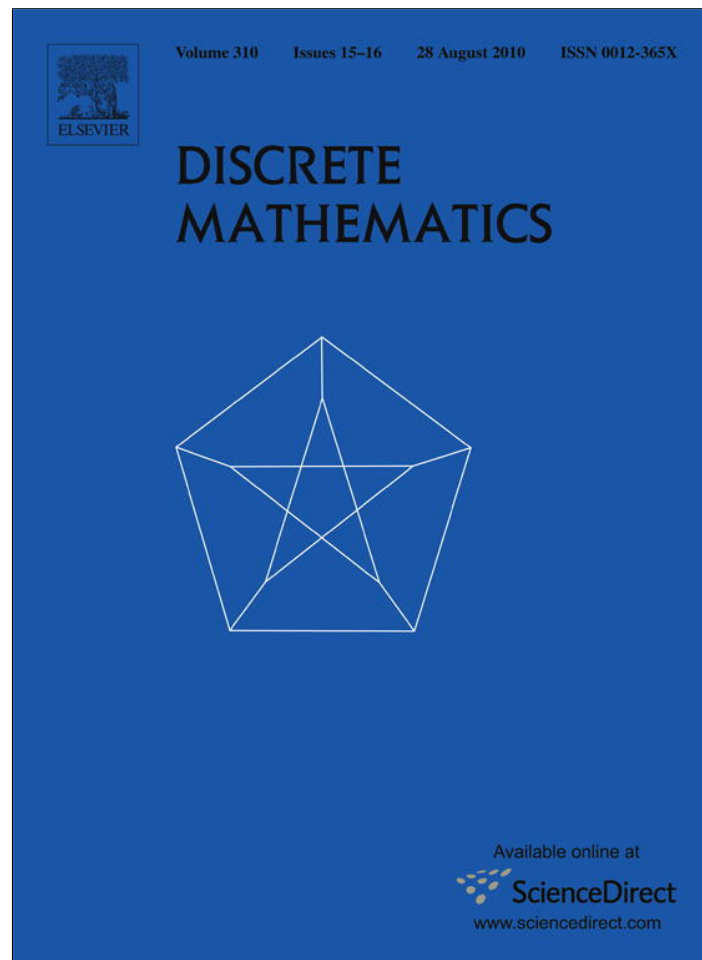


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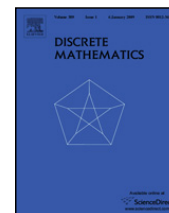
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The Chvátal–Erdős condition for supereulerian graphs and the Hamiltonian index

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ABSTRACT

A classical result of Chvátal and Erdős says that the graph G with connectivity $\kappa(G)$ not less than its independence number $\alpha(G)$ (i.e. $\kappa(G) \geq \alpha(G)$) is Hamiltonian. In this paper, we show that the 2-connected graph G with $\kappa(G) \geq \alpha(G) - 1$ is one of the following: supereulerian, the Petersen graph, the 2-connected graph with three vertices of degree two obtained from $K_{2,3}$ by replacing a vertex of degree three with a triangle, one of the 2-connected graphs obtained from $K_{2,3}$ by replacing a vertex of degree two with a complete graph of order at least three and by replacing at most one branch of length two in the resulting graph with a branch of length three, or one of the graphs obtained from $K_{2,3}$ by replacing at most two branches of $K_{2,3}$ with a branch of length three. We also show that the Hamiltonian index of the simple 2-connected graph G with $\kappa(G) \geq \alpha(G) - t$ is at most $\lfloor \frac{2t+2}{3} \rfloor$ for every nonnegative integer t . The upper bound is sharp.

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1. Introduction

We denote by $\kappa(G)$ and $\alpha(G)$ the connectivity and the independence number of G respectively, considering only the simple graph G with $\kappa(G) \geq 2$ in the following. Chvátal and Erdős gave the following well-known sufficient condition for a graph to be Hamiltonian.

Theorem 1 ([7]). *If $\kappa(G) \geq \alpha(G)$, then G is Hamiltonian.*

There are many extensions of **Theorem 1**, one of which aims at showing what properties a graph has when it satisfies a stronger Chvátal–Erdős condition.

Theorem 2 ([9]). *Let t be a nonnegative integer. If $\kappa(G) \geq \alpha(G) + t$, then any system of disjoint paths of total length at most t is contained in a Hamiltonian cycle of G .*

An immediate consequence of **Theorem 2** is that G is Hamilton connected if $\kappa(G) \geq \alpha(G) + 1$. A graph G is called *trivial* if it has only one vertex, and is called *even* if each vertex of G has even degree. A graph is called *eulerian* if it is connected and even.

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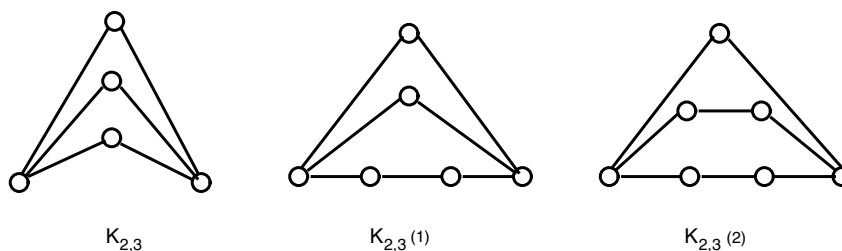


Fig. 1. Three graphs $K_{2,3}$, $K_{2,3}(1)$, $K_{2,3}(2)$.

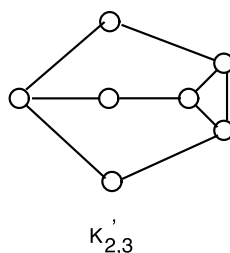


Fig. 2. The graph $K'_{2,3}$.

A graph is called *supereulerian* if it contains a spanning eulerian subgraph. Obviously, a Hamiltonian graph is supereulerian, whereas the converse is not generally true. The first aim of this paper is to consider whether a graph with a slightly weaker Chvátal–Erdős condition is supereulerian.

Theorem 3. *Let G be a 2-connected simple graph. If $\kappa(G) \geq \alpha(G) - 1$, then exactly one of the following holds.*

- (a) G is supereulerian;
- (b) $G \in \{\text{the Petersen graph}, K_{2,3}, K_{2,3}(1), K_{2,3}(2), K'_{2,3}\}$; see Figs. 1 and 2;
- (c) G is one of the two 2-connected graphs obtained from $K_{2,3}$ and $K_{2,3}(1)$ by replacing a vertex whose neighbors have degree three in $K_{2,3}$ and $K_{2,3}(1)$ with a complete graph of order at least three.

