emptyset$. Then the arcs in $\partial^+(X)$ induce an underlying complete bipartite graph with a vertex bipartition $\{X, Y\}$. It follows from $|\partial^+(X)| \leq k$ that we must have

$$|X|(n - |X|) = |X| \cdot |Y| \leq k, \text{ and } |X| + |Y| = n \geq k + 1.$$

As the minimum of $|X|(n - |X|)$ must be attained at the boundary point of the domain $1 \leq |X| \leq n - 1$, we observe that $k \leq n - 1 \leq |X|(n - |X|) \leq k$. It follows that we must have $n = k + 1$, and so by (1) that D must be a K_{k+1}^* , contrary to the assumption that D is not a complete digraph. This proves (i).

Therefore there must be an arc $(x, y) \in (X, Y)_{D^c}$. The existence of this arc and the fact that $D \in \mathcal{D}(k)$ imply that $|\partial_D^+(X)| = k$. This proves (ii).

To prove (iii), we again argue by contradiction and assume that for some $x \in X$ and $y \in Y$, $(y, x) \notin A(D)$. Then as $D \in \mathcal{D}(k)$, $\bar{\lambda}(D + (y, x)) \geq k + 1$. It follows that $D + (y, x)$ has a subdigraph H' with $\lambda(H') \geq k + 1$ and with $(y, x) \in A(H')$. Hence both $X \cap V(H') \neq \emptyset$ and $Y \cap V(H') \neq \emptyset$. As $\partial_{H'}^+(X \cap V(H')) \subseteq \partial_D^+(X)$, we have $k + 1 \leq |\partial_{H'}^+(X \cap V(H'))| \leq |\partial_D^+(X)| \leq k$, a contradiction. This proves (iii). ■

Lemma 2.3. *Suppose that $D \in \mathcal{D}(k) - \{K_{k+1}^*\}$, for any proper nonempty subset $X \subset V(D)$ such that $|\partial_D^+(X)| = k$, define $Y = V(D) - X$. One of the following must hold:*

- (i) $|X| = 1$ and $D[Y] \in \mathcal{D}(k)$, and, $|V(D) - X| \geq k + 1$.
- (ii) $|X| \geq k + 1$ and $D[X] \in \mathcal{D}(k)$, and, $|V(D) - X| = 1$.
- (iii) Both $D[X] \in \mathcal{D}(k)$ and $D[Y] \in \mathcal{D}(k)$, and, both $|X| \geq k + 1$ and $|V(D) - X| \geq k + 1$.

Proof. Let $Y = V(D) - X$. We make the following claims.

Claim 1. If $D[X]$ (or $D[Y]$, respectively) is a complete digraph, then $|X| \in \{1, k + 1\}$, (or $|Y| \in \{1, k + 1\}$, respectively).

By symmetry, we prove the case when $D[X]$ is complete. Let $m = |X|$. Then $D[X] = K_m^*$. Since $m - 1 = \lambda(K_m^*) \leq \bar{\lambda}(D) \leq k$, we have $m \leq k + 1$. Assume that $1 < m \leq k$. By Lemma 2.2(i), there exists an arc $(x, y) \in (X, Y)_{D^c}$. As $D \in \mathcal{D}(k)$, $D + (x, y)$ has a subdigraph H with $\lambda(H) \geq k + 1$ and $(x, y) \in A(H)$. Hence $X \cap V(H) \neq \emptyset$. Let $X' = X \cap V(H)$ and $m' = |X'|$, and denote $X' = \{x_1, x_2, \dots, x_{m'}\}$. For each x_i , let k_i denote the number of arcs in $(X', V(H) - X')_H$ incident with x_i . Then $\sum_{i=1}^{m'} k_i = |\partial_H^+(X')| \geq \lambda(H) \geq k + 1$, and so $m'(k + 1) \leq \sum_{i=1}^{m'} d_H^+(x_i) \leq \sum_{i=1}^{m'} (k_i + m' - 1) = (k + 1) + m'(m' - 1)$, which leads to $(m' - 1)(k + 1) \leq m'(m' - 1)$. If $m' > 1$, then $k + 1 \leq m' \leq k$, a contradiction. Therefore, we must have $m' = 1$. Assume that $X' = \{x_1\}$. Then all arcs in $(X, Y)_D$ must be incident with x_1 in X . Since $m \geq 2$, x_2 is not incident with any arcs in $(X, Y)_D$, and so for any $y \in Y$, $(x_2, y) \in A(D^c)$. As $D \in \mathcal{D}(k)$, $D + (x_2, y)$ must

