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ORIGINAL PAPER

On Spanning Disjoint Paths in Line Graphs

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Abstract Spanning connectivity of graphs has been intensively investigated in the study of interconnection networks (Hsu and Lin, Graph Theory and Interconnection Networks, 2009). For a graph *G* and an integer s > 0 and for $u, v \in V(G)$ with $u \neq v$, an (s; u, v)-path-system of *G* is a subgraph *H* consisting of *s* internally disjoint (u, v)-paths. A graph *G* is **spanning s-connected** if for any $u, v \in V(G)$ with $u \neq v$, *G* has a spanning (s; u, v)-path-system. The **spanning connectivity** $\kappa^*(G)$ of a graph *G* is the largest integer *s* such that *G* has a spanning (k; u, v)-path-system, for any $u, v \in V(G)$ with $u \neq v$, *G* has a spanning (s; u, v)-path-system. The **spanning connectivity** $\kappa^*(G)$ of a graph *G* is the largest integer *s* such that *G* has a spanning (k; u, v)-path-system, for any integer *k* with $1 \leq k \leq s$, and for any $u, v \in V(G)$ with $u \neq v$. An edge counter-part of $\kappa^*(G)$, defined as the supereulerian width of a graph *G*, has been investigated in Chen et al. (Supereulerian graphs with width *s* and *s*-collapsible graphs, 2012). In Catlin and Lai (Graph Theory, Combinatorics, and Applications, vol. 1, pp. 207–222, 1991) proved that if a graph *G* has 2 edge-disjoint spanning trees, and if L(G) is the line graph of

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E. Wei Department of Mathematics, Renming University of China, Beijing 100872, People's Republic of China *G*, then $\kappa^*(L(G)) \ge 2$ if and only if $\kappa(L(G)) \ge 3$. In this paper, we extend this result and prove that for any integer $k \ge 2$, if G_0 , the core of *G*, has *k* edge-disjoint spanning trees, then $\kappa^*(L(G)) \ge k$ if and only if $\kappa(L(G)) \ge \max\{3, k\}$.

Keywords Connectivity · Spanning connectivity · Hamiltonian linegraph · Hamiltonian-connected line graph · Supereulerian graphs · Collapsible graphs

1 Introduction

Graphs in this paper are finite and may have multiple edges but no loops. Terminology and notations not defined here are referred to [1]. In particular, for a graph $G, \delta(G), \kappa(G)$ and $\kappa'(G)$ represent the minimum degree, the connectivity and the edge connectivity of the graph G, respectively. A path with initial vertex u and terminal vertex v will be referred as a (u, v)-path. We use O(G) to denote the set of all odd degree vertices in G, and $D_i(G)$ the set of all vertices of degree i in G. A graph G is **Eulerian** if $O(G) = \emptyset$ and G is connected, and is **supereulerian** if G has a Eulerian subgraph H with V(H) = V(G). The maximum number of edge-disjoint spanning trees in a graph G is denoted by $\tau(G)$.

For an integer s > 0 and for $u, v \in V(G)$ with $u \neq v$, an (s; u, v)-path-system of G is a subgraph H consisting of s internally disjoint (u, v)-paths, and such an H is called a spanning (s, u, v)-path-system if V(H) = V(G). A graph G is **spanning s-connected** if for any $u, v \in V(G)$ with $u \neq v$, G has a spanning (s; u, v)-path-system. The **spanning connectivity** $\kappa^*(G)$ of a graph G is the largest integer s such that for any integer k with $1 \leq k \leq s$ and for any $u, v \in V(G)$ with $u \neq v$, G has a spanning (k; u, v)-path-system. A graph G is hamiltonian connected if for any $u, v \in V(G)$ with $u \neq v$, G has a spanning (k; u, v)-path-system. A graph G is hamiltonian connected if for any $u, v \in V(G)$ with $u \neq v$ G has a path P from u to v such that V(P) = V(G). Thus $\kappa^*(G) \geq 1$ if and only if G is hamiltonian-connected. The hamiltonian connectedness of graphs has been intensively studied, as shown in [8]. The spanning connectivity of a graph has also been studied, as can be seen in Chapters 14 and 15 of [11].