The line graph $L(G)$ of $G = (V(G), E(G))$ has $E(G)$ as its vertex set, where two vertices are adjacent in $L(G)$ if and only if the corresponding edges are incident in G . The m -iterated line graph $L^m(G)$ is defined recursively by $L^0(G) = G$, $L^1(G) = L(G)$, and $L^m(G) = L(L^{m-1}(G))$. The *Hamiltonian index* of a graph G , denoted by $h(G)$, is the smallest integer m such that $L^m(G)$ is Hamiltonian. In Section 2, we will show that the following known result is a consequence of Theorem 3.

Theorem 4 ([2]). *If $\kappa(G) \geq \alpha(G) - 1$, then $L(G)$ is Hamiltonian, i.e., $h(G) \leq 1$.*

The following result extends Theorem 1 for triangle-free graphs, showing that a graph with a weaker Chvátal–Erdős condition has a weaker result.

Theorem 5 ([8]). *If G is a 2-connected triangle-free graph of order n with $\kappa(G) \geq \frac{\alpha(G)+2}{2}$, then every longest cycle C in G is dominating (i.e., every edge of G has at least one end vertex in C), and G has a cycle of length at least $\min\{n - \alpha(G) + \kappa(G), n\}$.*

There are many other ways to extend Theorem 1 (see the survey paper [11] and update papers [1,13,14]). Our second aim in this paper is to extend Theorems 1 and 4 by presenting an upper bound for $h(G)$.

Theorem 6. *Let G be a 2-connected simple graph and let t be a nonnegative integer. If $\kappa(G) \geq \alpha(G) - t$, then $h(G) \leq \lfloor \frac{2t+2}{3} \rfloor$.*

In Section 2, we will provide some auxiliary results that are applied in Sections 3 and 4 to prove our main results. The sharpness of Theorem 6 is presented in the last section.

2. Preliminaries

2.1. Catlin's reduction method for supereulerian graphs

For a simple graph G , let $O(G)$ denote the set of odd degree vertices in G . In [3] Catlin defined collapsible graphs. Given a subset $R \subseteq V(G)$, with $|R|$ being even, a subgraph Γ of G is an R -subgraph if $O(\Gamma) = R$ and $G - E(\Gamma)$ is connected. A graph G is *collapsible* if, for any even subset R of $V(G)$, G has an R -subgraph. Catlin showed in [3] that every vertex of G lies in a unique maximal collapsible subgraph of G . If the graph G itself is collapsible, then we let $R = O(G)$. Obviously, $|R|$ is even. So, there exists an R -subgraph Γ of G such that $O(\Gamma) = R$ and $G - E(\Gamma)$ is connected. Hence, $G - E(\Gamma)$ must be a spanning eulerian subgraph. Note that C_4 is supereulerian but not collapsible. Thus every collapsible graph is supereulerian whereas

the converse is not generally true. If H is a connected subgraph of a graph G , then G/H denotes the graph by contracting H obtained from G by replacing H by a vertex v_H such that the number of edges in G/H joining any $v \in V(G) - V(H)$ to v_H in G/H equals the number of edges joining v in G to H . In [3], Catlin showed that any graph G has a unique collection of pairwise vertex-disjoint maximal collapsible subgraphs H_1, H_2, \dots, H_c such that $\bigcup_{i=1}^c V(H_i) = V(G)$. The *reduction* of G , denoted by G' , is the graph obtained from G by contracting each H_i ($1 \leq i \leq c$) to a single vertex v_{H_i} . We call H_i the pre-image of v_{H_i} , denoted by $R_{G'}^{-1}(v_{H_i})$. If H_i contains only one vertex, then v_{H_i} is called trivial. If the reduction of G is itself, then G is called a *reduced graph*.

Theorem 7 ([3]). *Let G be a connected graph and G' the reduction of G . Then, each of the following holds.*

- (a) G is *supereulerian* if and only if G' is *supereulerian*.
- (b) G' is *triangle-free* with $\delta(G') \leq 3$.

Theorem 8 ([12]). *If G is a 4-edge-connected graph, then G is *supereulerian*.*

The following theorem is obvious, where $\kappa'(G)$ denotes the edge-connectivity of G .

Theorem 9. *Let G' be the reduction of G . Then, each of the following holds.*

- (a) $\alpha(G') \leq \alpha(G)$.
- (b) If G' is nontrivial, then $\kappa'(G') \geq \kappa'(G)$.

Theorem 10 ([5]). *If G is a reduced graph and $\alpha(G) \geq 4$, then we have $(\delta(G)\alpha(G) + 4)/2 \leq |V(G)| \leq 4\alpha(G) - 5$.*

Theorem 11 ([6]). *Let G be the graph of order $n \leq 11$. If $\kappa'(G) \geq 3$, then G is either collapsible or the Petersen graph.*

2.2. Hamiltonian iterated line graphs

Let G be a graph. For any two subgraphs H_1 and H_2 of G , the *distance* $d_G(H_1, H_2)$ between H_1 and H_2 is defined to be the minimum of the distances $d_G(v_1, v_2)$ over all pairs with $v_1 \in V(H_1)$ and $v_2 \in V(H_2)$. If $d_G(e, H) = 0$ for an edge e of G , then we say that H *dominates* e . A subgraph H of G is called *dominating* if it dominates all edges of G . There is a characterization of a graph G with $h(G) \leq 1$ that involves the existence of a dominating eulerian subgraph in G .

Theorem 12 ([10]). *Let G be a graph with at least three edges. Then $h(G) \leq 1$ if and only if G has a dominating eulerian subgraph.*

By Theorem 12, all supereulerian graphs and the Petersen graph have Hamiltonian line graphs, and those graphs obtained from $K_{2,3}$ described in Theorem 3 have Hamiltonian line graphs, which shows that Theorem 4 is a consequence of Theorem 3.

For each integer $i \geq 0$, define $V_i(G) = \{v \in V(G) : d_G(v) = i\}$. A *branch* in G is a nontrivial path in which each inner vertex lies in $V_2(G)$ and neither end vertex lies in $V_2(G)$. We denote by $\mathcal{B}(G)$ the set of branches of G and by $\mathcal{B}_1(G)$ the subset of $\mathcal{B}(G)$ in which every branch has an end vertex in $V_1(G)$. For any subgraph H of G , we denote by $\mathcal{B}_H(G)$ the set of branches of G whose edges are all in H .

