

# Hamiltonian $N_2$ -Locally Connected Claw-Free Graphs

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**Abstract:** A graph  $G$  is  $N_2$ -locally connected if for every vertex  $v$  in  $G$ , the edges not incident with  $v$  but having at least one end adjacent to  $v$  in  $G$  induce a connected graph. In 1990, Ryjáček conjectured that every 3-connected  $N_2$ -locally connected claw-free graph is Hamiltonian. This conjecture is proved in this note. © 2004 Wiley Periodicals, Inc. *J Graph Theory* 48: 142–146, 2005

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## 1. INTRODUCTION

We use [1] for terminology and notations not defined here, and consider finite simple graphs only. Let  $G$  be a graph. Denote by  $d_G(v)$  the degree of a vertex

$v \in V(G)$ . For a vertex  $v$  of  $G$ , the neighborhood of  $v$ , that is, the induced subgraph on the set of all vertices that are adjacent to  $v$ , will be called the neighborhood of the *first type* of  $v$  in  $G$  and denoted by  $N_1(v, G)$ , or briefly,  $N_1(v)$  or  $N_G(v)$ . For notational convenience, we shall use  $N_G(v)$  to denote both the induced subgraph and the set of vertices adjacent to  $v$  in  $G$ . We define the neighborhood of the *second type* of  $v$  in  $G$  (denoted by  $N_2(v, G)$ , or briefly,  $N_2(v)$ ) as the subgraph of  $G$  induced by the edge subset  $\{e = xy \in E(G) : v \notin \{x, y\} \text{ and } \{x, y\} \cap N(v) \neq \emptyset\}$ . We say that a vertex  $v$  is *locally connected* if  $N(v)$  is connected; and  $G$  is *locally connected* if every vertex of  $G$  is locally connected. Analogously, a vertex  $v$  is  *$N_2$ -locally connected* if its second type neighborhood  $N_2(v)$  is connected; and  $G$  is called  *$N_2$ -locally connected* if every vertex of  $G$  is  $N_2$ -locally connected. It follows from the definitions that every locally connected graph is  $N_2$ -locally connected. A graph  $G$  is *claw-free* if it does not contain  $K_{1,3}$  as an induced subgraph. The following theorems give the hamiltonicity of a locally and  $N_2$ -locally connected graph.

**Theorem 1.1** (Oberly and Sumner, [6]). *Every connected locally connected claw-free graph on at least three vertices is hamiltonian.*

**Theorem 1.2** (Ryjáček, [8]). *Let  $G$  be a connected,  $N_2$ -locally connected claw-free graph without vertices of degree 1, which does not contain an induced subgraph  $H$  isomorphic to either  $G_1$  or  $G_2$  (Fig. 1) such that  $N_1(x, G)$  of every vertex  $x$  of degree 4 in  $H$  is disconnected. Then  $G$  is Hamiltonian.*

We say that  $G$  is *vertex pancyclic* if it contains cycles of every length through every vertex.

**Theorem 1.3** (Li, [4]). *Let  $G$  be a connected,  $N_2$ -locally connected claw-free graph with  $\delta(G) \geq 3$ , which does not contain an induced subgraph  $H$  isomorphic to either  $G_1$  or  $G_2$  (Fig. 1). Then  $G$  is vertex pancyclic.*

In another recent paper [5], Li suggested a new relaxation of the locally connectedness condition for Hamiltonian claw-free graphs. The main purpose of this note is to prove the following theorem, conjectured by Ryjáček in [8].

**Theorem 1.4.** *Every 3-connected  $N_2$ -locally connected claw-free graph is Hamiltonian.*

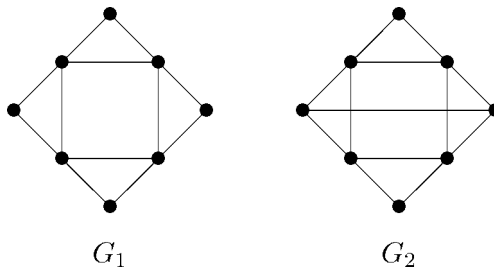


FIGURE 1.