The **line graph** of a graph *G*, denoted by L(G), has E(G) as its vertex set, where two vertices in L(G) are adjacent if and only if the corresponding edges in *G* are adjacent in L(G). Many interesting structure properties of a graph are closely related to the same properties of its line graph. Cai and Corneil [2] proved that the cycle double conjecture [18,20] holds for all 2-edge-connected graphs if and only if it holds for all 2-edge-connected line graphs. Chen et al. [6] proved that to solve Tutte's flow conjectures [13,22] in graphs, one only needs to prove the truth of these conjectures in line graphs. Thomassen's conjecture [21] that "every 4-connected line graph is hamiltonian" had attracted many researchers working on properties of line graphs. Catlin and Lai in [5] characterized line graphs L(G) with $\kappa^*(L(G)) \ge 2$ for graphs *G* with $\tau(G) \ge 2$.

Theorem 1.1 (Catlin and Lai [5]) Let G be a graph with $\tau(G) \ge 2$. Then $\kappa^*(L(G)) \ge 2$ if and only if $\kappa(L(G)) \ge 3$.

By the well known spanning tree packing theorem of Nash-Williams [17] and Tutte [23], every 2k-edge-connected graph must have k edge-disjoint spanning trees. Therefore, Theorem 1.1 implies the next theorem.

Theorem 1.2 (Zhan [24]) If $\kappa'(G) \ge 4$, then $\kappa^*(L(G)) \ge 2$.

Huang and Hsu [12] proved the following theorem, which extends Theorem 1.2 from k = 2 to all integers $k \ge 2$.

Theorem 1.3 (Huang and Hsu [12]) For any integer $k \ge 2$, if $\kappa'(G) \ge 2k \ge 4$, then $\kappa^*(L(G)) \ge k$.

In this paper, using a modified Catlin's reduction technique [3], we prove a new theorem which includes the theorems mentioned above as special cases.

Let *G* be a graph such that $\kappa(L(G)) \ge 3$. The **core** of *G*, denoted by G_0 , is defined as follows (see [19]). For each $v \in V(G)$, let $E_G(v)$ denote the set of edges incident with v in *G*. For any $u, v \in D_2(G)$, $E_G(u) \cap E_G(v) = \emptyset$. For each $v \in D_2(G)$, denote $E_G(v) = \{e'_v, e''_v\}$. Let

$$X_2(G) = \{e_v'' : v \in D_2(G)\}.$$

We define the **core** of *G* by

$$G_0 = (G - D_1(G))/X_2(G).$$

Here is our main theorem:

Theorem 1.4 Let $k \ge 2$ be an integer, and G be a connected graph with a core G_0 such that $\tau(G_0) \ge k$. Then $\kappa^*(L(G)) \ge k$ if and only if $\kappa(L(G)) \ge \max\{3, k\}$.

Remark When k = 2, Theorem 1.4 implies Theorem 1.1. Noting that $\kappa'(G_0) \ge \kappa'(G)$ and applying the spanning tree packing theorem of Nash-Williams [17] and Tutte [23], Theorem 1.3 follows from Theorem 1.4 immediately. In [9], infinitely many graphs *G* satisfying $\kappa'(G) = \tau(G)$ with minimum possible edges have been constructed, and all such graphs are characterized. For any of such graph *G*, we have $\kappa^*(L(G)) \ge k$ by Theorem 1.4; but we cannot make the same conclusion by Theorem 1.3. Hence, Theorem 1.4 is stronger than Theorem 1.3. The following examples give additional evidences that even edge-connectivity condition in Theorem 1.3 can be relaxed.

Let $k \ge 2$ and *H* be any graph with *k* edge-disjoint spanning trees. Obtain *G* from *H* by

- (i) subdividing every edge of H exactly once, and
- (ii) attaching a pendent edge at every vertex of H (not including the new vertices resulting from the subdivision operation).

Since the core of G is H, by Theorem 1.4, $\kappa^*(L(G)) \ge k$. But since $\kappa'(G) < 2$, such a conclusion cannot be made by using Theorem 1.3.

In the next section, we will prove a characterization of a graph *G* whose line graph L(G) satisfying $\kappa^*(L(G)) \ge k$, analogous to the characterization of Harary and Nash-Williams on hamiltonian line graphs [10]. A reduction method involving *s*-collapsible graphs will be presented in Sect. 3. In Sect. 4, we review some properties of the core of a graph. The results in Sects. 2, 3, and 4 will be applied to prove the main result in Sect. 5.