The following theorem can be considered as an analogue of Theorem 12 for $L^k(G)$.

Theorem 13 ([16]). *Let G be a connected graph that is not a path, and let $k \geq 2$ be an integer. Then, $h(G) \leq k$ if and only if $EU_k(G) \neq \emptyset$ where $EU_k(G)$ denotes the set of those subgraphs H of G which satisfy the following five conditions.*

- (I) H is an even subgraph.
- (II) $V_0(H) \subseteq \bigcup_{i=3}^{\Delta(G)} V_i(G) \subseteq V(H)$.
- (III) $d_G(H_1, H - H_1) \leq k - 1$ for every induced subgraph H_1 of H with $\emptyset \neq V(H_1) \subsetneq V(H)$.
- (IV) $|E(B)| \leq k + 1$ for every branch $B \in \mathcal{B}(G) \setminus \mathcal{B}_H(G)$.
- (V) $|E(B)| \leq k$ for every branch B in $\mathcal{B}_1(G)$.

Note that, if we only consider 2-connected graphs, then Condition (V) in the definition of $EU_k(G)$ is superfluous.

3. Proof of Theorem 3

Before presenting our proof, we start by proving three useful lemmas.

Lemma 14. *Let G' be a reduced graph such that $\kappa'(G') \geq 2$ and $\alpha(G') \leq 3$, then $\kappa(G') \geq 2$.*

Proof of Lemma 14. Assume, in contrast, that v is a cut vertex of G' , then $d_{G'}(v) \geq 4$ by $\kappa'(G') \geq 2$. Since G' is triangle-free, v has at least four neighbors in G' which form an independent set of G' , a contradiction. \square

Lemma 15. Let G' be a reduced graph such that $\kappa(G') = 2$ and $\alpha(G') = 3$. Then, exactly one of the following holds.

- (a) G' is supereulerian.
- (b) G' is isomorphic to one of the three graphs depicted in Fig. 1.

Proof of Lemma 15. Notice that a loop graph with one vertex is collapsible, and a two-vertex graph with parallel edges is also collapsible. Since G' is reduced, G' is a simple graph.

Let $k = \kappa(G')$, and $C = u_1 u_2 \cdots u_1$ be a closed trail with a maximal number of vertices in C (i.e., $|V(G') \setminus V(C)|$ is minimal). Then, C is a cycle (otherwise there is a vertex v of C with $d_C(v) \geq 4$ and hence the neighbors of v consist of an independent set of at least four vertices of G' by Theorem 7, contradicting that $\alpha(G') = 3$). We denote by u^+ (u^- , respectively) the successor (the predecessor, respectively) of u on C , and let $u^{++} = (u^+)^+$, $u^{--} = (u^-)^-$ and so on; a portion of C from u to v means $uu^+u^{++} \dots v^-v$. Since G' is triangle-free, we obtain that $u^-u^+ \notin E(G)$ for any vertex u in C . Since $\delta(G') \geq \kappa(G')$, and every graph with $\delta(G') \geq 2$ has a cycle of length at least $\delta(G') + 1$, C has at least $k + 1$ vertices.

If C is not a spanning cycle of G' , then there is a component H of $G' - V(C)$. We will prove that G' is isomorphic to one of those graphs depicted in Fig. 1.

Since $\kappa(G') = 2$, C has at least two vertices having neighbors in $V(H)$. Let x, y be any two vertices of C having neighbors in $V(H)$.

We claim that none of $\{x^-, x^+, y^-, y^+\}$ has a neighbor in H ; otherwise, say, x^+ has a neighbor in H , then we can obtain a cycle from C by adding an x, x^+ -path through H and by deleting the edge xx^+ , which contains more vertices than C , contradicting the maximality of $|V(C)|$. We have the following two facts.

Claim 1. $x^+y^+ \notin E(G')$ and $x^-y^- \notin E(G')$.

Proof of Claim 1. By symmetry, we only need to prove that $x^+y^+ \notin E(G')$. Since C is a cycle and $x \neq y^+$, $x^+y^+ \notin E(C)$. Suppose $x^+y^+ \in E(G') \setminus E(C)$, then we can obtain a cycle that contains more vertices than C by including x^+y^+ , the portions of C from x^+ to y and from y^+ to x , and an x, y -path through H , contradicting the maximality of $|V(C)|$. \square

Claim 2. Either $x^+ = y^-$ or $x^+y^- \in E(C)$ and either $x^- = y^+$ or $x^-y^+ \in E(C)$.

Proof of Claim 2. By symmetry, we only need to prove the second part of the claim. We will prove this by contradiction. If possible, suppose that neither $x^- = y^+$ nor $x^-y^+ \in E(C)$, then there are at least five vertices in the portion of C from y to x . Suppose first that $x^-y^+ \notin E(G')$; then, by Claim 1, $S = \{x^+, x^-, y^+\}$ is an independent set such that none of S is adjacent to H ; hence we can get an independent set of four vertices by adding one vertex of H to S , contradicting $\alpha(G') = 3$. Now, suppose that $x^-y^+ \in E(G')$; then, by the fact that G' is triangle-free and by the assumption that $x^-y^+ \notin E(C)$, $x^{--}y^+ \notin E(G')$ and there are at least six vertices in the portion of C from y to x . Note that x^{--} is not adjacent to any vertex of H ; otherwise, we can obtain a cycle that contains more vertices than C by including x^-y^+ , the portions of C from y^+ to x^{--} and from x^- to y , and an x^-, y -path through H , contradicting $\alpha(G') = 3$. Hence, $x^{--}x^+ \in E(G')$; otherwise $\{x^+, x^{--}, y^+, h\}$ is an independent set of four vertices, where $h \in V(H)$, contradicting $\alpha(G') = 3$. Note that $x^{--}x^+ \notin E(C)$. Hence the cycle obtained from C by deleting three edges $xx^+, x^{--}x^+, yy^+$, and by adding two edges x^+x^{--}, y^+x^- and an x, y -path through H , contains more vertices than C , contradicting the maximality of $|V(C)|$. This proves Claim 2. \square

By Claim 2, C is a cycle of length at most six and any vertex in $V(C) \setminus \{x, y\}$ is not adjacent to H . This implies that $\{x, y\}$ is a 2-vertex cut-set of G' . By the choice of C , H has at most two vertices; otherwise, suppose H has at least three vertices, then x and y must have different neighbors h_x and h_y in H , respectively, since $\kappa(G') = 2$. No matter whether h_x and h_y are adjacent or not, there must be a path $P(h_x, h_y)$ of H between h_x and h_y with length at least two, since $\kappa(G') = 2$. Hence, we will obtain a longer cycle than C by including the path $P(h_x, h_y)$, xh_x, yh_y and a longest portion of C between x and y , contradicting the maximality of $|V(C)|$. So $|V(H)| \leq 2$.