have a subdigraph H'' such that $\lambda(H'') \geq k + 1$ and $(x_2, y) \in A(H'')$. Since $m \leq k$, $d_{H''}^+(x_2) \leq |\partial_{D[X]}^+(x_2) \cup \{x_2, y\}| = (m - 1) + 1 \leq k$, contrary to the assumption that $\lambda(H'') \geq k + 1$. Therefore if $D[X]$ is complete, then $|X| \in \{1, k + 1\}$. This proves Claim 1.

Claim 2. If $D[X]$ (or $D[Y]$, respectively) is not complete digraph, then $D[X] \in \mathcal{D}(k)$, (or $D[Y] \in \mathcal{D}(k)$, respectively).

Again by symmetry, it suffices to show that $D[X] \in \mathcal{D}(k)$. The case for showing $D[Y] \in \mathcal{D}(k)$ is similar and will be omitted. Since $D[X]$ is not complete, $A(D[X]^c) \neq \emptyset$. For any $e \in A(D[X]^c)$, since $A(D[X]^c) \subseteq A(D^c)$ and since $D \in \mathcal{D}(k)$, it follows by definition of $\mathcal{D}(k)$ that $D + e$ has a subdigraph H_e with $e \in A(H_e)$ and with $\lambda(H_e) \geq k + 1$. If for any $e \in A(D[X]^c)$, we always have $V(H_e) \subseteq X$, then by definition of $\mathcal{D}(k)$, we have $D[X] \in \mathcal{D}(k)$ and so $|X| \geq k + 1$. Hence by contradiction, we assume that there exists an arc $(u, v) \in A(D[X]^c)$ such that $D + (u, v)$ has a subdigraph H with $\lambda(H) \geq k + 1$ and with $(u, v) \in A(H)$, and such that $V(H) \cap Y \neq \emptyset$. Then as $u, v \in X$ and as $V(H) \cap Y \neq \emptyset$, we have $\partial_H^+(X \cap V(H)) = (X \cap V(H), Y \cap V(H))_H \subseteq (X, Y)_D$. It follows that $k + 1 \leq \lambda(H) \leq |\partial_H^+(X \cap V(H))| \leq |(X, Y)_D| = |\partial_D^+(X)| = k$, a contradiction. This proves Claim 2.

Claim 3. If $|X| = 1$, then $|Y| \geq k + 1$; and If $|Y| = 1$, then $|X| \geq k + 1$.

By symmetry, we shall assume that $|X| = 1$ to prove $|Y| \geq k + 1$. The other case can be done with symmetric arguments. By contradiction, we assume that $X = \{x\}$ and $Y = \{y_1, y_2, \dots, y_h\}$ with $h \leq k$. By Lemma 2.2 (iii), for each i with $1 \leq i \leq h$, we have $(y_i, x) \in A(D)$. Since D is simple, and since $|(X, Y)_D| = |\partial_D^+(X)| = k$, we must have $h = k$ and $(x, y_i) \in A(D)$ for every i with $1 \leq i \leq k$. It follows that $D = K_{k+1}^*$, contrary to the assumption that $D \neq K_{k+1}^*$. This proves Claim 3.

With these claims, we now to prove the lemma. If $|X| = 1$ or $|Y| = 1$, then by Claim 3, Lemma 2.3 (i) or (ii) must hold. Assume that both $|X| > 1$ and $|Y| > 1$. If $D[X]$ is a complete digraph, then by Claim 1, $D[X] \cong K_{k+1}^* \in \mathcal{D}(k)$. If $D[X]$ is not a complete graph, then by Claim 2, $D[X] \in \mathcal{D}(k)$ as well. Hence if $|X| = 1$, then $D[Y] \in \mathcal{D}(k)$, which implies $|Y| \geq k + 1$. Similarly, in any case, $D[X] \in \mathcal{D}(k)$ and $|X| \geq k + 1$. This proves the lemma. ■