2 Spanning Connectivity in Line Graphs

In this section, we shall follow the idea of Harary and Nash-Williams in [10] to determine a relationship between dominating (k; e', e'')-trail systems in G and spanning (k; e', e'')-path-systems in L(G). We view a trail of G as a vertex-edge alternating sequence

$$v_0, e_1, v_1, e_2, \dots, e_k, v_k$$
 (1)

such that all the e_i are distinct and for each $i = 1, 2, ..., ke_i$ is incident with both v_{i-1} and v_i . When the edge-vertex incidence is understood from the context for such a trail, we often use an edge sequence $e_1e_2...e_k$ to denote the same trail in (1). All the vertices in $\{v_1, v_2, ..., v_{k-1}\}$ are **internal vertices** of the trail in (1). For edges $e', e'' \in E(G)$, an (e', e'')-**trail** of G is a trail of G whose first edge is e' and whose last edge is e''. As an example, the trail in (1) is an (e_1, e_k) -trail. An (e', e'')-trail T of G is **dominating in** G if every edge of G is incident with an internal vertex of T; and a **spanning** (e', e'')-**trail** T of G is a dominating (e', e'')-trail T of G such that V(T) = V(G). A **dominating** (k; e', e'')-**trail systems** in G is a subgraph H consisting of k edge-disjoint (e', e'')-trail $(T_1, T_2, ..., T_k)$ such that every edge of G is incident with an internal vertex of T_i for some $i(1 \le i \le k)$.

Theorem 2.1 Let $s \ge 1$ be an integer, and G a graph with $|E(G)| \ge 3$. The following are equivalent.

- (i) $\kappa^*(L(G)) \ge s;$
- (ii) For any edge e', $e'' \in E(G)$, G has a dominating (k; e', e'')-trail-system, for all $1 \le k \le s$.

Proof Assume that $\kappa^*(L(G)) \ge s$. By the definition of κ^* , for any positive integer $k \le s$, and for any e' and e'' in E(G), L(G) has a spanning (k; e', e'')-path-system (P_1, P_2, \ldots, P_k) .

Denote $P_i = e_1^i e_2^i \dots e_{n_i}^i$, where each $e_j^i \in E(G) = V(L(G))$, and where $e_1^i = e'$ and $e_{n_i}^i = e''$, for $i = 1, 2, \dots, k$. By the definition of a line graph, *G* has a longest (e', e'')-trail $T_i = e_{i_1}^i e_{i_2}^i \dots e_{i_{n(i)}}$ such that $e_{i_1}^i = e', e_{i_{n(i)}} = e''$ and $i_1, i_2, \dots, i_{n(i)}$ is a subsequence of $1, 2, \dots, n_i$. Since $P_i = e_1^i e_2^j \dots e_{n_i}^i$ is a path in L(G), by the definition of a line graph and by the maximality of $|V(T_i)|$, for any j with $1 \le j < n(i)$, if $i_{j+1} > i_j + 1$ and if $v_j \in V(G)$ is the vertex in the trail $e_{i_1}^i e_{i_2}^j \dots e_{i_{n(i)}}^i$ incident with both e_{i_j} and $e_{i_{j+1}}$, then any edge e_i^i with $i_j < t < i_{j+1}$ must be incident with v_j in G. It follows that (T_1, T_2, \dots, T_k) is a dominating (k; e', e'')-trail-system of G. Conversely, we assume that (ii) holds to prove (i). Suppose $\{T_1, T_2, \ldots, T_k\}$ is a dominating (k; e', e'')-trail-system of G for any k with $1 \le k \le s$. By the definition of dominating (k; e', e'')-trail-systems, for any edge $e \in E(G) - \bigcup_{i=1}^{k} E(T_i)$, there exists an i such that e is incident with an internal vertex of T_i . Therefore, we can partition $E(G) - \bigcup_{i=1}^{k} E(T_i)$ into a disjoint union of subsets X_1, X_2, \ldots, X_k such that edges in X_i are incident with internal vertices of T_i . It follows by the definition of line graphs that in L(G), the vertex subset $E(T_i) \cup X_i$ induces a subgraph in L(G) which contains an (e', e'')-path P_i of L(G). Since every edge of G must be in an $E(T_i) \cup X_i$, (P_1, P_2, \ldots, P_k) is a spanning (e', e'')-path system of L(G).

3 Reductions and s-Collapsible Graphs

Throughout this paper, we shall adopt the convention that any graph G is 0-edge-connected, and always assume that $s \ge 1$ is an integer.

Definition 3.1 A graph *G* is *s*-collapsible if for any subset $R \subseteq V(G)$ with $|R| \equiv 0 \pmod{2}$, *G* has a spanning subgraph Γ_R such that

- (i) both $O(\Gamma_R) = R$ and $\kappa'(\Gamma_R) \ge s 1$, and
- (ii) $G E(\Gamma_R)$ is connected.