Claim 3. $|V(H)| = 1$ and $G' - V(C)$ does not have more than one component.

Proof of Claim 3. Noticing that $4 \leq |V(C)| \leq 6$, we only need to discuss the following two cases.

Case I: $|V(C)| = 6$.

By Claim 2, we may assume $C = xx^+y^-yy^+x^-x$. Then $|V(H)| = 1$; otherwise, since $\kappa(G') = 2$, $|V(H)| = 2$ and G' is triangle-free, we may assume that $V(H) = \{h_x, h_y\}$ with $xh_x, yh_y \in E(G')$ and $xh_y \notin E(G')$, and hence $\{x, y^-, y^+, h_y\}$ is an independent set of G' , contradicting $\alpha(G') = 3$. Hence, x and y share the same neighbor v_H in H . It follows that $xy \notin E(G')$ since G' is triangle-free. Therefore C is an induced cycle of G' by the fact that G' is triangle-free and by Claims 1 and 2.

If $G' - V(C)$ has one more component K , then $|V(K)| = 1$ by the same argument on H . Moreover, the only vertex v_K in K is only adjacent to two vertices of C and the two vertices must be $\{x, y\}$ or $\{x^-, y^-\}$ or $\{x^+, y^+\}$ by the maximality of $|V(C)|$. If the two vertices are $\{x, y\}$, then $\{v_H, v_K, x^-, x^+\}$ is an independent set of order 4, a contradiction. If the two vertices are $\{x^-, y^-\}$, then $\{v_H, v_K, x^+, y^+\}$ is an independent set of order 4, a contradiction. If the two vertices are $\{x^+, y^+\}$, then $\{v_H, v_K, x^-, y^-\}$

is an independent set of order 4, a contradiction. Hence, $G' - V(C)$ does not have more than one component in this situation and G' is the graph $K_{2,3}(2)$ depicted in Fig. 1.

Case II: $|V(C)| = 4$ or 5.

We claim $|V(H)| = 1$ (otherwise, since $\kappa(G') = 2$ and $|V(H)| = 2$, we may assume that $V(H) = \{h_x, h_y\}$ with $xh_x, yh_y, h_xh_y \in E(G')$ and obtain a longer cycle than C by including the path xh_xh_y and a longest portion of C between x and y , contradicting the maximality of $|V(C)|$). Since G' is triangle-free, $xy \notin E(G')$. It follows that C is an induced cycle of G' , by Claims 1 and 2. Note that the vertex v_H of H is not adjacent to any vertex of $V(C) \setminus \{x, y\}$.

If $G' - V(C)$ has one more component K , then $|V(K)| = 1$ by the same argument on H . Moreover, the only vertex v_K in K is only adjacent to two nonadjacent vertices of C .

If $|V(C)| = 5$, then v_K is only adjacent to the two vertices $\{x, y\}$ or $\{x, y^-\}$ or $\{x^-, y^-\}$ or $\{x^-, x^+\}$ or $\{x^+, y\}$. Notice that $x^- = y^+$. By symmetry, we only need to discuss the following three cases: if the two vertices are $\{x, y\}$, then $\{v_H, v_K, x^-, y^-\}$ is an independent set of order 4, a contradiction; if the two vertices are $\{x, y^-\}$, then $\{v_H, v_K, x^-, x^+\}$ is an independent set of order 4, a contradiction; if the two vertices are $\{x^-, x^+\}$, then $v_Hx^-v_Kx^+y^-yv_H$ is a longer cycle than C , a contradiction. Hence $G' - V(C)$ does not have more than one component in this situation; and then G' is the graph $K_{2,3}(1)$ depicted in Fig. 1.

If $|V(C)| = 4$, then v_K is only adjacent to the two vertices $\{x, y\}$ or $\{x^-, y^-\}$. If the two vertices are x, y , then $\{v_H, v_K, x^-, y^-\}$ is an independent set of order 4, a contradiction; if the two vertices are $\{x^-, y^-\}$, then $v_Hx^-v_Ky^-yv_H$ is a longer cycle than C , a contradiction. Hence $G' - V(C)$ does not have more than one component in this situation, and then $G' = K_{2,3}$. So we finish the proof of Claim 3. \square

Above all, G' is isomorphic to one of those graphs depicted in Fig. 1. This completes the proof of Lemma 15. \square

Lemma 16. Let G' be the reduction of a simple graph G . Then, the following statements hold.

- (a) If v is a nontrivial vertex of G' with $d_{G'}(v) = 2$, then there exists at least one vertex v' of $R_{G'}^-(v)$ such that v' is not adjacent to any vertex of $\bigcup_{s \in V(G') \setminus \{v\}} R_{G'}^-(s)$.
- (b) if u is a nontrivial vertex of G' with $d_{G'}(u) = 3$, then for any pair of vertices $\{x, y\} \subseteq N_{G'}(u)$, there exists at least one vertex u' of $R_{G'}^-(u)$ such that u' is not adjacent to any vertex of $R_{G'}^-(x) \cup R_{G'}^-(y)$.

Proof of Lemma 16. Note that, by the definition of contraction, $R_{G'}^-(s)$ is collapsible for any nontrivial vertex s of $V(G')$; hence, since G is simple, $R_{G'}^-(s)$ has at least three vertices. Therefore, since $d_{G'}(v) = 2$, there are at most two vertices v'', v''' of $R_{G'}^-(v)$ that are adjacent to one or two vertices of $\bigcup_{s \in N_{G'} \setminus \{v\}} R_{G'}^-(s)$ by the definition of contraction. Hence, there exists at least one vertex $v' \in R_{G'}^-(v) \setminus \{v'', v'''\}$ that is not adjacent to any vertex of $\bigcup_{s \in V(G') \setminus \{v\}} R_{G'}^-(s)$. This proves the first statement. For a nontrivial vertex u of G' with $d_{G'}(u) = 3$, there is exactly one vertex of $R_{G'}^-(u)$ that is adjacent to some vertex of $R_{G'}^-(s)$ for any $s \in N_{G'}(u)$ by the definition of contraction. Therefore, since $|V(R_{G'}^-(s))| \geq 3$ for any nontrivial vertex s of G' , the second statement holds. This completes the proof of Lemma 16. \square

Now we finish the proof of Theorem 3.