## 2. PROOF OF THE MAIN THEOREM

Our approach is to firstly apply the line graph closure (invented by Ryjáček in [8]) to convert the problem into a line graph problem. Then apply techniques in supereulerian graphs to solve the corresponding line graph problem.

The *line graph* of a graph  $G$ , denoted by  $L(G)$ , has  $E(G)$  as its vertex set, where two vertices in  $L(G)$  are adjacent if and only if the corresponding edges in  $G$  have a vertex in common. Let  $G$  be the line graph  $L(H)$  of a graph  $H$ . If  $L(H)$  is  $k$ -connected, then  $H$  is *essentially  $k$ -edge-connected*, which means that the only edge-cut sets of  $H$  having less than  $k$  edges are the sets of edges incident with some vertex of  $H$ .

In [9], Ryjáček defined the closure  $cl(G)$  of a claw-free graph  $G$  by recursively completing the neighborhood of any locally connected vertex of  $G$ , as long as this is possible. The closure  $cl(G)$  is a well-defined claw-free graph and its connectivity is at least equal to the connectivity of  $G$ . The *circumference* of  $G$  is the length of a longest cycle in  $G$ .

**Theorem 2.1** (Ryjáček, [9]). *Let  $G$  be a claw-free graph and  $cl(G)$  its closure. Then*

- (i) *there is a triangle-free graph  $H$  such that  $cl(G)$  is the line graph of  $H$ ,*
- (ii) *both graphs  $G$  and  $cl(G)$  have the same circumference.*

Let  $O(G)$  denote the set of all vertices in  $G$  with odd degree. A graph  $G$  is *eulerian* if  $O(G) = \emptyset$  and  $G$  is connected. A spanning closed trail of  $G$  is called a *spanning Eulerian subgraph* of  $G$ . A subgraph  $H$  of  $G$  is *dominating* if  $G - V(H)$  is edgeless. If a closed trail  $C$  of  $G$  satisfies  $E(G - V(C)) = \emptyset$ , then  $C$  is called a *dominating Eulerian subgraph*.

**Theorem 2.2** (Harary and Nash-Williams, [2]). *The line graph  $G = L(H)$  of a graph  $H$  is Hamiltonian if and only if  $H$  has a dominating Eulerian subgraph.*

Theorem 2.2 reveals the relationship between a dominating Eulerian subgraph in  $H$  and a Hamiltonian cycle in  $L(H)$ .

Theorem 2.3 below provides a sufficient condition for a graph to have a spanning Eulerian subgraph (therefore a dominating Eulerian subgraph), which is originally conjectured by Paulraja ([7]).

**Theorem 2.3** (Lai, [3]). *Let  $G$  be a 2-connected graph with  $\delta(G) \geq 3$ . If every edge of  $G$  is in an  $m$ -cycle of  $G$  ( $m \leq 4$ ), then  $G$  has a spanning Eulerian subgraph.*

**Lemma 2.4.** *Let  $G$  be an  $N_2$ -locally connected graph and let  $x$  be a locally connected vertex of  $G$  such that  $G[N_G(x)]$  is not complete. Let  $N' = \{uv : u, v \in N_G(x), uv \notin E(G)\}$  and let  $G'$  be the graph with vertex set  $V(G') = V(G)$  and with edge set  $E(G') = E(G) \cup N'$ . Then  $G'$  is  $N_2$ -locally connected.*

**Proof.** Let  $w \in V(G')$ . If  $w = x$ , then  $N_2(w, G')$  is connected since  $N_{G'}(x)$  is complete. So we may assume that  $w \neq x$ . Since  $G$  is  $N_2$ -locally connected,  $N_2(w, G)$  is connected. If  $E(N_2(w, G')) - E(N_2(w, G)) = \emptyset$ , then  $E(N_2(w, G')) = E(N_2(w, G))$  and  $N_2(w, G')$  is connected. Thus we assume that  $E(N_2(w, G')) - E(N_2(w, G)) \neq \emptyset$ . Let  $e = uv \in E(N_2(w, G')) - E(N_2(w, G))$ . Since  $e = uv \in E(N_2(w, G'))$ , we have  $w \notin \{u, v\}$ , and so  $uv \in E(G')$ . Without loss of generality, we assume that  $wu \in E(G')$ .