Thus a collapsible graph defined in [3] is a 1-collapsible graph in Definition 3.1. A spanning subgraph Γ_R of G with both properties in Definition 3.1 is an (s, R)-sub-graph of G. Let C_s denote the collection of all s-collapsible graphs. Then C_1 is the collection of all collapsible graphs [3]. By definition, for $s \ge 1$, any (s + 1, R)-sub-graph of G is also an (s, R)-subgraph of G. This implies that

$$C_{s+1} \subseteq C_s$$
, for any positive integer *s*. (2)

For a graph G, and for $X \subseteq E(G)$, the **contraction** G/X is obtained from G by identifying the two ends of each edge in X and then by deleting the resulting loops. If H is a subgraph of G, then we write G/H for G/E(H), and we use v_H to denote the vertex in G/H onto which H is contracted.

Proposition 3.2 ([7,15]) Let $s \ge 1$ be an integer. Then C_s satisfies the following.

(C1) $K_1 \in \mathcal{C}_s$

- (C2) If $G \in \mathcal{C}_s$ and if $e \in E(G)$, then $G/e \in \mathcal{C}_s$.
- (C3) If H is a subgraph of G and if $H, G/H \in C_s$, then $G \in C_s$.

Let *G* be a graph, and s > 0 be an integer. For any distinct $u, v \in V(G)$, an $(\mathbf{s}; \mathbf{u}, \mathbf{v})$ **trail-system** of *G* is a subgraph *H* consisting of *s* edge-disjoint (u, v)-trails. A graph is **supereulerian with width s** if for any $u, v \in V(G)$ with $u \neq v$, *G* has a spanning (s; u, v)-trail-system. The **supereulerian width** $\mu'(G)$ of a graph *G* is the largest integer *s* such that *G* is supereulerian with width *k* for any integer *k* with $1 \leq k \leq s$. A reduction method on applying *s*-collapsible graphs to study $\mu'(G)$ has been developed in [7,15]. **Lemma 3.3** ([7,15]) Let $s \ge 1$ be an integer. If a graph $G \in C_s$, then $\mu'(G) \ge s + 1$.

A graph is C_s -reduced if it contains no nontrivial subgraph in C_s . It is shown in [7] that every graph *G* has a unique collection of maximally *s*-collapsible subgraphs H_1, H_2, \ldots, H_c , and the graph $G'_s = G/(\bigcup_{i=1}^c E(H_i))$ is C_s -reduced, which is called the C_s -reduction of *G*.

Lemma 3.4 ([7,15]) Let $s \ge 1$ be an integer, G be a graph and H be a subgraph of G such that $H \in C_s$. Each of the following holds.

- (i) $G \in \mathcal{C}_s$ if and only if $G/H \in \mathcal{C}_s$.
- (ii) $\mu'(G) \ge s + 1$ if and only if $\mu'(G/H) \ge s + 1$.

Let F(G, s) denote the minimum number of additional edges that must be added to G to result in a graph Γ with $\tau(\Gamma) \ge s$. The quantity of F(G, s) has been determined in [16], whose matroidal versions are proved in [14,15].

Theorem 3.5 ([7,15]) Let $s \ge 1$ be an integer. If $F(G, s + 1) \le 1$, then $G \in C_s$ if and only if $\kappa'(G) \ge s + 1$.

Theorem 3.6 (Catlin et al. Theorem 1.3 of [4]) Let G be a connected graph and t an integer. If $F(G, 2) \le 2$, then $G \in C_1$ if and only if the C_1 -reduction of G is not a member in $\{K_2\} \cup \{K_{2,t} : t \ge 1\}$.

4 Facts on the Core of a Graph

Throughout this section, we assume that *G* is a connected graph satisfying $\kappa(L(G)) \ge 3$. For any $e', e'' \in E(G)$, let G(e', e'') be the graph obtained from *G* by replacing e' = u'v' by a path $u'v_{e'}v'$ and by replacing e'' = u''v'' by a path $u'v_{e''}v''$, where $v_{e'}$ and $v_{e''}$ are new vertices added to the graph when subdividing e' and e'', respectively.

Proposition 4.1 (Shao, Lemma 1.4.1 and Proposition 1.4.2 of [19]) Let *G* be a connected graph with $\kappa(L(G)) \ge 3$, and let G_0 denote the core of *G*. Each of the following holds.