Proof of Theorem 3. Let G' be the reduction of G . If $G' = K_1$, then G is supereulerian by Theorem 7. So, we only need to consider the case when G' is not K_1 .

If $\kappa(G) \geq 4$, then G is supereulerian by Theorem 8.

If $\kappa(G) = 3$, then, by hypothesis, we have that $\alpha(G) \leq 4$. We let $\alpha(G) = 4$ since, otherwise, G is Hamiltonian by Theorem 1. Theorem 9 gives that $\alpha(G') \leq 4$ and $\kappa'(G') \geq \kappa'(G) \geq \kappa(G) \geq 3$. First, suppose that $\alpha(G') = 4$, then from Theorem 10 we deduce that $|V(G')| \leq 11$. Hence, G' is either collapsible or the Petersen graph by Theorem 11. If G' is collapsible, then G' is supereulerian and G is hence supereulerian by Theorem 7. If G' is the Petersen graph, then we claim that every vertex of G' is trivial. Assume, in contrast, that there exists a nontrivial vertex w of G' . Let $N_{G'}(w) = \{x, y, z\}$. By Lemma 16(b), there exists at least one vertex $w' \in R_{G'}^-(w)$ such that w' is not adjacent to any vertex of $\bigcup_{v \in \{x, y\}} R_{G'}^-(v)$. Let $N_{G'}(z) \setminus \{w\} = \{u, v\}$. Taking $x' \in R_{G'}^-(x), y' \in R_{G'}^-(y), u' \in R_{G'}^-(u)$ and $v' \in R_{G'}^-(v)$, we obtain an independent set $\{w', x', y', u', v'\}$ of G , contradicting that $\alpha(G) = 4$. This shows that G' has no nontrivial vertex. Hence, $G = G'$, i.e., G is the Petersen graph. Now, suppose $\alpha(G') \leq 3$. Since $\kappa'(G') \geq 3, \kappa(G') \geq 2$ by Lemma 14. In the case when $\kappa(G') \geq \alpha(G')$, we have that G' is Hamiltonian (hence supereulerian) by Theorem 1. So G is supereulerian by Theorem 7. It remains the case when $\kappa(G') = 2$ and $\alpha(G') = 3$. By Lemma 15, G' is isomorphic to one of the three graphs in Fig. 1, which contradicts that $\kappa'(G') \geq 3$.

It remains to consider the case that $\kappa(G) = 2$ and $\alpha(G) = 3$. Theorem 9 gives that $\alpha(G') \leq 3$ and $\kappa'(G') \geq \kappa'(G) \geq \kappa(G) = 2$. Hence, $\kappa(G') \geq 2$ by Lemma 14. We only need to consider the case that $\kappa(G') = 2$ and $\alpha(G') = 3$ since, in all other cases, G' is Hamiltonian and G is hence supereulerian. By Lemma 15, G' is either supereulerian or isomorphic to one of $\{K_{2,3}, K_{2,3}(1), K_{2,3}(2)\}$. If G' is supereulerian, then so is G by Theorem 7.

Now, suppose that G' is isomorphic to one of $\{K_{2,3}, K_{2,3}(1), K_{2,3}(2)\}$. Then, we claim that G' has at most one nontrivial vertex. Assume, in contrast, that G' has two nontrivial vertices u and v . Note that G is a simple graph. Firstly, suppose $d_{G'}(u) = d_{G'}(v) = 2$. Then, by Lemma 16(a), there exist at least two vertices u' and v' in $R_{G'}^-(u)$ and $R_{G'}^-(v)$, respectively,

such that they are not adjacent to any vertex of $\bigcup_{x \in V(G') \setminus \{u, v\}} R_{G'}^-(x)$. Taking any pair of vertices y' in $R_{G'}^-(y)$ and z' in $R_{G'}^-(z)$ for the two vertices y and z of degree 3 in G' , we obtain that $\{u', v', x', y'\}$ is an independent set of G , contradicting $\alpha(G) = 3$. Secondly, suppose $d_{G'}(u) = d_{G'}(v) = 3$, and x, y are vertices with degree two in G' and $xy \notin E(G')$. Then, by Lemma 16(b), there exist at least two vertices x' and y' in $R_{G'}^-(x)$ and $R_{G'}^-(y)$, respectively, such that they are not adjacent to any of $R_{G'}^-(u) \cup R_{G'}^-(v)$. It follows that $\{x', y', u', v'\}$ is an independent set of G (where $u' \in R_{G'}^-(u)$ and $v' \in R_{G'}^-(v)$), contradicting $\alpha(G) = 3$. Finally, suppose that $d_{G'}(u) \neq d_{G'}(v)$, say, $d_{G'}(u) = 3, d_{G'}(v) = 2$. Take a pair of vertices $\{x, y\} \subseteq V_2(G') \setminus \{v\}$ such that $xy \notin E(G')$. Hence, by Lemma 16(b), there exists at least one vertex u' (say) of $R_{G'}^-(u)$ that is not adjacent to any vertex of $R_{G'}^-(x) \cup R_{G'}^-(y)$, and there is a vertex v' of $R_{G'}^-(v)$ that is not adjacent to any vertex of $\bigcup_{s \in V(G') \setminus \{v\}} R_{G'}^-(s)$. Taking $x' \in R_{G'}^-(x)$ and $y' \in R_{G'}^-(y)$, we obtain an independent set $\{x', y', v', u'\}$ of G , contradicting $\alpha(G) = 3$. This proves that G' has at most one nontrivial vertex if $G' \in \{K_{2,3}, K_{2,3}(1), K_{2,3}(2)\}$.

