

On the edge arboricity of a random graph

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Abstract

The edge arboricity $a(G)$ of a graph G is the minimum number of acyclic subgraphs whose union covers the edge set of G . In this note we show that if the edge probability is given by $p^3 n = c \log n$, then almost every graph has

$$a(G) = \left\lceil \frac{|E(G)|}{n-1} \right\rceil$$

provided the constant c is sufficiently large.

Dedicated to Roger Entringer on the occasion of his 60th birthday

1 Introduction

The edge arboricity $a(G)$ of a graph G is minimum number of acyclic subgraphs whose union covers the edge set of G . Nash-Williams [Na64] proved that

$$a(G) = \max_{H \subseteq G} \left\lceil \frac{|E(H)|}{|V(H)| - 1} \right\rceil \quad (1.1)$$

where the maximum runs over all non-trivial induced subgraphs H of G . The first two authors showed [CaC91] that when the edge

probability p is fixed, almost all graphs G have the property that $|E(H)|/(|V(H)| - 1)$ attains its maximum in (1.1) if and only if $G = H$. Following closely the method of [CaC91], we will extend that result for $p = p(n) \rightarrow 0$.

Our sample space consists of all labeled graphs G with n vertices. The vertex set of G is $V(G) = \{1, 2, \dots, n\}$ and the edge set is $E(G)$. Given the edge probability $0 < p < 1$, the probability of a graph G with M edges is defined by

$$P(G) = p^M(1 - p)^{N-M} \quad (1.2)$$

where $N = \binom{n}{2}$, the number of slots available for edges. Thus the sample space consists of Bernoulli trials and the edges are selected independently with probability p . Suppose \mathcal{Q} is a set of graphs of order n with some specified property Q . If the probability $P(\mathcal{Q})$ approaches 1 as n goes to infinity, then we say that *almost all graphs have property Q or the random graph has property Q a.s.* (almost surely).

For background material and notation not provided here one can consult the introductory book on random graphs [Pa85] and for the strongest and many of the most recent results we use the extensive and comprehensive treatise [Bo85].

2 Edge arboricity

For any non-trivial, connected graph G of order n , define

$$\gamma(G) = \max_{H \subseteq G} \frac{|E(H)|}{|V(H)| - 1}, \quad (2.1)$$

where the maximum is taken over all non-trivial subgraphs H of G . We use the following elementary inequalities frequently:

$$\frac{|E(G)|}{n-1} \leq \gamma(G) \leq \lceil \gamma(G) \rceil \leq a(G). \quad (2.2)$$

Let $\mathcal{F}(G)$ be the family of non-trivial subgraphs H of G such that

$$\gamma(G) = \frac{|E(H)|}{|V(H)| - 1}. \quad (2.3)$$

Thus these graphs H achieve the maximum value in (2.1) and it is also easy to see that $\gamma(H) = \gamma(G)$.

Theorem 2.1. With edge probability defined by $p^3 n = c \log n$, if the constant c is at least 28, almost surely $\mathcal{F}(G) = \{G\}$ and hence the edge arboricity is

$$a(G) = \left\lceil \frac{|E(G)|}{n-1} \right\rceil. \quad (2.4)$$

Proof: Suppose G is any connected graph of order $n > 1$. Let H be in the family $\mathcal{F}(G)$ and set $r = |V(H)|$. First we find a lower bound for r in terms of the number of edges of G . By the definition of H we have

$$\gamma(H) = \gamma(G) = \frac{|E(H)|}{r-1}. \quad (2.5)$$

Now we combine (2.2) and (2.5) and use the fact that H has order r to obtain .

$$r = \frac{2}{r-1} \binom{r}{2} \geq \frac{2}{r-1} |E(H)| = 2\gamma(G) \geq 2 \frac{|E(G)|}{n-1}. \quad (2.6)$$

Next we can use Chebyshev's inequality to derive an approximation for the number of edges in a random graph G from which we can determine a lower bound for $|E(G)|$. See, for example, a special case in exercise 3.1.2 of [Pa85]. For a slightly more general result, we have the following. For any positive sequence $\varepsilon_n \rightarrow 0$,

$$|E(G)| \geq p \binom{n}{2} (1 - \varepsilon_n), \quad (2.7)$$

provided that $\varepsilon_n^2 p n^2 \rightarrow \infty$.

By hypothesis our edge probability is well beyond the threshold for connectedness (see [Bo85] or [Pa85]) so we can assume that almost all graphs are connected.

Combining (2.6) and (2.7) we observe that for almost all graphs, the number r of vertices in a graph H from the family $\mathcal{F}(G)$ satisfies

$$r \geq p n (1 - \varepsilon_n), \quad (2.8)$$

provided that the condition in (2.7) on ε_n is satisfied.