Definition 2.4. Let $H \in \mathcal{D}(k)$ and let $\{v_1, v_2, \dots, v_k\} \subset V(H)$ be a subset of k distinct vertices. Let u be a vertex not in $V(H)$. Define a digraph $[H, K_1]_k$ ($[K_1, H]_k$, respectively) as follows:

- (i) $V([H, K_1]_k) = V([K_1, H]_k) = V(H) \cup \{u\}$.
- (ii) $A([H, K_1]_k) = A(H) \cup \{(v_1, u), (v_2, u), \dots, (v_k, u)\} \cup (\bigcup_{v \in V(H)} \{(u, v)\})$.
 $A([K_1, H]_k) = A(H) \cup \{(u, v_1), (u, v_2), \dots, (u, v_k)\} \cup (\bigcup_{v \in V(H)} \{(v, u)\})$,
 respectively).

Note that each of $[H, K_1]_k$ and $[K_1, H]_k$ represents a family of graphs as the set $\{v_1, v_2, \dots, v_k\} \subset V(H)$ may vary. For notational convenience, we often use $[H, K_1]_k$ ($[K_1, H]_k$, respectively) to denote any member in the family $[H, K_1]_k$ ($[K_1, H]_k$, respectively).

Definition 2.5. Let $H_1, H_2 \in \mathcal{D}(k)$ with H_1 and H_2 being vertex disjoint, and let $\{u_1, u_2, \dots, u_k\} \subset V(H_1)$ be a multiset of $V(H_1)$ and $\{v_1, v_2, \dots, v_k\} \subset V(H_2)$ be a multiset of $V(H_2)$ such that all the arcs $(u_1, v_1), (u_2, v_2), \dots, (u_k, v_k)$ are distinct. Define a digraph $D = D(H_1, H_2; k)$ as follows.

- (i) $V(D) = V(H_1) \cup V(H_2)$.
- (ii) $A(D) = A(H_1) \cup A(H_2) \cup \{(u_1, v_1), (u_2, v_2), \dots, (u_k, v_k)\} \cup \left(\bigcup_{u \in V(H_1), v \in V(H_2)} \{(v, u)\}\right)$.

Let $[H_1, H_2]_k$ be the family of all such digraphs $D = D(H_1, H_2; k)$ defined above. For notational convenience, we often use $[H_1, H_2]_k$ to denote any member in the family $[H_1, H_2]_k$.

Corollary 2.6. *Let $D \in \mathcal{D}(k) - \{K_{k+1}^*\}$ be a digraph. Then there exists a nonempty proper subset $X \subseteq V(D)$ such that one of the following holds.*

- (i) $|X| = 1$, and for some $H \in \mathcal{D}(k)$, $D = [H, K_1]_k$.
- (ii) $|V(D) - X| = 1$ and for some $H \in \mathcal{D}(k)$, $D = [K_1, H]_k$.
- (iii) For some $H_1, H_2 \in \mathcal{D}(k)$, $D = [H_1, H_2]_k$.

Proof. Since $D \in \mathcal{D}(k) - \{K_{k+1}^*\}$, by Lemma 2.2, there exists a proper nonempty subset $X \subset V(D)$ such that $|\partial_D^+(X)| = k$. Let $Y = V(D) - X$. By Lemma 2.3, one of the conclusions of Lemma 2.3 must hold.

If Lemma 2.3(i) holds, then $|X| = 1$ and $D[Y] \in \mathcal{D}(k)$. Let $X = \{x\}$ and $H = D[Y]$. As $|(\{x\}, Y)_D| = k$ and as, by Lemma 2.2(iii), for each $y \in Y$, $(y, x) \in A(D)$, it follows by Definition 2.4 that $D = [K_1, H]_k$, implying Corollary 2.6 (ii). If Lemma 2.3(ii) holds, then $|Y| = 1$ and $D[X] \in \mathcal{D}(k)$. Let $H = D[X]$. With a symmetric argument, Corollary 2.6 (i) holds. If Lemma 2.3(iii) holds, then let $H_1 = D[X]$ and $H_2 = D[Y]$. As $|(X, Y)_D| = k$, and $H_1, H_2 \in \mathcal{D}(k)$, and by Lemma 2.2(iii), $D = [H_1, H_2]_k$. ■

3. THE EXTREMAL FUNCTION

In this section, we shall determine the extremal function as shown in Theorem 3.2 below. This clearly implies Theorem 1.2. We start with a definition.