If G' has no nontrivial vertices, then $G = G'$ and G is isomorphic to one of those graphs shown in Fig. 1. Now, suppose that G' has exactly one nontrivial vertex u . Note that G is a simple graph. Hence, if $d_{G'}(u) = 2$, then, by Lemma 16(a), there is at least one vertex u' of $R_{G'}^-(u)$ such that u' is not adjacent to any vertex in $\bigcup_{x \in V(G') \setminus \{u\}} R_{G'}^-(x) = V(G') \setminus \{u\}$; if $d_{G'}(u) = 3$ then, by Lemma 16(b), for any pair of vertices in $N_{G'}(u)$, say s and t , there exists a vertex in $R_{G'}^-(u)$ which is adjacent to neither s nor t , since $|V(R_{G'}^-(u))| \geq 3$. Hence, $G' = K_{2,3}(2)$ has no nontrivial vertices. For $G' \in \{K_{2,3}, K_{2,3}(1)\}$, $R_{G'}^-(u)$ must be a complete graph. Suppose, in contrast, that there exist two nonadjacent vertices s, t in $R_{G'}^-(u)$. If $d_{G'}(u) = 2$, then the vertices s, t and the other two trivial vertices (that are not adjacent to u in G') of degree two in G' induce an independent set in G , contradicting $\alpha(G) = 3$. If $d_{G'}(u) = 3$, then there exists one vertex in $R_{G'}^-(u)$ that is not adjacent to any one of the three independent trivial vertices of degree two in G' , contradicting $\alpha(G) = 3$. Hence, $R_{G'}^-(u)$ is a complete graph. If $G' = K_{2,3}$ and $d_{G'}(u) = 3$, then $R_{G'}^-(u)$ is a triangle and the original graph G is isomorphic to $K'_{2,3}$ (otherwise, by a way similar to the proof of Lemma 16(b), we can prove that there exists a vertex in $R_{G'}^-(u)$ that is not adjacent to any one of the three nonadjacent vertices of degree two in G' , contradicting $\alpha(G) = 3$). If $G' = K_{2,3}(1)$, then $d_{G'}(u) = 2$ (otherwise, by Lemma 16(b) and by the fact that there exists a vertex of degree two that is not adjacent to u , there exists a vertex in $R_{G'}^-(u)$ that is not adjacent to any one of the three independent vertices of degree two in G' , contradicting $\alpha(G) = 3$). Note that there exists a vertex w of $R_{G'}^-(u)$ such that it is not adjacent to any vertex of $V(G') \setminus \{u\} (= \bigcup_{z \in V(G') \setminus \{u\}} R_{G'}^-(z))$. We have that any vertex of $N_{G'}(u)$ having degree three in G' (otherwise $\{w\} \cup (V_2(G') \setminus \{u\}) = \{w\} \cup (\bigcup_{z \in V_2(G') \setminus \{u\}} R_{G'}^-(z))$ is an independent set of four vertices in G , contradicting $\alpha(G) = 3$). This completes the proof of Theorem 3. \square

4. Proof of Theorem 6

Before presenting our proof of Theorem 6, we need some additional terminologies and notation. A subset S of $\mathcal{B}(G)$ is called a *branch cut* if the deletion of all edges and all inner vertices (of degree two) in any branch of S will induce more components than G has. A minimal branch cut is called a *branch-bond*. It is easily shown that, for a connected graph G , a subset S of $\mathcal{B}(G)$ is a branch-bond if and only if the deletion of all edges and all inner vertices (of degree two) in any branch of S will induce exactly two components. We denote by $\mathcal{BB}(G)$ the set of branch-bonds of G . A branch-bond $S \in \mathcal{BB}(G)$ is called *odd* if S consists of an odd number of branches. The *length* of a branch-bond $S \in \mathcal{BB}(G)$, denoted by $l(S)$, is the length of a shortest branch in it. Define $\mathcal{BB}_3(G) = \{S \in \mathcal{BB}(G) : |S| \geq 3 \text{ and } S \text{ is odd}\}$. Define $h_3(G) = \max\{l(S) : S \in \mathcal{BB}_3(G)\}$ if $\mathcal{BB}_3(G)$ is not empty; 0, otherwise. By $S \Delta T$ we denote the symmetric difference $(S \setminus T) \cup (T \setminus S)$. Note that, if H is an even subgraph of G and C is a cycle of G , then $G[E(H) \Delta E(C)]$ is also an even subgraph of G , but $G[E(H) \Delta E(C)]$ may have more components than H .

From a result of [15], one can easily obtain the following.

Theorem 17 ([17]). *Let G be a 2-connected graph. Then,*

$$h_3(G) - 1 \leq h(G) \leq h_3(G) + 1.$$

The following characterization of eulerian graphs involves branch-bonds.

Theorem 18 ([15]). *A connected graph is eulerian if and only if each branch-bond contains an even number of branches.*

The following two results will be used in the proof of Theorem 6.

Theorem 19 ([4]). *If G is a connected graph such that $\delta(G) \geq 3$, then*

$$h(G) \leq 2.$$

Lemma 20. *If H is an even subgraph of G , then every branch of $\mathcal{B}_H(G)$ lies in a cycle of H .*

Proof. It immediately follows from the definitions of a branch and an even subgraph. \square

Now, we present the proof of our second main result.

Proof of Theorem 6. By Theorems 1 and 4, we only need to consider the case when $t \geq 2$. Let $m = \lfloor \frac{2t+2}{3} \rfloor$. Then, $m \geq 2$. Since $h(G) \leq 2$ for the graph G with $\delta(G) \geq \kappa(G) \geq 3$ by Theorem 19, we can assume that $\kappa(G) = 2$. If $h_3(G) \leq m - 1$, then $h(G) \leq m$ by Theorem 17. So, it remains to consider the case when $h_3(G) \geq m$. Let \mathcal{B}_0 be a branch-bond in $\mathcal{B}_{\mathcal{B}_3}(G)$ such that $h_3(G) = \min\{|E(B)| : B \in \mathcal{B}_0\}$. Note that there are at least $\lfloor \frac{3h_3(G)}{2} \rfloor$ vertices in $G[\cup_{B \in \mathcal{B}_0} V(B)]$ forming an independent set of G , which contains at most one end vertex of \mathcal{B}_0 , and hence $\alpha(G) \geq \lfloor \frac{3h_3(G)}{2} \rfloor$. We claim that $h_3(G) \leq m + 1$ since otherwise $\alpha(G) \geq \lfloor \frac{3h_3(G)}{2} \rfloor \geq \frac{3h_3(G)-1}{2} \geq \frac{3m+5}{2}$, which contradicts the fact that

$$\alpha(G) \leq \kappa(G) + t = 2 + t < 2 + \frac{3m + 3 - 2}{2} = \frac{3m + 5}{2}. \tag{4.1}$$

We will prove that there is a subgraph $H \in EU_m(G)$ that implies that $h(G) \leq m = \lfloor \frac{2t+2}{3} \rfloor$ by Theorem 13.