**Case 1.**  $uv \in E(G)$ .

By  $e = uv \notin E(N_2(w, G))$ , we have  $wu, wv \notin E(G)$ . Since  $wu \in E(G')$  by the assumption,  $w, u \in N_G(x)$ . So  $xu \in E(N_2(w, G))$ . Therefore adding a new edge  $uv$  to  $N_2(w, G)$  does not change its connectivity, and so  $N_2(w, G')$  is connected.

**Case 2.**  $uv \notin E(G)$ .

Since  $uv \in E(G')$ , we have  $u, v \in N_G(x)$ . If  $w \in N_G(x)$ , then  $xu, xv \in E(N_2(w, G))$ . Thus adding a new edge  $uv$  to  $N_2(w, G)$  does not change its connectivity, and so  $N_2(w, G')$  is connected. If  $w \notin N_G(x)$ , then we have  $wu \in E(G)$  since  $wu \in E(G')$  (otherwise,  $w \in N_G(x)$ , a contradiction). Thus  $xu \in E(N_2(w, G))$ . So adding a new edge  $uv$  to  $N_2(w, G)$  does not change its connectivity, and therefore  $N_2(w, G')$  is connected. ■

**Proof of Theorem 1.4.** By Theorem 2.1(ii), the graph  $G$  is Hamiltonian if and only if its closure  $cl(G)$  is Hamiltonian. By Lemma 2.4 and as  $cl(G)$  is both 3-connected and claw free, the graph  $cl(G)$  is also a 3-connected  $N_2$ -locally connected claw-free graph. By Theorem 2.1, we may assume that for a triangle-free graph  $H$ ,  $G = cl(G) = L(H)$ .

An edge  $e = uv$  is called a *pendant edge* if either  $d_G(u) = 1$  or  $d_G(v) = 1$ .

**Claim 1.** Let  $e = uv \in E(H)$ . If  $e$  is not a pendant edge, then  $e$  is in some 4-cycle of  $H$ .

**Proof.** Since  $H$  is triangle free, we have  $N_H(u) \cap N_H(v) = \emptyset$ . Let  $v_e \in V(G)$  corresponds to the edge  $e \in E(H)$  in terms of the line graph. Since  $e$  is not a pendant edge and  $G$  is claw free,  $N_G(v_e)$  is the union of two disjoint cliques. Suppose they are  $L_1, L_2$ . Since  $G$  is 3-connected, there exists at least one path  $w_1 w_2 \cdots w_n$  which is edge disjoint with  $G[V(L_1) \cup V(L_2) \cup \{v_e\}]$  in  $G - v_e$  with  $w_1 \in V(L_1), w_n \in V(L_2)$ . Since  $G$  is  $N_2$ -locally connected, we have that  $n = 3$ . Thus  $v_e w_1 w_2 w_3 v_e$  is an induced 4-cycle of  $G$ , which corresponds to a 4-cycle in  $H$  containing  $e$ . ■

Let  $\tilde{H}$  be the graph obtained from  $H$  by deleting the vertices of degree 1 or 2 and replacing each path  $xyz$  in  $H$  with  $d_H(y) = 2$  by an edge  $xy$ . Then it is straightforward to verify the following claim.

**Claim 2.** If  $\tilde{H}$  has a spanning Eulerian subgraph, then  $H$  has a dominating Eulerian subgraph.

Since  $G$  is 3-connected,  $\tilde{H}$  is 3-edge-connected. Let  $B$  be an arbitrary block of  $\tilde{H}$ . Since  $\tilde{H}$  is 3-edge-connected,  $\delta(B) \geq 3$ . By Claim 1, every edge of  $B$  lies in a cycle of  $B$  of length at most 4. By Theorem 2.3 and since  $B$  is 2-connected,  $B$  has a spanning Eulerian subgraph. Since every block of  $\tilde{H}$  has a spanning Eulerian subgraph,  $\tilde{H}$  has a spanning Eulerian subgraph. By Claim 2,  $H$  has a dominating Eulerian subgraph. By Theorem 2.2,  $cl(G)$  is Hamiltonian. ■

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