- (i) G_0 is uniquely defined.
- (ii) $\delta(G_0) \ge \kappa'(G_0) \ge 3$.
- (iii) If G_0 is supereulerian, then L(G) is hamiltonian.
- (iv) If for any $e', e'' \in E(G_0), G_0(e', e'')$ has a spanning $(v_{e'}, v_{e''})$ -trail, then L(G) is hamiltonian-connected.

In this section, we extend some of Shao's results above for later applications in our proofs. For any integer k > 0, and for any $e', e'' \in E(G_0)$, define $G_0^k(e', e'')$ be the graph obtained from $G_0(e', e'')$ by, for any $v \in \{v_{e'}, v_{e''}\}$, replacing each edge incident with v in $G_0(e', e'')$ by a set of $\lceil k/2 \rceil$ parallel edges. As examples, $G_0^1(e', e'') = G_0^2(e', e'') = G_0(e', e'')$.

Lemma 4.2 Let k, l and s be integers such that $s \ge 1l \ge 2$ and $k \ge 2$.

- (i) ([7,15]) Let lK_2 is the loopless connected graph with two vertices and l edges. Then $lK_2 \in C_s$ if and only if $l \ge s + 1$. More generally, if T is a tree with $|E(T)| \ge 2$ and if lT is the graph obtained from T by replacing every edge of T by a set of l parallel edges. Then $lT \in C_s$ if and only if $l \ge s + 1$.
- (ii) If $G_0 \{e', e''\} \in \mathcal{C}_{k-1}$, then $G_0^k(e', e'') \in \mathcal{C}_{k-1}$.

Proof (ii). Let $G' = G_0 - \{e', e''\}$. By the definition of $G_0^k(e', e'')$, $G_0^k(e', e'')/G' = lK_{1,2}$ with $l \ge k$. By Lemma 4.2(i), $G_0^k(e', e'')/G' = lK_{1,2} \in C_{k-1}$. Since $G' \in C_{k-1}$, it follows by Proposition 3.2 (C3) that $G_0^k(e', e'') \in C_{k-1}$. □

Theorem 4.3 Let G be a graph with core G_0 , and let $k \ge 3$ be an integer. Each of the following holds.

- (i) If for any $e', e'' \in E(G_0)$ with $e' \neq e'', G_0(e', e'') \in C_1$, then $\kappa^*(L(G)) \ge 2$.
- (ii) If for any $e', e'' \in E(G_0)$ with $e' \neq e'', G_0^k(e', e'')$ has a spanning $(k; v_{e'}, v_{e''})$ -trail system, then G_0 has a spanning (k; e', e'')-trail system.
- (iii) If for any distinct edges e' = u'v' and e'' = u''v'' in $E(G_0)$, G_0 has a spanning (k; e', e'')-trail system $(T_1, T_2, ..., T_k)$ such that for any $v \in \{u', v', u'', v''\}$, there exists an i with $1 \le i \le k$, and such that T_i contains v as an internal vertex, then $\kappa^*(L(G)) \ge k$.
- (iv) If for any $e', e'' \in E(G_0)$ with $e' \neq e'', G_0 \{e', e''\} \in C_{k-1}$, then $\kappa^*(L(G)) \ge k$.
- *Proof* (i) Since $G_0(e', e'') \in C_1$, by Lemma 3.3, $G_0(e', e'')$ has a spanning $(v_{e'}, v_{e''})$ -trail. Thus by Proposition 4.1(iv), $\kappa^*(L(G)) \ge 2$.
 - (ii) Let H'' be a spanning $(k; v_{e'}, v_{e''})$ -trail system of $G_0^k(e', e'')$. Then H'' is an edge disjoint union of $(v_{e'}, v_{e''})$ -trails T'_1, T'_2, \ldots, T'_k . For each $i = 1, 2, \ldots, k$, let

$$T_i = G_0[E(T'_i - \{v_{e'}, v_{e''}\}) \cup \{e', e''\}].$$

Then each T_i is an (e', e'')-trail, and (T_1, T_2, \ldots, T_k) is a spanning (k; e', e'')-trail system of G_0 .