Since the edgeless graph with vertex set $\cup_{i=3}^{\Delta(G)} V_i(G)$ satisfies (I) and (II), we may let H be a subgraph such that

- (1) H satisfies (I) and (II);
- (2) subject to (1), H contains a maximum number of branches of length at least $m + 2$;
- (3) subject to (1) and (2), $\max_{\emptyset \neq V(H_1) \subseteq V(H)} d_G(H_1, H - H_1)$ is minimized;
- (4) subject to the above, H contains a minimum number of induced subgraphs F for which $d_G(F, H - F) = \max_{\emptyset \neq V(H_1) \subseteq V(H)} d_G(H_1, H - H_1)$.

We claim that $H \in EU_m(G)$. It suffices to prove that H satisfies (III) and (IV) since G is 2-connected. It suffices to prove the following two claims.

Claim 4. $d_G(H_1, H - H_1) \leq m - 1$ for any subgraph H_1 of H .

Proof of Claim 4. Assume, in contrast, that there is an induced subgraph F of H with $d_G(F, H - F) \geq m \geq 2$, then any shortest path between F and $H - F$ is a branch of G and hence both F and $H - F$ are even subgraphs of G (by the condition (II) of the definition of H and by the fact that the shortest path has length at least two). Since G is 2-connected, there are at least two branches P_1, P_2 between F and $H - F$ and we may hence choose a cycle C containing both P_1 and P_2 such that the length of a longest branch in $\mathcal{B}_C(G) \cap \mathcal{B}_H(G)$ is as small as possible. Hence, $|E(P_i)| \geq m$. By the choice of C , we have the following facts.

Fact 4.1. If B is a longest branch in $\mathcal{B}_C(G) \cap \mathcal{B}_H(G)$, then $|E(B)| \geq m$ and there exists at least one branch $B' \in \mathcal{B}_H(G) \setminus \mathcal{B}_C(G)$ of length at least $|E(B)|$ such that the end vertices of both B and B' are all in the same one of $\{H, H - F\}$.

Proof of Fact 4.1. We prove the first statement. Assume, in contrast, that $|E(B)| < m$. Then the subgraph H' obtained from the subgraph $G[E(H) \Delta E(C)]$ by adding the remaining vertices of $\cup_{i=3}^{\Delta(G)} V_i(G)$ as isolated vertices of H' is a subgraph satisfying (1) and (2) such that $\max_{\emptyset \neq V(H_1) \subseteq V(H')} d_G(H_1, H' - H_1)$ is less than $\max_{\emptyset \neq V(H_1) \subseteq V(H)} d_G(H_1, H - H_1)$, contradicting the condition (3) of the choice of C .

We then prove the second statement. Let B be a longest branch of $\mathcal{B}(G) \cap \mathcal{B}_H(G)$. By symmetry, say, $B \subseteq F$. By Lemma 20, there exists a cycle C' of F such that $B \subseteq C'$ and C' contains a branch $B' \in \mathcal{B}_F(G) \setminus \mathcal{B}_C(G)$ of length at least $|E(B)|$, since otherwise the subgraph H' obtained from $G[E(H) \Delta E(C')]$ by adding the remaining vertices of $\cup_{i=3}^{\Delta(G)} V_i(G)$ as isolated vertices of H' is a subgraph satisfying (1) and (2) such that $\max_{\emptyset \neq V(H_1) \subseteq V(H')} d_G(H_1, H' - H_1)$ is less than $\max_{\emptyset \neq V(H_1) \subseteq V(H)} d_G(H_1, H - H_1)$, contradicting the condition (3) of the choice of C . This completes the proof of Fact 4.1. \square

Fact 4.2. There are no pair of branches of length at least $m + 2$ such that all of their end vertices belong to the same one of $\{F, H - F\}$.

Proof of Fact 4.2. Assume, in contrast, that B_1, B_2 are two branches of length at least $m + 2$ such that all of their end vertices belong to the same one of $\{F, H - F\}$. Then there exist an independent set containing at least $\lfloor \frac{m+2}{2} \rfloor + 2 \lfloor \frac{m}{2} \rfloor$ inner vertices in the four branches B_1, B_2, P_1, P_2 and hence $\alpha(G) \geq 2 \lfloor \frac{m+2}{2} \rfloor + 2 \lfloor \frac{m}{2} \rfloor + 1 \geq \frac{3m+5}{2}$ by Fact 4.1, which contradicts (4.1). This completes the proof of Fact 4.2. \square

Then we have the following fact.

Fact 4.3. There are at least two branches B_1, B_2 of length m or $m + 1$ that are in $\mathcal{B}(G) \setminus \{P_1, P_2\}$ such that $E(B_i) \in E(C)$ and all end vertices of both B_1 and B_2 are in the same one of $\{F, H - F\}$.

Proof of Fact 4.3. Otherwise, by Facts 4.1 and 4.2, there is at most one branch B_1 of length m or $m + 1$ that is in $\mathcal{B}(G) \setminus \{P_1, P_2\}$ such that $E(B_1) \in E(C)$ and both endvertices of B_1 are in the same one of $\{F, H - F\}$. Then the subgraph H' obtained from the subgraph $G[E(H) \Delta E(C)]$ by adding the remaining vertices of $\cup_{i=3}^{\Delta(G)} V_i(G)$ as isolated vertices of H' is a subgraph satisfying (1) and (2) such that either $\max_{\emptyset \neq V(H_1) \subseteq V(H')} d_G(H_1, H' - H_1)$ is fewer than $\max_{\emptyset \neq V(H_1) \subseteq V(H)} d_G(H_1, H - H_1)$ or H' has a fewer

number of induced subgraphs F for which $d_G(F, H' - F) = \max_{\emptyset \neq V(H_1) \subseteq V(H')} d_G(H_1, H' - H_1)$ than H (in this case, H' satisfies (3)), which shows that H does not satisfy the condition (4) of the choice of C , a contradiction. This completes the proof of Fact 4.3. \square

Since G is 2-connected, every component of H has at least one vertex which is not adjacent to any inner vertex of branches of length at least m in $\mathcal{B}_C(G)$. Hence, by the definition of odd branch-bonds, there are at least two branches of length m or $m + 1$ in $\mathcal{B}(G) \setminus \mathcal{B}_0$. Hence, by Fact 4.3, there is an independent set containing at least $2\lfloor \frac{m}{2} \rfloor + 1 \geq 3$ vertices, any element of which is not adjacent to any inner vertex of any branch of \mathcal{B}_0 , and so $\alpha(G) \geq \lfloor \frac{3h_3(G)}{2} \rfloor + 3 \geq \frac{3m-1}{2} + 3 \geq \frac{3m+5}{2}$, which contradicts (4.1). This proves Claim 4. \square

Claim 5. $|E(B)| \leq m + 1$ for any branch $B \in \mathcal{B}(G) \setminus \mathcal{B}_H(G)$.