Definition 3.1. *Let $\mathcal{M}(k)$ be the family of digraphs such that*

- (i) $K_{k+1}^* \in \mathcal{M}(k)$ and such that
- (ii) a digraph $D \neq K_{k+1}^*$ is in $\mathcal{M}(k)$ if and only if for some $H \in \mathcal{M}(k)$, $D = [H, K_1]_k$ or $D = [K_1, H]_k$.

Theorem 3.2. *Let $D \in \mathcal{D}(n, k)$. Then*

$$|A(D)| \leq k(2n - k - 1) + \binom{n - k}{2}. \tag{2}$$

Furthermore, each of the following holds.

- (i) If $k = 1$, then every digraph $D \in \mathcal{D}(1)$ satisfies equality in (2).
- (ii) If $k \geq 2$, then a digraph $D \in \mathcal{D}(k)$ satisfies equality in (2) if and only if $D \in \mathcal{M}(k)$.

Proof. By Definition 3.1, every digraph D in $\mathcal{M}(k)$ satisfies $\bar{\lambda}(D) \leq k$. It is routine to verify that every digraphs $D \in \mathcal{M}(k)$ is k -maximal. A straightforward inductive computation shows that if $D \in \mathcal{M}(k)$ with $n = |V(D)| \geq k + 1$, then

$$|A(D)| = k(2n - k - 1) + \binom{n - k}{2}. \tag{3}$$

To proceed the proof of the theorem, we first prove (2) by induction on $n = |V(D)|$. If $n = k + 1$, then $D = K_{k+1}^*$. As $|A(K_{k+1}^*)| = k(k + 1)$, we observe that (2) holds when $n = k + 1$. Assume that $n > k + 1$ and (2) holds for smaller values of n . Since $n > k + 1$, D cannot be a complete digraph. By Corollary 2.6, we make the following claims.

Claim 1. If Corollary 2.6 (i) or (ii) holds, then (2) holds as well.

Without loss of generality, we assume that $D = [H, K_1]_k$ for some $H \in \mathcal{D}(k)$ with $V(K_1) = \{v\}$. As $|V(H)| = n - 1$, by Definition 2.4, $|\partial_D^+(v)| = k$ and $|\partial_D^-(v)| = n - 1$. It follows by induction on n ,

$$\begin{aligned} |A(D)| &= |A(H)| + k + (n - 1) \leq k(2(n - 1) - k - 1) + \binom{(n - 1) - k}{2} \\ &\quad + k + (n - 1) \\ &= k(2n - k - 1) - 2k + \binom{n - k - 2}{i=1} i + 2k + (n - k - 1) \\ &= k(2n - k - 1) + \binom{n - k}{2}. \end{aligned}$$

Claim 2. If Corollary 2.6 (iii) holds, then (2) holds. Furthermore, if $k \geq 2$, then (2) holds with strict inequality.

Without loss of generality, we may assume that $D = [H_1, H_2]_k$ for some $H_1, H_2 \in \mathcal{D}(k)$. Let $n_1 = |V(H_1)|$ and $n_2 = |V(H_2)|$. It follows by induction on n ,

$$\begin{aligned} |A(D)| &= |A(H_1)| + k + n_1n_2 + |A(H_2)| \\ &\leq k(2n_1 - k - 1) + \binom{n_1 - k}{2} + k + n_1n_2 + k(2n_2 - k - 1) + \binom{n_2 - k}{2} \\ &= k(2n - k - 1) + \binom{n - k}{2} - k^2 + n_1n_2 + \binom{n_1 - k}{2} + \binom{n_2 - k}{2} \\ &\quad - \binom{n - k}{2}. \end{aligned}$$

We further observe the following.

$$\begin{aligned} &2 \left[n_1n_2 + \binom{n_1 - k}{2} + \binom{n_2 - k}{2} - \binom{n - k}{2} \right] \\ &= 2n_1n_2 + (n_1 - k)(n_1 - k - 1) + (n_2 - k)(n_2 - k - 1) - (n - k)(n - k - 1) \\ &= 2n_1n_2 + \left[\sum_{i=1}^2 (n_i^2 - 2n_i k + k^2 - n_i + k) \right] - (n^2 - 2nk + k^2 - n + k) \\ &= (n_1 + n_2)^2 - 2nk + 2k^2 - (n_1 + n_2) + 2k - n^2 + 2nk - k^2 + n - k = k(k + 1). \end{aligned}$$

As $k \geq 2$, we have $-k^2 + \frac{k(k+1)}{2} = -\frac{k^2-k}{2} < 0$, and so $|A(D)| \leq k(2n - k - 1) + \binom{n-k}{2} - \frac{k^2-k}{2}$. This implies Claim 2.