(iii) By Theorem 2.1, it suffices to show that for any $e', e'' \in E(G)$ with $e' \neq e''G$ has a dominating (k; e', e'')-trail system. By the assumption of (iii),

for any $e', e'' \in E(G_0)$ $(e' \neq e'')$, G_0 has a spanning (k; e', e'') – trail system with the property stated in *(iii)*. (3)

Let $e', e'' \in E(G)$ be two distinct edges. Let $e \in \{e', e''\}$. If $e \in E(G - D_1(G))$ and e is not incident with a vertex $z \in D_2(G)$, then let f(e) = e, which is an edge in $E(G_0)$. If $e \in E(G - D_1(G))$ and e is incident with a vertex $z \in D_2(G)$, then we may assume that $e \in E(G_0)$ and that the edge in $E_G(z) - \{e\}$ has been contracted in obtaining G_0 , and define f(e) = e, which is again an edge in $E(G_0)$. If e is incident with a vertex $z \in D_1(G)$, then denote e = zw, where $w \notin D_1(G)$. Define $f(e) \in E_G(w) - \{e\}$ so that $f(e') \neq f(e'')$. This can be done as $\kappa(L(G)) \ge 3$, when $z \in D_1(G)$, w must be incident with at least 4 edges in G. In any case, $f(e) \in E(G_0)$. Since f(e'), $f(e'') \in E(G_0)$ and $f(e') \ne f(e'')$, by (3), G_0 has a spanning (k; f(e'), f(e''))-trail system $(T'_1, T''_2, \ldots, T'_k)$ satisfying the assumption of (iii).

For each $i \in \{1, 2, ..., k\}$, let $X_2(T'_i)$ be the set of all edges $e \in E(T'_i)$ such that for some vertex $z \in D_2(G)$, $X_e := E_G(z) = \{e, f\}$. Define

$$T_i = G[(E(T'_i) - X_2(T'_i)) \bigcup \left(\bigcup_{e \in X_2(T'_i)} X_e\right) \bigcup \{e', e''\}].$$

In other words, T_i is obtained from T'_i by replacing each $e \in E(T'_i)$ that is incident with a vertex $z \in D_2(G)$ by the path consisting with both edges incident with z in G, and then extending the resulting trail to an (e', e'')-trail. It follows that (T_1, T_2, \ldots, T_k) is a (k; e', e'')-trail system that contains all vertices of degree at least 3 in G, such that for any $v \in \{u', v', u'', v''\}$, there exists an $i, (1 \le i \le k)$, such that T_i contains v as an internal vertex. Thus every edge not in $\bigcup_{i=1}^k E(T_i)$ must be incident with an internal vertex of some T_i , and so (T_1, T_2, \ldots, T_k) is a dominating (k; e', e'')-trail system.

(iv) By Theorem 4.3(iii), it suffices to show that the hypothesis of Theorem 4.3(iii) will hold.

Suppose that for any $e', e'' \in E(G_0)$ with $e' \neq e'', G_0 - \{e', e''\} \in C_{k-1}$. By Lemma 4.2, $G_0^k(e', e'') \in C_{k-1}$. It follows by Lemma 3.3 that $G_0^k(e', e'')$ has a $(k; v_{e'}, v_{e''})$ -trail system $(T_1', T_2', \ldots, T_k')$. Denote e' = u'v', e'' = u''v'' in $E(G_0)$. By the definition of $G_0^k(e', e'')$, there are at most $\lceil k/2 \rceil$ of these $(v_{e'}, v_{e''})$ -trails that contain the one of the $\lceil k/2 \rceil$ edges parallel to $v_{e'}v'$. This implies that for any $v \in \{u', v''\}$, at least one T_i' will use v as an internal vertex. Similarly, for any $v \in \{u'', v''\}$, at least one T_i' will use v as an internal vertex. Define

$$T_i = G_0[E(T'_i - \{v_{e'}, v_{e''}\}) \cup \{e', e''\}], \ (1 \le i \le k).$$

Then $(T_1, T_2, ..., T_k)$ is a spanning (k; e', e'')-trail system satisfying the hypothesis of Theorem 4.3(iii), and so $\kappa^*(L(G)) \ge k$. This completes the proof of the theorem.

5 Proof of Theorem 1.4

In this section, we shall prove the following slightly stronger result, which implies Theorem 1.4.

Proof By the Menger's theorem (Theorem 9.1 of [1]), for any graph G, we always have

$$\kappa(G) \ge \kappa^*(G). \tag{4}$$

By the definition of hamiltonian-connectivity, we know that every hamiltonian-connected graph *G* with at least 4 vertices must have connectivity at least 3. This, together with (4), implies that if $\kappa^*(L(G)) \ge k \ge 2$, then $\kappa(L(G)) \ge \max\{3, k\}$.