At this point we need an estimate for the number of edges in H . Using Theorem 8, p. 44 of [Bo85], we can conclude that $|E(H)|$ is almost surely quite close to $p \binom{r}{2}$. In particular, we can conclude that

$$\gamma(H) = \frac{|E(H)|}{r-1} \leq \frac{r}{2} \left\{ p + \left(\frac{7p \log n}{r} \right)^{1/2} \right\}, \quad (2.9)$$

almost surely, provided that

$$r \geq (252/p) \log n. \quad (2.10)$$

And this latter condition will be met if the lower bound in (2.8) exceeds the right side of (2.10), i.e. we just need

$$pn(1 - \varepsilon_n) \geq (252/p) \log n. \quad (2.11)$$

On solving this equation for p , we find that all required conditions on p are met if p is defined as in the hypothesis.

Now we are ready to compare n and r by using the lower bound on $\gamma(G)$ in (2.6) and (2.7) and the upper bound on $\gamma(H)$ from (2.9). Since $\gamma(H) = \gamma(G)$, we have

$$\frac{r}{2} \left\{ p + \left(\frac{7p \log n}{r} \right)^{1/2} \right\} \geq \frac{pn}{2} (1 - \varepsilon_n). \quad (2.12)$$

On substituting the expression from the hypothesis for p in this inequality, after a few steps we find that

$$n - r \leq c_0 (n \log n / p)^{1/2}, \quad (2.13)$$

for suitable ε_n and where c_0 is a constant greater than $\sqrt{7}$.

Now suppose that there is a vertex v of G that is not also in H . We are going to find an upper bound for the degree of v in G that is too far from the mean to hold for almost all graphs. This will imply that such vertices almost surely do not exist. Define H_v to be the subgraph of G induced by v together with the r vertices of H . By the definition in (2.1),

$$|E(H_v)| \leq \gamma(H_v)r. \quad (2.14)$$

But since H achieves the maximum value in (2.1),

$$\gamma(H_v) \leq \gamma(H). \quad (2.15)$$

Combining (2.14), (2.15) and (2.5), we have

$$|E(H_v)| \leq \gamma(H_v)r \leq \gamma(H)r = |E(H)| + \gamma(H). \quad (2.16)$$

This implies that the degree of v in H is at most $\gamma(H)$, and hence the degree of v in G is at most $n - r + \gamma(H)$, i.e., almost surely

$$\deg_G v \leq n - r + \gamma(H). \quad (2.17)$$

Using the bounds in (2.9) and (2.13), we find

$$\deg_G v \leq c_0(n \log n/p)^{1/2} + \frac{r}{2} \left\{ p + \left(\frac{7p \log n}{r} \right)^{1/2} \right\}. \quad (2.18)$$

And after a bit of work on the right side of (2.18) in which the value of c_0 depends on $c > 28$, we have

$$\deg_G v \leq (1 - \varepsilon)pn, \quad (2.19)$$

for large n and sufficiently small $\varepsilon > 0$.

This last inequality contradicts a theorem of Erdős and Rényi which states that if $pn/\log n \rightarrow \infty$, then almost surely the degrees of all vertices satisfy

$$(1 - \varepsilon)pn < \deg_G v < (1 + \varepsilon)pn. \quad (2.20)$$

where $\varepsilon > 0$ is arbitrary (see p. 66 of [Pa85]). //

We suspect that the theorem gives the right value for the edge arboricity for much lower edge probabilities but the family $\mathcal{F}(G)$ may not consist of G alone.

3 Tree packing number

The *tree packing number* $t(G)$ of a connected graph G is the maximum number of edge-disjoint spanning trees contained in G . It can be used as a measure of network vulnerability and is closely related to the edge arboricity $a(G)$. And the same method of [CaC91] can be applied here with the same result. Tutte [Tu61] and Nash-Williams [Na61] proved that

$$t(G) = \lfloor \eta(G) \rfloor, \quad (3.1)$$

where

$$\eta(G) = \min_{E \subseteq E(G)} \frac{|E|}{c(G - E) - 1} \quad (3.2)$$

and $c(G - E)$ is the number of components of $G - E$.

For any graph satisfying $\mathcal{F}(G) = \{G\}$, we always have $\gamma(G) = \eta(G)$ (see [CaGHL92]). But it can be shown that almost surely $\gamma(G)$ is not an integer and hence random graphs for which $\mathcal{F}(G) = \{G\}$, have

$$a(G) = t(G) + 1. \quad (3.3)$$

Of course, we only have found the values of these packing and covering numbers for random graphs when p is defined as in Theorem 2.1.

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