Claim 3. If $k \geq 2$, and if $D \in \mathcal{D}(k)$ satisfies equality in (2), then $D \in \mathcal{M}(k)$.

Let $D \in \mathcal{D}(k)$ be a digraph satisfying equality in (2). We argue by induction on $n = |V(D)| \geq k + 1$ to prove Claim 3. If $n = k + 1$, then $D = K_{k+1}^* \in \mathcal{M}(k)$. Hence we assume that $n > k + 1$ and that Claim 3 holds for smaller values of n .

If D satisfies Corollary 2.6(iii), then by Claim 2, D cannot satisfy equality in (2), contrary to the assumption of Claim 3. Hence D must satisfy Corollary 2.6(i) or (ii). By symmetry, we assume that for some digraph $H \in \mathcal{D}(k)$ and vertex $v \in V(D)$, we have $D = [\{v\}, H]_k$. By Definition 2.4, we have $|A(D)| = |A(H)| + k + n - 1$. As D satisfies equality in (2), this implies that

$$|A(H)| = k(2n - k - 1) + \binom{n - k}{2} - (n + k - 1) = k(2(n - 1) - k - 1) + \binom{(n - 1) - k}{2}.$$

It follows by induction that $H \in \mathcal{M}(k)$. By Definition 3.1, $D \in \mathcal{M}(k)$, and so Claim 3 is proved by induction.

Claim 4. If $k = 1$, then for any $D \in \mathcal{D}(k)$, we always have

$$|A(D)| = k(2n - k - 1) + \binom{n - k}{2} = 2(n - 1) + \frac{(n - 1)(n - 2)}{2}. \tag{4}$$

We again argue by induction. Since $|A(K_2^*)| = 2$, (4) holds when $n = k + 1 = 2$. Assume that $n > 2$ and (4) holds for smaller values of n . Hence D is not a complete digraph. By Corollary 2.6, we have the following observations.

Suppose that Corollary 2.6(i) or (ii) holds. Then $D \in \{[K_1, H]_k, [H, K_1]_k\}$ for some $H \in \mathcal{D}(1)$. By induction,

$$|A(D)| = 2[(n - 1) - 1] + 1 + (n - 1) + \frac{(n - 2)(n - 3)}{2} = 2(n - 1) + \frac{(n - 1)(n - 2)}{2}.$$

Thus (4) is justified in this case.

Now suppose that for some $H_1, H_2 \in \mathcal{D}(1)$, we have $D = [H_1, H_2]_k$. Let $n_1 = |V(H_1)|$ and $n_2 = |V(H_2)|$. By induction, and with the same computation in the proof of Claim 2, we have

$$\begin{aligned} |A(D)| &= |A(H_1)| + 1 + n_1n_2 + |A(H_2)| = \sum_{i=1}^2 \left[2(n_i - 1) + \frac{(n_i - 1)(n_i - 1)}{2} \right] \\ &+ 1 + n_1n_2 \\ &= 2(n_1 + n_2 - 1) - 1 + \frac{(n - 1)(n - 2)}{2} + \\ &\quad \frac{2n_1n_2 + (n_1 - 1)(n_1 - 2) + (n_2 - 1)(n_2 - 2) - (n - 1)(n - 2)}{2} \\ &= 2(n_1 + n_2 - 1) - 1 + \frac{(n - 1)(n - 2)}{2} + \frac{1 + 1}{2} = 2(n_1 + n_2 - 1) \\ &+ \frac{(n - 1)(n - 2)}{2}. \end{aligned}$$

Therefore, (4) is also justified in this case, and so Claim 4 is proved by induction.

Now Theorem 3.2 follows from (3), and Claims 1, 2, 3 and 4. ■

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