It remains to prove that for $k \ge 2$, if $\tau(G) \ge k$ and if $\kappa(L(G)) \ge \max\{3, k\}$, then $\kappa^*(L(G)) \ge k$.

First assume that k = 2. By Theorem 4.3, it suffices to show that for any pair of distinct edges $e', e'' \in E(G_0), G_0(e', e'') \in C_1$. We argue by contradiction and assume that $G_0(e', e'') \notin C_1$. Let G'_0 denote the C_1 -reduction of $G_0(e', e'')$. Since $\tau(G_0) \ge 2$, it follows that $F(G_0(e', e''), 2) \le 2$, and so $F(G'_0, 2) \le 2$. By Theorem 3.6, $G'_0 \in \{K_2, K_{2,t}, (t \ge 1)\}$. Since $G_0(e', e'')$ has no cut edges, neither does G'_0 . Hence $G'_0 = K_{2,t}$ for some $t \ge 2$. By Proposition 4.1, $\kappa'(G_0) \ge 3$, and so we must have t = 2, and $v_{e'}$ and $v_{e''}$, the two vertices newly added when subdividing e' and e'', are two nonadjacent vertices of G'_0 . By the definition of a core, the other two vertices in $V(G'_0) - \{v_{e'}, v_{e''}\}$ must be nontrivial vertices, and so $\{e', e''\}$ must be an essential edge cut of G, contrary to the assumption that $\kappa(L(G)) \ge 3$. This settles the case when k = 2.

Now we assume that $k \ge 3$. By Theorem 4.3, it suffices to show that the hypothesis of Theorem 4.3(ii) or (iii) holds. We shall assume that

$$G_0 - \{e', e''\} \notin \mathcal{C}_{k-1},\tag{5}$$

to prove Theorem 4.3(ii) holds.

Let $e', e'' \in E(G_0)$ be two distinct edges such that e' = u'v' and e'' = u''v''. Let G' denote the \mathcal{C}_{k-1} -reduction of $G_0 - \{e', e''\}$. Since $\tau(G_0) \ge k$, $F(G_0 - \{e', e''\}, 2) \le 2$.

Case 1 $F(G_0 - \{e', e''\}, k) \le 1.$

Since G' is a contraction of $G_0 - \{e', e''\}$, $F(G', k) \leq F(G_0 - \{e', e''\}, k) \leq 1$. By (5) and by Lemma 3.4, $G' \notin C_{k-1}$. By Theorem 3.5, $G' \in C_{k-1}$ if and only if $\kappa'(G') \geq k$. As $G' \notin C_{k-1}\kappa'(G') \leq k-1$. Since $F(G', k) \leq 1$, there must be an edge $f \notin E(G')$ such that $\tau(G' + f) \geq k$, and so $\kappa'(G' + f) \geq \tau(G' + f) \geq k$. This implies that G' must have an edge-cut X consisting of k - 1 edges $\{e_1, e_2, \ldots, e_{k-1}\}$, which is also an edge cut of $G_0 - \{e', e''\}$.

Since $\tau(G_0) \ge k$, G_0 has k spanning trees, denoted as T_1, T_2, \ldots, T_k . Since $F(G_0 - \{e', e''\}, k) \le 1$ and $G_0 - \{e', e''\}$ has an edge cut X with |X| = k - 1, we may assume that $e' \in E(T_k)$, and $e'' \notin \bigcup_{i=2}^{k-1} E(T_i)$. Hence we may assume that $e_i \in E(T_i)$, for $1 \le i \le k - 1$. Since T_k is a spanning tree of G_0 , we may assume that $T_k - e'$ has a (u', u'')-path P_k . Since $G_0[\bigcup_{i=1}^{k-1} E(T_i)]$ is a spanning subgraph of $G_0 - \{e, e''\}$ that has k - 1 edge-disjoint spanning trees, it follows by Theorem 3.5 and Lemma 3.3 with s = k - 2 that $G_0 - (\{e', e''\} \cup E(P_k))$ has a spanning (k - 1; u', u'')-trail system $(P_1, P_2, \ldots, P_{k-1})$. Let $P'_i = G_0[E(P_i) \cup \{e', e''\}], 1 \le i \le k$. Then $(P'_1, P'_2, \ldots, P'_k)$ is a spanning (k; e', e'')-trail system of G_0 . Hence Theorem 4.3(ii) holds.