Proof of Claim 5. Assume, in contrast, that there is a branch $B_0 \in \mathcal{B}(G) \setminus \mathcal{B}_H(G)$ with $|E(B_0)| \geq m + 2$. Let u and v be two end vertices of B_0 and $S(u, B_0)$ be a branch-bond containing B_0 such that any branch of $S(u, B_0)$ has u as an end vertex. Since G is 2-connected, $|S(u, B_0)| \geq 2$.

By the following algorithm, we first find a cycle of G that contains B_0 and then obtain a contradiction.

Algorithm B_0 . 1. If $|S(u, B_0)|$ is even, then select a branch $B_1 \in S(u, B_0) \setminus (\mathcal{B}_H(G) \cup \{B_0\})$ by Theorem 18. Otherwise, since $|E(B_0)| \geq m + 2$, (we can) select a branch $B_1 \in S(u, B_0)$ with

$$|E(B_1)| = l(S(u, B_0)) \leq h_3(G) \leq m + 1$$

(obviously $B_1 \neq B_0$) and let $u_1 (\neq u)$ be the other end vertex of B_1 . If $u_1 = v$, then set $t := 1$ and stop. Otherwise $i := 1$.

2. Select a branch-bond $S(u, u_i, B_0)$ in G which contains B_0 but not any of $\{B_1, B_2, \dots, B_i\}$ such that any branch in $S(u, u_i, B_0)$ has exactly one end vertex in $\{u, u_1, u_2, \dots, u_i\}$. If $|S(u, u_i, B_0)|$ is even, then, by Theorem 18, (we can) select a branch

$$B_{i+1} \in S(u, u_i, B_0) \setminus (\mathcal{B}_H(G) \cup \{B_0\}).$$

Otherwise, since $|E(B_0)| \geq h_3(G) + 1$, (we can) select a branch $B_{i+1} \in S(u, u_i, B_0)$ such that

$$|E(B_{i+1})| = l(S(u, u_i, B_0)) \leq h_3(G)$$

(obviously $B_{i+1} \neq B_0$), and let u_{i+1} be the end vertex of B_{i+1} that is not in $\{u, u_1, u_2, \dots, u_i\}$.

3. If $u_{i+1} = v$, then set $t := i + 1$ and stop. Otherwise replace i by $i + 1$ and return to step 2.

Note that $|\mathcal{B}(G)|$ is finite, and $d_G(v) \geq 2$ implies that Algorithm B_0 will stop after a finite number of steps. Note that $G[\bigcup_{i=0}^t E(B_i)]$ is connected. Furthermore, since $u_t = v$ and $|S(u, u_i, B_0)| \geq 2$, $G[\bigcup_{i=0}^t E(B_i)]$ has a cycle C_0 of G which contains B_0 . Let H' be the subgraph of G obtained from $G[E(H) \Delta E(C_0)]$ by adding the remaining vertices of $\bigcup_{i=3}^{\Delta(G)} V_i(G)$ as isolated vertices in H' . By the choice of B_i , $|E(B)| \leq h_3(G) \leq m + 1$ for $B \in \mathcal{B}_H(G) \cap \{B_1, B_2, \dots, B_t\}$. Hence, since B_0 is in $\mathcal{B}_{C_0}(G)$, H' is a subgraph with the conditions (I) and (II) that has fewer branches of length $m + 2$ than H , which shows that H does not satisfy (2), a contradiction. This proves Claim 5. \square

Thus we, in fact, have proved that G has a subgraph in $EU_m(G)$ by Claims 4 and 5. Hence, $h(G) \leq m = \lfloor \frac{2t+2}{3} \rfloor$ by Theorem 13. This also completes the proof of Theorem 6. \square

5. Sharpness

In this section, we discuss sharpness. We show that Theorem 6 is sharp by describing a graph satisfying the conditions of Theorem 6 and having the Hamiltonian index of exactly $\lfloor \frac{2t+2}{3} \rfloor$. By Theorems 1, 3 and 12, we only need to show the sharpness in the case when $t \geq 2$.

Let H_1, H_2 be two vertex-disjoint complete graphs of order at least four and P_1, P_2, P_3 be three vertex-disjoint paths with length $m + 1 \geq 3$. Now, obtain the graph G_0 by identifying three end vertices of P_1, P_2, P_3 with three vertices of H_1 , and identifying the other end vertex of P_i with three vertices of H_2 , respectively; see Fig. 3.

Let C be a longest cycle of G_0 . Then, $C \in EU_m(G_0)$. So $h(G_0) \leq m$ by Theorem 13, and $h(G_0) \geq h_3(G_0) - 1 = m$ by Theorem 17. Hence, $h(G_0) = m$. Note that if m is even, then $t = \alpha(G_0) - 2 = \lfloor \frac{3m}{2} \rfloor$ and $h(G_0) = m = \lfloor \frac{2t+2}{3} \rfloor$. For the case when m is odd, consider the reduction G'_0 of G_0 , which is obtained from G_0 by contracting H_1 and H_2 , respectively. Then $t = \alpha(G'_0) - 2 = \lfloor \frac{3m}{2} \rfloor$ and $h(G'_0) = m = \lfloor \frac{2t+2}{3} \rfloor$ for m is odd.

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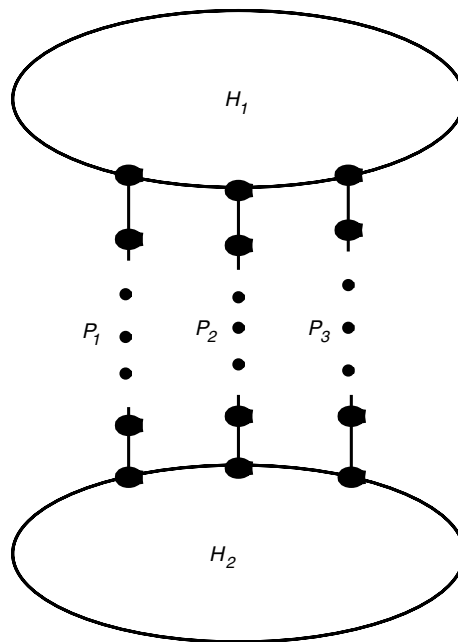


Fig. 3. The graph G_0 .

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