Case 2 $F(G_0 - \{e', e''\}, k) = 2.$

Let T_1, T_2, \ldots, T_k denote k edge-disjoint spanning trees of G_0 . Since $F(G_0 - \{e', e''\}, 2) = 2$, we must have $e', e'' \in \bigcup_{i=1}^k E(T_i)$. Choose (T_1, T_2, \ldots, T_k) , among

all such choices of edge-disjoint spanning trees, such that

$$|\{T_i : E(T_i) \cap \{e', e''\} \neq \emptyset\}| \text{ is minimized.}$$
(6)

Subcase 2.1 For some $i, e', e'' \in E(T_i)$.

We may assume that $e', e'' \in E(T_k)$. Since T_k is a spanning tree of G_0 , we may assume (renaming the end vertices of e' and e'' if needed) that $T_k - \{e', e''\}$ has a (u', u'')-path P_k . Since $G_0[\bigcup_{i=1}^{k-1} E(T_i)]$ is a spanning subgraph of $G_0 - \{e, e''\}$ that has k - 1 edge-disjoint spanning trees, by Theorem 3.5 and Lemma 3.3 with s = k - 2that $G_0[\bigcup_{i=1}^{k-1} E(T_i)]$ has a spanning (k - 1; u', u'')-trail system $(P_1, P_2, \ldots, P_{k-1})$. Let $P'_i = G_0[E(P_i) \cup \{e', e''\}] 1 \le i \le k$. Then $(P'_1, P'_2, \ldots, P'_k)$ is a spanning (k; e', e'')-trail system of G_0 . Hence Theorem 4.3(ii) holds.

Subcase 2.2 For any $i, |\{e', e''\} \cap E(T_i)| \le 1$.

We may assume that $e' \in E(T_{k-1})$ and $e'' \in E(T_k)$. Let T'_{k-1} and T''_{k-1} be the two components of $T_{k-1} - e'$, and let T'_k and T''_k be the two components of $T_k - e''$. We may further assume that $V(T'_{k-1}) \cap V(T''_k) \neq \emptyset$.

We shall show first that $V(T'_{k-1}) = V(T'_k)$ and so $V(T''_{k-1}) = V(T''_k)$. Since T'_{k-1} and T''_{k-1} are the two components of $T_{k-1} - e'$, we may assume that $u' \in V(T'_{k-1})$ and $v' \in V(T''_{k-1})$. Similarly, we may assume that $u'' \in V(T'_k)$ and $v'' \in V(T''_{k-1})$ since T_{k-1} and T_k are spanning trees of G_0 , if $V(T'_{k-1}) \neq V(T'_k)$, then either $v'' \in V(T''_{k-1})$ or $u'' \in V(T''_{k-1})$. Arguing by contradiction, we assume that $V(T'_{k-1}) \neq V(T'_k)$, and by symmetry, we assume further that $v'' \in V(T'_{k-1})$. It then follows that $T'_{k-1} + e''$ has a unique cycle C'' which contains at least one edge $e''' \in E(T'_{k-1})$. Redefine $L_{k-1} = T_{k-1} + e'' - e'''$ and $L_k = T_k - e'' + e'''$. Then $(T_1, \ldots, T_{k-2}, L_{k-1}, L_k)$ is a set of k edge-disjoint spanning trees violating (6). Hence we must have

$$V(T'_{k-1}) = V(T'_k) \text{ and } V(T''_{k-1}) = V(T''_k).$$
(7)

By (7), both $(T_{k-1} \cup T_k)[V(T'_{k-1})]$ and $(T_{k-1} \cup T_k)[V(T''_{k-1})]$ are graphs with 2 edgedisjoint spanning trees. By Theorem 3.5 with s = 1, both are in C_1 , and so by Lemma 3.3, $(T_{k-1} \cup T_k)[V(T'_{k-1})]$ has a spanning (u', u'')-trail P_{k-1} and $(T_{k-1} \cup T_k)[V(T''_{k-1})]$ has a spanning (v', v'')-trail P_k . Since $T_1, T_2, \ldots, T_{k-2}$ are spanning trees of G_0 , each T_i has a (u', u'')-path $P_i(1 \le i \le k-2)$. Let $P'_i = G_0[E(P_i) \cup \{e', e''\}], 1 \le i \le k$. Then $(P'_1, P'_2, \ldots, P'_k)$ is a spanning (k; e', e'')-trail system of G_0 . Hence Theorem 4.3(ii) holds.

Since in all the cases, either Theorem 4.3(ii) or Theorem 4.3(iii) must hold, and so by Theorem 4.3, $\kappa^*(L(G)) \ge k$. This completes the proof.

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